LINCOS

DESIGN OF A LANGUAGE FOR COSMIC INTERCOURSE

PART I

Dr. HANS FREUDENTHAL

Professor of Mathematics
University of Utrecht



1960

No part of this book may be reproduced in any form, by print, photoprint, microfilm or any other means without written permission from the publisher.



aug 2025 / worcester, ma / madrid, spain
salitter.org @salitters instagram
thank
 you

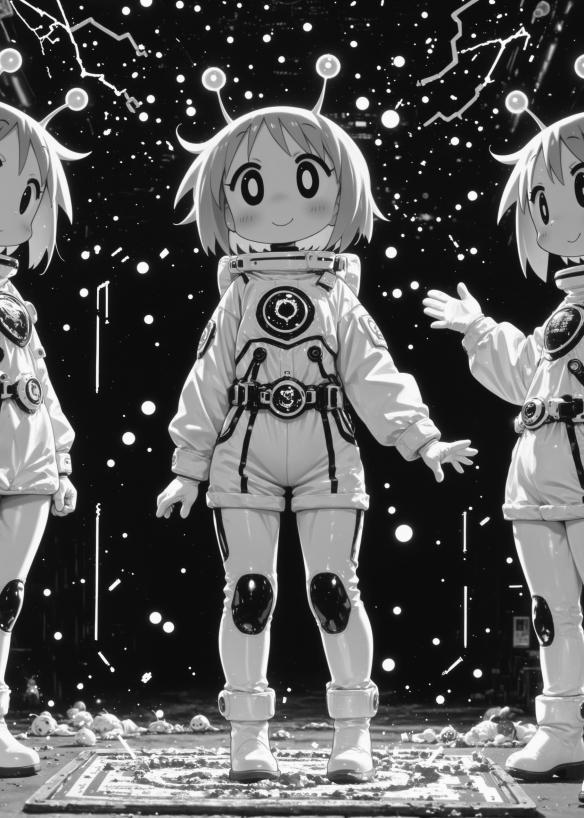












INTRODUCTION

001. Scientists, artists and artisans tend to develop a terminology of their own. They use common language as a vernacular that will be enriched, impoverished, and modified in order to serve special purposes. The transformation may affect not only the vocabulary but also the syntax of the vernacular, though essential syntactical modifications are rather unusual.

It is certain, though far from generally admitted, that the meaning of a linguistic term should be determined by its contexts, it is a matter of fact that there cannot be any reasonably uniform opinion about the meaning of a word, if people cannot agree about the truth of the majority of contexts in which the word occurs. It is a common historical feature of many sciences, that their first representatives tried to create a terminology before the stock of known facts was large enough to provide a sufficiently large context for the elements of that terminology. Even now this is a serious drawback for philosophy and the arts. When people do not understand each other, this is usually said to be because 'they speak different languages'. This may be correct as long as "understanding" means a mere linguistic phenomenon. If understanding means intelligence, one may posit the inverse thesis with at least as much justification: People speak different languages because they do not understand each other. (Note that the word language has here a rather unusual meaning.)

As long as there is no reasonable agreement about facts and terminology among those who work in a given field, it seems wise to stick as closely as possible to the vernacular, and not to create any new term until it has become possible to define it in a satisfactory way by a sufficiently rich context. In any case, one should try to abstain from playing upon words as contexts.

In the following pages I shall use common English as a vernacular, and I shall endeavour not to use technical terms that are not generally agreed upon or else to provide for a sufficiently rich context or even something of a definition if I use technical terms of my own. So verbs like "to designate" and "to mean" will have the same meaning as they have in the Vernacular and not that given them in recent publications which might be characterized by the term "formalist semantics". The use of a few words will be explained below.

- 0 02. The meaning of the word "language" is fairly unambiguous as long as there is no reference to any special language. Generally people agree that somebody who uses English words when he explains some astronomical experiments, is speaking English, but sometimes, especially if one wishes to stress some divergences between the vernacular and the more technical idiom, one says that he uses a special astronomical language. Though borderlines may always be disputed (one may quarrel about the question whether as early as 1250 the monk Robert of Gloucester wrote English and whether nowadays "Beach-la-mar" is still English) "language" in the first sense is one of the best defined notions of contemporary English, whereas in the second sense it is extremely vague. So I will provisionally stick to the customary unsophisticated use of the word "language".
- 0 03. We shall have occasions to speak of natural languages. This is a term which may give rise to misunderstanding, as it seems to evoke the obsolete idea of one natural language in the dawn of mankind, as the germ of our present languages. When nature is opposed to society, then language is rather a social then a natural phenomenon. "Natural language", however, as meant by us, is opposed to "artificial language", and as the best-known example of such an "unnatural" language, we may mention Esperanto. Of course natural languages are not absolutely natural, any more than artificial languages are absolutely artificial. Sanskrit is a noteworthy transitional form, though it is perhaps more artificial than natural, but even in our present-day languages conscious influences have directed the spontaneous growth. Orthography in particular has often been the object of attack and the source of a conscious linguistic policy. On the other hand spontaneous growth has been observed in Esperanto, the only living artificial language.
- 0 04. When dealing with artificial languages we must decide whether codes are to be counted among languages. First of all we shall explain the sense in which we will use the word "code". In the European languages "code" may denote a collection of laws of jurisprudence, and in some of them "code" is also the system of rules of etiquette. This circumstance may give rise to a terminology in which codes are distinguished from languages by their highly conventional character. In this terminology every artificial language would be a code.

I shall not accept this terminology. I shall distinguish codes from languages not by a genetic, but rather by a formal criterion. If the meaning of a message is to be kept secret, sender and receiver may use a language that is unknown to other people. Inventing and using brand-new languages is difficult. So they will try to parasitize on a common language. By a well-

chosen system of superficial transformations they will change that common language into something like a language, that can only be understood by people who are acquainted with the system of transformation rules. This collection of rules is a code in the proper sense (as a juridical code is a collection of juridical rules). If a code X is employed to transform plain English into a less understandable text, we say that this text is written "in Code X", just as the original text was written "in English". So "Code X", which was the name of a system of "coding" rules, becomes the name of the linguistic matter which arises from transforming plain English according to the transformation rules of "Code X". I do not object to this ambiguous use of "Code X". The strong dependence of "code language" on ordinary language is stressed by the fact that the same word is used for both the system of "translation" rules and for the result of applying those rules. There is, however, still another word that means only the system of coding rules, viz. the word "key".

The relation between a natural language and a code is not the same as that between two languages. Translating from one language into another is quite different from coding and decoding. Coding and decoding can be done by formal substitution, whereas translating presupposes understanding. Of course this is again a gradual difference. The relation between spoken and written language is similar to that between a language and a code. Writing is like coding, and reading like decoding. Yet the rules that govern the relations between a spoken and a written language, are much more complicated than the rules of any cryptography. The rules of writing and reading are not purely formal such as those of coding and decoding. They can hardly be handled by people who do not understand the language in question, even in languages with a nearly phonetic spelling as Italian and Dutch. Coding and decoding machines can be much simpler than machines for writing down spoken language and for reading written language. Furthermore written language is much more independent of its spoken foster-mother than coded language is of its plain substratum. Spoken and written language may widely differ, and spoken language may undergo profound influences from written language. So it is not surprising that one speaks of spoken English and of written English, as though they were really two languages. The relation between plain English and coded English is quite different from this. Calling coded English a language would not match our customary use of the word "language". Yet I would agree to coded English being called a coded language.

005. A few other linguistic or quasi linguistic phenomena need to be considered. If a traffic sign shows the picture of two children going to school, or that of a fence, the car driver understands these pictures as if

they were linguistic announcements of a school and of a level-crossing gate. The skull on a bottle is to be read as "poison", an arrow or a fingerpost may mean "this way", a stroke of lightning is a warning against high tension wires. Pictures of children, fences, skulls, arrows, strokes of lightnings do not belong to customary English, but everybody knows how to translate them. The analogy with a code is striking. Of course there are differences: codes proper are to be secret, and they are fitted for more general purposes than the traffic-signs code. The common feature of these codes and the codes proper is the existence of a key for coding and decoding. Coding and decoding do not presuppose any understanding, but only the mechanical use of coding and decoding rules which can be collected in some list. An astrologer who casts someone's horoscope is considering the constellations as coded messages which can be transferred into plain language by adepts who possess the key. Telling the cards, interpreting dreams, looking for the "sigilla" of herbs as practised by herbalists of earlier times in order to recognize the medicinal value of plants, are methods of decoding something that is considered as a message.

0 06. We rejected the denotations "astronomical language", "medical language", and so on for the idiom used by astronomers, physicians, and so on. We now resume this discussion because there is at least one exceptional case: the case of mathematics. Is there something that, in the frame of a reasonable terminology, may be called mathematical language?

A strong argument for an affirmative answer may be found in what happens if a book be translated from one language into another, let us say from English into French. If this book is a mathematical textbook or treatise, there will be parts that need not be conveyed into, viz. the mathematical expressions and formulae. If we seek for other examples of texts that are excempt from translation, the result will be meagre. First of all: the occurrences of Arabic figures. Yet Arabic figures are, to say least of it, so near to mathematical expressions that they can hardly be considered as a new instance. The proper names Shakespeare, Oxford, and Portugal will reappear in the French translation in their original form, but "Pliny", "London" and "England" will be changed into "Pline", "Londres", and "Angleterre". Clearly it is not on principle that proper names are not translated. Or rather: the French translation of "Shakespeare" is again "Shakespeare", whereas the English and the French name of Pliny differ from each other. (Note that in a Russian translation all proper names would at least be transcribed, whereas mathematical expressions and formulae do not undergo any change.)

Another example: it is highly probable that a translator will leave quotations unchanged, especially if the language of such a quotation

differs from that of the surrounding text. The sentence "Before crossing the Rubicon Caesar said: Alea est iacta" will be translated "Avant de traverser le Rubicon, César dit: Alea est iacta". It is, however, not sure that a quotation of Hamlet's "to be or nor to be" would again read "to be or not to be" in a French translation. (Perhaps adherents of formalist semantics might say that, as compared with proper names, quotations are not a new instance of texts not to be translated, because a quotation is a proper name, namely of the thing quoted. Yet I do not feel the need to use the term "proper name" in a broader sense than usual.)

Even if mathematical expressions and formulae were proper names or quotations (I am sure that they are not), it would be a strange thing that they are never involved in the effective translation procedure. There are not many other examples of so strong an immunity. The best-known case is that of the formulae of chemistry. Physics is less striking. The resistant parts of a textbook on a subject of physics will be mathematical formulae, denotations of measures and units, and some reaction formulae like those of chemistry. The last instance of invariance with respect to translation I will mention, are staves.

I shall restrict myself to invariant parts that are neither proper names nor quotations, and I shall ask the question what is the reason for this invariance, e.g. why the mathematical formulae of an English textbook reappear unchanged in its French version. A priori two different answers are possible: The mathematical formulae are neither English nor French, or they are both. In the case of chemistry and music we may likewise choose between these two solutions.

We prefer the first answer: that the mathematical expressions and formulae belong to a language different from that of the surrounding context, because there are more reasons to adhere to this interpretation. The syntactical structure of "mathematical language" differs enormously from that of all natural languages. The main points of departure are "punctuation" and "treatment of variables".

007. One may fairly well maintain that punctuation is one of the numerous redundant elements in natural languages. There is hardly any interdependence between punctuation in spoken and in written language. Punctuation signs do not match the means of spoken punctuation, pauses and intonation. The question mark is the only punctuation sign that might reflect a kind of intonation, but its actual use is much more syntactic than phonetic. The punctuation signs which might mark pauses are placed according to conventional — in the main syntactic — rules which do not cover the use of pauses in spoken language. Attempts at phonetic punctuation may be found in the t'amim (accents) of the masoretic bible text,

and in the neumata added to the Greek gospel, but both were adjusted to the special use of language in liturgy and they developed into rather musical systems not related to spoken but to sung language.

In speech punctuation still accomplishes syntactic tasks. Perhaps the inability of written language to reproduce spoken punctuation has been one of the reasons why nowadays spoken and written language differ. Complicated syntactic patterns took the place of spoken punctuation. Sumerian reflects a linguistic stage of mankind where connectives were nearly unknown, and nevertheless syntactic patterns of subordination were possible by chain constructions in which a whole sentence could be subject or object of a superordinated sentence - this is still a characteristic of infant language. Such patterns are unthinkable without a well-developed punctuation system. Since then this system has been abandoned in favour of a system of subordinating by means of connectives. The invention of relative clauses has been the most important step in this direction, though originally relative clauses are merely interpolations independent of the sentence on which they seem to depend; indeed the relative pronoun is etymologically interrogative or demonstrative. Relative clauses in the proper sense are still unknown in Accadian and in Biblical Hebrew. The words sha and asher that correspond to our relative pronouns cannot literally be translated by relative pronouns. Sentences such as "a nation whose language you do not understand", reads in a literal translation "a nation asher you do not understand its language". So Accadian "sha" and Hebrew "asher" are still rather punctuations than connectives. It is true that contemporary English (and Scandinavian languages) also know relative clauses without connectives, but this case is quite different from that of the ancient Semitic languages. "The man I have seen ..." would read in an interlinear version of Biblical Hebrew "the man asher I have seen him . . . "; in English the absence of the personal object "him" betrays a dependence of the second clause; in Hebrew it is syntactically independent.

The principle of syntactic structure by means of punctuation has been abandoned in natural languages (at least in those of Western civilisation), but it has been revived in the language of mathematical formulae. As punctuation marks pairs of brackets are used. Pairs of brackets tell us that in $\frac{1}{4} \left[(a+b)^2 - (a-b)^2 \right]$

we start by calculating the sum and difference of a and b, that these numbers are then squared, that the squares are subtracted, and that the difference is to be divided by 4. We should arrive at clumsy linguistic structures, if we tried to translate such expressions into the vernacular. Of course we can read the above expression as "one fourth of square bracket open a plus b bracket closed square..." but this

would be to infringe lexicological and syntactic laws of the vernacular. It is noteworthy that in the vernacular the mathematical process of putting in brackets is described by a reference to the graphical means (the brackets) by which this process is represented in written formulae. Compare this situation with the one we would meet if the vernacular did not possess a word meaning "equal" and speaking about equality were only possible by using graphical descriptions of the mathematical equality sign such as "pair of horizontal lines".

Brackets are not the only means of punctuation used in mathematical expressions. By mere convention algebraic connectives are at the same time punctuation marks. The separative power of the different connectives is fixed by a few rules, some of which are explicitly stated in the textbooks, whereas others are tacitly assumed. The signs of equality and inequality, and of "smaller than" and "larger than" are the strongest separators; the separative power of a plus-sign is greater than that of a times-sign; algebraic sums (such as a+b-c-d+e) are to be read as if all open brackets stand at the utmost lefthand; and so on.

In the vernacular there are no rules about the use of connectives as punctuation signs. In the sentences "we visited Dutch towns and villages" and "we got lessons in Dutch history and mathematics" the apparent syntactic pattern of "Dutch towns and villages" is the same as that of "Dutch history and mathematics" though on semantic grounds we are to bracket "towns and villages" in the first case and "Dutch history" in the second. The sentence "He studied the history of medecine in Bologna" does not tell us whether he was in Bologna when studying history of medecine or whether he studied that part of the history of medecine that was played in Bologna. Again the decisive point is the place of the brackets.

Little attention is paid to the need of formal punctuation in the vernacular, whereas in the language of mathematical expressions it is the backbone of the syntactic structure.

0 08. The other feature that distinguishes mathematical language from natural languages is to be found in the treatment of variables. Philologists usually ignore the notion of variable though they are reasonably acquainted with other notions discovered by logistical analysis of language such as constant, connective, relation, function and so on. Indeed the variable as a syntactic element is not easy to uncover in natural languages, precisely because in natural languages almost all syntactic units may be used as variables though mostly ranging over a very restricted domain; it depends on the context whether a unit is to be considered as a variable or not. Furthermore in almost all instances the variable or at least its binding is hidden and can be recognized only by logical analysis.

It is probable that generic names (appellatives) were originally (ambiguous) proper names, i.e. every appellative was really a proper name of every individual belonging to the genus designated by that name, e.g. ant is a proper name of every individual ant. This conception is reflected in the Bible when "Adam" or "Israel" or "Kanaan" is used not only for one single person, but also for a group of people and for every member of this group. Every man bears the name Adam = Human, or "Son of Adam" = "Son of Human" in virtue of his descent from the first person that bore the name Adam.

The semantic and syntactic change from proper name to generic name has been made in an early stage of linguistic development, as it is still made by children in their early youth (only a few vocatives such as "father", and "mother" will strongly resist the transformation from proper name into appellative).

The semantic transformation of proper names into ambiguous proper names and of ambiguous proper names into appellatives has not been accompanied by systematic syntactical transformations. "Man" will be an ambiguous proper name in the sentence "He killed a man" and a generic name in "man is mortal". Many languages have suffixes, prefixes or connectives that transform ambiguous proper names into generic names, but their use is neither obligatory nor always possible. I may say "man is mortal" though there is a "mankind", which could be used in this case, but I have to say "a mouse is a rodent", because there is no "mousekind". There are suffixes which in one case can transform an ambiguous name into a generic name and which in other cases will alter the sense. "Priesthood" may be the genus "priest", but "brotherhood" will never be the genus "brother".

Ambiguous proper names are an instance of what logical analysis calls variables. In the sentence "I bought a car" we recognize "a car" as a variable that has been bound by the fact that there is some car that I have bought, thus by existential binding. In the sentence "A car is a carriage" one can interpret "a car" as a generic name or as a variable that has been generalized over the genus "car". There are other kinds of binding, usually not considered in logical analysis, but as important as the preceding ones, e.g. the interrogative binding such as that met in "Which car did you buy?"

The linguistic equivalent of a variable may be still more complicated. In "I bought three apples" and in "three apples (of a certain kind) go to a pound" we may consider "three apples" as a variable ranging over all triples of apples; in the first case the binding is existential, in the second case it is generalizing. The use of "little money" in "I have little money" and "with little money you can succeed at Smith's" is quite analogous.

Pronouns are still other instances of variables. "I" is a variable bound

to the person who utters it by the fact that he does so. "You" is a variable bound to the person that the speaker is addressing. "Now" is a variable bound to the moment at which it is pronounced. "Here" is a variable bound to the place where it is pronounced. One might speak of demonstrative binding in these instances, because the demonstrative pronouns "this" and "that" are variables of this kind with an unusually large domain. Formally it is, however, possible to reduce demonstrative bindings to existential bindings.

These examples show that natural languages are equipped with a large stock of variables which all have different domains. Natural languages are unsystematic in dealing with variables; and they are far from a clear-cut solution of the problem of indicating the means by which a variable is bound. The makers of language acted instinctively, uninfluenced by rational analysis. The use of ambiguous proper names as generic names was a discovery. Rational analysis has turned it into a problem. The eleatics and sophists were troubled by the "one that can be many", but the first who consciously justified generic names as a new kind, was Socrates - in Aristotle's appraisal of Socrates' work stress is laid on this very fact. In Socrates' tracks Plato and Aristotle arrived at a frankly ontological solution of the problem. A generic name is said to be the proper name of a thing that is no less real than individuals are. These things are called ideas by Plato and genera by Aristotle. Historically this ontological interpretation seems to be provoked by the rather superficial and incidental methods which natural languages have applied in order to settle the relation between ambiguous proper names and generic names.

It would be an enormous task to analyse how the various natural languages operate variables. We have already noticed that every variable has its own domain, "car" cannot be used for other things than cars, "he" ranges over the set of males only, "yesterday" is a temporal variable, "there" varies over places and "thither" over directions. With the help of numerals one can form variables such as "five fingers" and "one hundred people". "One gallon" can hardly be interpreted as a variable, but "one gallon of milk" can. "One shilling" behaves differently in "this costs one shilling" and in "I gave him one shilling". In the first case it is not a variable, but something else, in the second case it will be a variable (at least if the sentence really means that he got a shilling, and not that the thing he got was worth one shilling). It is unnecessary to point once more to the indifference natural languages show to the character of binding, whether it be existential or generalizing.

We shall now compare this use of variables in natural languages with that in the language of mathematical expressions. It is a striking principle in this language that all variables have the same domain and are freely interchangeable. Incidentally the domain may be restricted and the interchangeability may be partially blocked. This can be done in a formal way, e.g. if in the first chapter of a book on linear algebra the author announces that Latin letters will be used for vectors and Greek letters for scalars. The reader of a paper on elementary geometry is assumed to know that variables written as Latin capitals range over the set of points of three-space and variables written as Greek capitals range over the set of angles. The reader of a paper on complex functions may conclude from the context which letters are related to complex numbers and which to functions. If a theorem is pronounced, it is understood that all restrictions imposed on the variables belong to the premises and may easily be added, if they do not occur in the text as it stands. This does not affect the principle of uniform domain and free interchangeability.

It is true that this principle is not always strictly observed. The best-known infringement is related to a notation of function that in thirty years may be considered old-fashioned though it is now very common. If a function f is called f(x), this may have been done in order to state a special relation of this function f with the particular variable x, so that f(y) will be another function than f(x) though f and f may range over the same domain. In the case of functions of more variables this notation will be a means of distinguishing between "the functions f(x, y) and f(y, x)". It is well-known that this notation leads to serious inconsistencies, especially when partial derivatives occur. It is a remarkable feature in the development of mathematical language that mathematics are becoming more and more repugnant to such infringements.

By the device of free interchangeability of variables mathematical language avoids a difficulty that is vividly felt in natural languages. In natural languages there is not more than one variable that ranges over the set of stones (books, trees, numbers and so on) viz. the variable "stone" ("book", "tree", "number", and so on). So if I wish to speak about two or more stones (books, trees, numbers, and so on) in the same context, I am committed to clumsy circumlocutions such as "this stone, and that stone", "one stone and another", "the former and the latter", "the first, the second, the third" (and analogous constructions). Incidentally such variables may be distinguished by letters, "stone A, stone B, stone C", and this is just what mathematics has done systematically, because the simultaneous occurrence of many variables meaning variable points or numbers has been a not unusual phenomenon in mathematical texts from the oldest times. Finally the need for a large stock of variables has led to the principle of free interchangeability.

It will be clear from this exposition that natural languages and mathematical language differ widely as to the management of variables.

0 09. In 7 and 8 we have seen that mathematical expressions have a syntax of their own which, owing to a rapid evolution, may successfully deny its descent from the syntaxes of natural languages. So I feel justified when I speak of a "mathematical language" as opposed to natural languages and to medical or juridical idioms which lack almost all characteristics of an independent language. Even in the case of chemical formulas one may wonder whether it would be reasonable to state that they constitute a more or less independent language. The syntax of these formulas, though very poor, is patterned after that of natural languages; for every chemical element there is only one variable that may vary through atoms of that kind (called H for the hydrogen atom), and the only respect in which the syntax of chemical formulas is superior to syntaxes of natural languages is a good system of punctuation.

It is true that mathematical language as written in textbooks and papers still parasitizes on natural languages. The text surrounding the mathematical formulas is usually written in an idiom that bears the characteristics of the vernacular, to which it belongs in the ordinary sense. This is why I spoke of a language of mathematical expressions. The mathematical expressions about which the author reasons are written in this language, whereas the reasonings are usually worded in the vernacular.

So far I have not yet mentioned the logistic language created by G. Peano and perfected by B. Russell and A. N. Whitehead. There have been earlier attempts, but Peano was the first to design a linguistic pattern more adequate to mathematical reasoning than common language. Logistic language used and developed consciously the peculiar features of the language of mathematical expressions. Of course words that mean logical connectives had to be introduced, but their number could be cut down, as syntactic structure is mainly got by means of punctuation: variables are treated as in the language of mathematical expressions, systematically and according to the principle of free interchangeability. Further, words are created to account for the kind of binding that has been applied to a variable, whether existential or generalizing.

Russell and Whitehead's purpose was different from Peano's. Peano's idea had been to create a workable language that should supersede the vernacular in mathematical texts. Yet in the work of Russell and Whitehead the stress shifted to the foundations of mathematics. Logistical language proved useful in foundational research, and improvements tended to increase that utility. Instead of the colloquial language of mathematicians logistics became the subject-matter of foundational exploration. Mathematicians not specialized in this domain did not mind whether logistic language could serve any other purpose than the very special ones of foundational research. Many mathematicians even distrusted

logistics, because practicians of this discipline used to link it to some school of philosophy and because certain details created an impression of hairsplitting and clumsiness. For years I too shared this false impression. It is true that the actual systems are a little too rigid. Though logistics has influenced the idiom of mathematical texts we still have little experience at putting mathematical expressions into a frankly logistical frame. Mathematics, particularly in its youngest branches, has become an intricate machinery. I think logistics may be a great help to better understanding.

0 10. As logistical language has not been much applied to mathematics, it is not to be wondered at that no attempts have been made (at least as far as I know) to apply logistics on a reasonable scale in some other domain, let us say physics or daily life. Before logistics came into existence, people were more optimistic. Leibniz fostered the idea of a "characteristica universalis", and under his influence as well as independently of him, 17th and 18th-century scholars developed systems of this kind. At that time the syntactic problems were not yet posed, still less solved. The attempts did not lead to languages but rather to rational systems of concepts, forerunners of the decimal vocabulary of our days.

At present affairs would appear to be much more promising. Thanks to logistics the syntactic problems are solved or, in any case, elucidated. One can imagine numerous ways of tackling the subject. After several unsuccessful attempts I finally became convinced that it is just the difficulty of choice which causes the trouble, and that the only thing which matters is to find a starting point. Seeking in history how analogous situations were met, I came to the conclusion that one should start with a concrete, sharply-defined and rather narrow problem.

011. In spite of Peano's original idea, logistical language has never been used as a means of communication. If there was something to be communicated, even logisticians kept to the vernacular. Logistical language became a subject-matter. For the sake of a better understanding of its intrinsic properties one tried to isolate it as strictly as possible. The bonds with reality were cut. It was held that language should be treated and handled as if its expressions were meaningless. Thanks to a reinterpretation, "meaning" became an intrinsic linguistic relation, not an extrinsic one that could link language to reality.

A language that is to operate without regard to the meaning of its expressions, e.g. by people who do not understand it, must fulfil very high requirements. It must be fully formalized. Formal rules of constitution and transformation will determine the structure and handling of the linguistic expressions.

Formalization of language is a highly important problem. Coding as treated in 0 04 is an example of formal translation. Indeed, in such cases it is imperative that a code works without any regard to the meaning of the text to be coded or decoded.

This is a rather unproblematic example, because inside the coded text there is no need for operations, any more than for rules of operation. The coded language has a static character. The best-known example of an "individual" really operating within a fully formulaized language is the computing machine. Independently of any closer examination of the concept of meaning, we may assert that the machine does not understand the meaning of the linguistic expressions on which it is operating, in the sense in which we understand them. The machine works according to formal rules, whereas the human being who runs the machine knows the meaning of the input and output text. The subject matter on which the machine operates is language only from the point of view of the human being, just because it is he who attaches some meaning to it.

0 12. In view of the logistic tendency towards formalization it is not surprising that logisticians have paid less attention to language as a means of communication. In 0 10 I have explained why in the construction of a characteristica universalis we have to start with a concrete, sharply-defined, and rather narrow problem. After these explanations it will be clear why I think that this problem must be a problem of communication, and more precisely communication ab ovo. My purpose is to design a language that can be understood by a person not acquainted with any of our natural languages or even their syntactic structures. The messages communicated by means of this language will contain not only mathematics, but in principle the whole bulk of our knowledge. I shall assume that the receiver of these messages has understood their language if he is able to operate on it. He will be able to do so only if he has grasped the meaning of the expressions of this language, because it will be a moderately formalized language that cannot be handled on the ground of formal rules only.

The problem is not yet sufficiently outlined. Communicating with individuals who are not acquainted with our natural languages is a problem that is daily solved in the intercourse with babies and infants. We succeed in teaching them our language though we have started with a tabula rasa of lexicologic and syntactic knowledge. This linguistic education takes place in surroundings where things and situations are shown, and named at the same time, by people who are acquainted with the language. Showing will mostly be unintentional; the act of naming will mostly include that of showing, even if the thing or the situation is not named for the purpose of being shown.

In order to narrow the problem, I propose to exclude or at least to restrict excessively the opportunities of showing. I shall use showing as little as possible as a means of explaining the meaning of linguistic expressions.

On the other hand I shall suppose that the person who is to receive my messages is human or at least humanlike as to his mental state and experiences. I should not know how to communicate with an individual who does not fulfil these requirements. Yet I shall not suppose that the receivers of my messages must be humans or humanlike in the sense of anatomy and physiology.

013. It was in this way that I arrived at the problem of designing a language for cosmic intercourse. Lincos, the name of this language, is an abbreviation of "lingua cosmica". As a linguistic vehicle I propose to use radio signals of various duration and wave length. These two dimensions will suffice. As compared with our acoustic vehicle, intensity as a third dimension is lacking. I rejected it because I think it less fit for radiographical use. But this is not a serious loss, because the wealth of information is much greater in the ether than in the acoustic medium, even if we dispose of the dimensions of duration and wave-length only.

I have not made a study of communication techniques, so I shall not dwell on technical details. When I designed Lincos, I had in mind using unmodulated waves. Perhaps modulated waves would be better. It would require only slight changes to adapt the program of messages to the use of modulated waves.

I have not examined which distances can actually be covered by radio messages. It is very probable that we can reach other planets of our solar system by means of radio signals, but I doubt whether with the powers now available we can communicate with inhabitants of other solar systems. Nevertheless I have taken into account the possibility that the receiver will belong to one of the neighbouring solar systems.

Of course I do not know whether there is any humanlike being on other celestial bodies, and even if there were millions of planets in the universe inhabited by humanlike beings, it is possible that our nearest neighbour lives at a distance of a million light-years and, as a consequence, beyond our reach.

On the other hand it is not unthinkable that inhabitants of other planets have anticipated this project. A language for cosmic intercourse might already exist. Messages in that language might unceasingly travel through the universe, maybe on wavelengths that are intercepted by the atmosphere of the earth and the ionosphere, but which could be received on a station outside. On such an outpost we could try to switch into the

cosmic conversation. Long ago I thought cosmic radiation to be a linguistic phenomenon, but at the present state of our knowledge this seems to be very improbable. It is not easy to state a priori how to distinguish messages from purely physical phenomena. But should the case really arise, we shall know how to answer the question: we should try to understand the message.

This, I suppose, intelligent beings in the universe will do if they receive our messages. They will try to decipher them and to translate them into their own language. This task might be easier than that of terrestrial decipherers who have to discover the key of a code. Indeed our objective is just the opposite of that of the sender of coded information. We want to communicate with everybody who might receive our messages, whereas the sender of a coded message wishes to keep secret the information contained in his message. But in spite of our efforts even intelligent receivers might interpret our messages as physical phenomena or as music of the spheres.

- 0 14. It has been one of our presumptions that we cannot start cosmic intercourse by showing concrete things or images of concrete things. One might observe that television images are sent by radio and that we could try to do likewise. Television is founded on the principle that an area is decomposed into a set of spots and that these spots are arranged in a simply ordered (temporal) sequence according to a fixed rule. As long as we cannot tell the receiver the principle and the code of decomposition and arrangements, we cannot expect that he will understand television messages and transform them into images. Synthesizing television messages into images may be taught by means of analytic geometry. In our Lincos program analytic geometry is a rather advanced stage. When this is reached, television, even three-dimensional, will be possible. But by then our linguistic proficiency will be great enough to dispense with television. I tend to believe that verbal description is superior to showing by television. I have not however paid sufficient attention to this problem as I am mainly interested in linguistic problems.
- 015. Throughout our exposition we will talk of Lincos as if it were speech (not writing). We shall say that Lincos phonemes (not letters) are radio-signals of varying wave-length and duration, and that Lincos words are constituted of phonemes (not of letters). If we aim at this constitution, we shall speak of phonetica (not of spelling).

This is no more than a convention, but it will prove very useful. We shall need to consider two versions of Lincos, Lincos proper, and a coded version. It will be convenient to use the term spoken Lincos for Lincos

proper, and the term written Lincos for coded Lincos. Moreover this choice of names is justified by the fact that the broadcast language will be "spoken Lincos" whereas in printed expositions such as the present book "written Lincos" will be used.

Written Lincos will not bear a one-to-one graphic-phonetic relation to spoken Lincos. On the contrary the relation will be global; every spoken Lincos word is reflected in written Lincos by a rather arbitrary graphical sign or group of graphical signs. This is a realistic procedure. A phonetic spelling of Lincos would be rather clumsy and illegible. Moreover for the time being there is no point in knowing how the different Lincos words will be pronounced (i.e. translated into groups of radio signals) when they are broadcast. The reader of the present book, who meets with words of written Lincos such as 10, +, \wedge , Num, PAN, Q, and so on will understand that these graphical complexes are codings of spoken Lincos words. It is quite unimportant how these spoken Lincos words are constituted, and I have never concerned myself with this question. In any case the spoken Lincos word corresponding to Num need not be composed of phonetic groups which correspond to N, u, and m. In this book we shall stick to the convention that e.g. 'Num' is not the proper Lincos word but the written (or printed) image of it. But we shall not always stress this feature, and sometimes, in order to avoid clumsy language we shall even behave as if the Lincos word were 'Num'.

0 16. The notations in written Lincos are borrowed from mathematics (10, +), from logistics (\land) , from other sciences (\lozenge) , and from Latin (e.g. Num is an abbreviation of Latin numerus). As a rule such abbreviations will be of three letters; we also use contractions of groups written of Lincos words, such as PAN (from Pau Ant Nnc). Finally there are Lincos words written in entirely arbitrary fashion, such as Ha, Hb, Hc,... which mean names of humans.

As pointed out before the spoken Lincos correlates of the written Lincos words will be rather arbitrary. A certain restriction, however, should be made. Lincos will be more easily understood if it has a systematic phonetics. It would be pleasant if the various syntactic and semantic categories could be distinguished by certain striking phonetic characteristics. A receiver who has to interpret the Lincos messages and who has made some progress in understanding them, will be grateful if he can see at first glance whether a word means a variable, a connective, a set, a human, a chemical element, and so on. Probably this will even help deciphering. Disconnected fragments of such a systematization exist in all natural languages. A noteworthy attempt to amalgamate these fragments has been made in Esperanto. There the syntactic and semantic categories emerging from Western

languages have been accepted without much criticism and have been characterized by the traditional means of stems and inflections.

Establishing the principles of an adequate Lincos phonetics would be premature. No more than a rough systematization would be available at the present time. A few syntactic and semantic categories can be clearly distinguished at the moment, but their number is rather small. We have still to wait for a more extensive vocabulary of Lincos. Then we can resume this question.

0 17. In a more essential way the choice of Lincos phonetic means is restricted on account of the use of ideophonetic words. We have agreed to abstain as much as possible from showing, but we cannot entirely abstain from it. Our first messages will show numerals, as an introduction to mathematics. Such an ostensive numeral meaning the natural number n consists of n peeps with regular intervals; from the context the receiver will conclude that it aims at showing just the natural number n. Ostensive numerals will soon be superseded by algorithmic ones. An ostensive numeral is ideophonetic (as a hieroglyph both showing and meaning an eve is ideogrammatic). Other ideophonetic words are the so called timesignals introduced in the second chapter. These words have no simple correlates in other languages. One might even doubt whether they are really words, or rather signs, that mean nothing but themselves. More things will be shown in the Lincos broadcasting program, but then the means of showing are still less linguistic. With a common name such occurrences might be called noises. Some of them mean nothing but a noise, others mean unintelligible speaking, and still others mean events which are usually accompanied by a noise, such as in a radio play the noise of smashing a door means smashing a door.

The phonetic structure of an ideophonetic word is strongly determined by its meaning. It cannot be arbitrarily changed. This is even true of the less linguistic pieces of our program. Of course a noise meaning a noise may be replaced by another noise, but in any case the substitute ought to be again a noise.

There is still another kind of truly linguistic phenomena which might be reckoned among ideophonetics: the means of punctuation. We have already had occasion to point out the importance of punctuation in mathematical and logistical language. Punctuation will be the principal means of showing Lincos syntactic structure. Lincos punctuation will be explained in 1 05 0. As punctuation marks we will use pauses, because pauses are self-explanatory to a high degree. It is a debatable point whether purposive pauses are to be considered as words. If we agree to this terminology, we may count Lincos punctuation marks among the

ideophonetic words. From a syntactic point of view they are a special kind of connectives.

For punctuation in written Lincos we have adopted the Russell-Whitehead system of square dots, though in a more economical version (see 1050).

0 18. The syntax and the lexicon of connectives of Lincos have been mainly borrowed from logistical language, but at the same time we have tried to mitigate the rigidity and clumsiness of logistical language. Looking back at the result, I must confess that I am not satisfied, mainly because I often lacked the courage to go far enough. (Formalists will surely claim that I have gone too far.)

Lincos should be a means of communication. Lincos syntax is not an aim in itself, but it should serve the general aim of Lincos. Lincos will be far from fully formalized. It will be used by people who are supposed to understand what they are saying, and to try to understand what they perceive, not by machines which must operate according to formal rules. So we can dispense with full formalization, even where this is possible. To a high extent comprehension may be favoured by formalization, but it may also be hampered by needless exaggerations.

The liberties we have taken are very usual in common language, the common idiom of mathematics included. Words and clauses which can readily be supplied are often omitted, especially when pieces from a question would have to be repeated in the answer to this question. Sometimes we account for the omitted words or clauses by a word meaning "and so on", but this is not compulsory. Variables bound by a definite or indefinite article are repeated without the clause by means of which they have been bound. In the common mathematical idiom an existentially bound variable is allowed to reappear as an instance out of the set to which it has been bound. We have adopted this habit, sometimes justifying it by a more formal device (see 3 05 2), but sometimes disregarding even this formality. Looking back I regret that I did not take more care for uniformity in this matter. It would have been easy to introduce words meaning "this", "that", "the afore mentioned", and so on, and it would have been a real gain. On the other hand it is not so important a question as to justify a total revision of the project.

Definite and indefinite articles, and other words which conceal existentially bound variables, may only be used if it is clear where the binding takes place. In many instances of our Lincos program this will be decided not on syntactic but on semantic grounds.

Lincos speech has been assimilated to common language by introducing words meaning "thus" (3 04 5), "suppose that" (3 05 2, 3 05 3), "because" (3 12 0), and others, and by using them in a not too formal way. "And"

and "or" as used in "John and Mary have come" and in "John or Mary will come" are represented by Lincos words, of course different from the usual syntactic connectives meaning "and" and "or" (1 29 2). A rather verbal translation of "three people" such as occurring in "three people have come" is possible in Lincos (3 28 1). Modalities have been dealt with in a quite unformal way, disregarding the highly developed formal theory of modality (3 35 1–3 35 7). Nor have I respected the formalism of quotations and quasi-quotations; in Lincos one may speak about things and events in a recklessly autonymous fashion, i.e. by using copies of these things or events as names. I have decided on this conduct only after a thorough analysis which led to the result that the usual precautions are unsatisfactory and superfluous (2 04 0–2 06 0). The reader will have remarked already that we have not adopted the rules of formalist semantics in the use of quotation marks, or any other conventions of this kind.

I have lowered the degree of formalization of Lincos purposely and deliberately though without any intention of abusing this argument. It is not meant as a means of evading criticisms. In all cases where more formalization would have favoured comprehension and comprehensibility, I shall plead guilty. I feel convinced that this will happen rather often, and that even very serious mistakes may have occurred.

019. I have dwelt upon divergences from the syntax of logistical language. This might mislead the reader. In fact these divergences are far from essential, surely less essential than the concurrences and the additions which have proved indispensable.

One of this additions is concerned with the syntactic structure of oblique or reported speech. As far as I know logisticians have always treated oblique speech in another context, viz. that of the problem of intentionality. In my conviction this is a pseudo-problem of formalist semantics which is unable to deal with language as a means of communication adequately.

The syntactic problem of oblique speech arises if a certain person A says that a certain person B said that p, and if p contains variables which do not belong to B's but to A's discourse (e.g. if B said "I am ill" and A changes the variable "I" into "he"). Syntactical means have to be devised in order to ascertain whether such a variable is genuine or belongs to the discourse of another person. The devices used in common languages are too unformal; they cannot serve in Lincos. Our solution is quite different. In Lincos hidden variables and words that look like variables will always be genuine, whereas overt variables belong to the discourse in which they are bound. We shall see that such variables cannot be avoided. The problem has been fully treated in 2 04 0-2 06 0, 3 01 5 and 3 07 0.

Hitherto in logistics little attention has been paid to sentence patterns other than the propositional one. The interrogative mode must however be accepted if one considers mathematics not as a stock of true propositions, but as an art of discovering. Moreover language as a means of communication cannot dispense with the interrogative mode. Interrogative sentences will be formed by a new kind of binding variables, the interrogative binding. We introduce a connective '?', which prefixed to a variable x in a context '? x...' produces a question that may be translated "wanted an x with the property..." (see 1 12 1 and 3 01 1, 3 02 1, 3 02 2 and so on). Of course many more problems of syntax have had to be solved. In fact, whenever a new word is introduced, its syntactic status has to be settled. This requires a careful analysis in every single case. Fortunately the variety of syntactic status is rather small, thanks to the existence of a few syntactic paradigms, such as the words meaning "to happen", "to say", "to know", "second". For instance the unit of time is a paradigm

I have tried to outline the syntactic status of every word sharply, but I am not sure whether I have always strictly observed these syntactic rules. In any case it is still too early to write a syntax of Lincos.

for all units of measure; "second" is handled as if it were a (linear) function

from the set of real numbers to the set of durations.

0 20. Lexicology is much more embarrassing than syntax. In a field which can be organized by simple rules, lexicology reduces as it were to syntax. A vocabulary of natural numbers can easily be designed on the basis of some positional system. Number vocabularies of soldiers are known in modern armies, and it would be easy to draw up similar vocabularies comprising all humans. Geographical coordinates can serve to set up a vocabulary of geographical positions; a pair of integer coordinates can be used to name chemical elements. But these are exceptional cases. We are still far from constructing the rational system of notions of which Plato and Leibniz might have dreamed, perhaps we are even farther than they were. We have not the faintest notion of the character of the enormous preparatory work that we ought to do if we want to realize those dreams. The system of stems and inflexion of our natural languages is more historical than rational, and even in languages with a relatively simple history such as German, this system is a logical chaos. In Esperanto inflexion has been intentionally organized, but not on logical grounds. (This is not at all a reproach. Nobody can tell us what conditions good means of communication are to fulfil. Under exceptional circumstances a rational structure may be advantageous. In daily life common languages have proved to be extremely useful in spite of their chaotic structure.)

As long as we have not the faintest idea of a rational system of notions,

lexicological engineering will be a mere succession of haphazard attempts. When introducing a few hundred Lincos words successively, I have not been guided by lexicological rules. I have made a rather sporadic use of inflexion as a means of formation of words. Two of the three affixes I have introduced are unproblematic; they transform a cardinal into a ordinal number and into a multiplicative number respectively. With the third I feel less satisfied. Its function is to transform a verb into a nomen actionis. If such a word really means the action and not also the result of the action as is often the case in common languages, it can be probably dispensed with.

0 21. Lincos has to be taught to the receiver. Therefore our program to be broadcast cannot have the character of a newsreel. On the contrary, in the beginning we shall communicate facts which may be supposed to be known to the receiver. In this way we can try to facilitate the work of deciphering.

As our means of showing are heavily restricted, we cannot start with a concrete subject matter. Mathematics is the most abstract subject we know and at the same time a subject that may be supposed to be universally known to humanlike intelligent beings. So we have decided to start our program with mathematics. Natural numbers are introduced by ostensive numerals in contexts of elementary arithmetics. The first texts of this kind will be true formulas of the type a+b=c, a=b, a>b, a< b, where a, b, c are to be replaced by natural numbers.

Of course we cannot communicate all formulas of this type because of their infinite number. We shall communicate a number of such formulas that is large enough to justify the supposition that an intelligent receiver can conclude from these texts what the words (complexes of signs) +, =, >, < mean. By a finite number of instances the relation between natural numbers described by a+b=c, or a=b, or a>b, or a< b is not uniquely fixed from a formalist point of view. A machine that does not know any more about the relation a+b=c, but its truth value for natural a, b, c with $1 \le a$, b, $c \le 1000$ cannot solve similar problems with numbers a, b or c > 1000. Yet the receiver will be supposed not to be a machine but humanlike. Children learn arithmetics by a finite number of instances, and a set of rather unformal rules. They succeed in generalizing the instances in a unique way. Arithmetics yields unique results though the definitions leave room for an infinite diversity. If I ask anybody to continue the sequence 1, 3, 5, 7, 9, 11, I may be sure that he will do it with 13, 15, 17 and so on. So I am sure that a humanlike being will conclude from a sufficiently large number of texts like 3>2, 7>1, 9>3, that '>' has its usual meaning.

It is perhaps superfluous to insist on this point. A hundred years ago such assertions would have been self-evident. Since then we have got into the habit of giving formal general definitions and proofs even for arithmetical and algebraical motions and formulas which could be defined and proved in a less formal way. It is not long ago that in arithmetics and algebra a common method of defining and proving was the method I have called elsewhere the quasi-general definition and the quasi-general proof. Euclid proves the existence of infinitely many prime numbers by showing that three prime numbers being given, one can construct a fourth different from them. The proof is formulated in such a way that everybody understands immediately that it intends to show how n prime numbers being given, a (n+1)-th prime number can be constructed. Euclid could not sufficiently formalize such a proof or even such an assertion because he lacked a functional symbol which meant numbering a general finite set, such as our system of numbering by indices, which we owe to a gradual development that started after Vieta and lasted nearly three hundred years.

Euclid's proof for the assertion A_n is a formal proof only for n=3. For A_n it is "quasi-general", i.e. every reader can be expected to be able to recognize the general pattern of the proof, and to generalize it for all natural numbers.

Up to the end of the 19th century it was commonly held that even geometrical proofs are quasi-general: a geometrical theorem is proved by means of one figure in such a way that the proof can be transferred to any other figure of the same kind.

Even now the quasi-general proof is still indispensable as a didactic tool, and the same is true of the quasi-general definition. In school mathematics, in sciences, and in daily life most definitions are quasi-general. One defines by means of instances, though in such a way that the other can generalize the definition as desired by the person who has given the definition. The method of quasi-general definitions (and even of quasi-general proofs) has been a fundamental principle in the construction and presentation of the Lincos vocabulary, in the chapter "Mathematics" as well as in the later chapters. The reader will meet a good many examples of this method. However, we may remark that our definitions will never have the logical form of a definition. Things will not be defined, but described. We shall not say "this is called a clock", but "this is a clock".

I will mention another didactic principle which has often been applied: If a set (a predicate, a relation) is to be introduced and named, we shall start not with a formal definition, but with a quasi-general definition, displaying a sufficiently large stock both of elements belonging to the set and of elements not belonging to the set (possessing the predicate, fulfilling the relation). So we shall not introduce the predicate "integer" until other

predicates can be named. We shall not introduce a word meaning action, until a large number of actions is known, and so on.

From the aforesaid it is clear that mathematical notions will not be introduced axiomatically. Of course when the Lincos vocabulary has developed far enough, one can add an axiomatic, and one can join this axiomatic to the preceding nonformalist mathematics by a semantic: We have renounced this, because it would not touch our main problem. In the chapter "Mathematics" we shall restrict ourselves to the fundamental notions. We shall even drop notions of calculus which are needed in other chapters, because at present it cannot be our aim to write a textbook of calculus in Lincos.

0 22. This is closely connected to another feature of our project.

The definite program to be broadcast will consist of a large number of pieces called program texts. The collection of texts to be found in the present book is to be considered as an abstract from the definite program. It would be no use publishing soemthing like the definite program in full. We have restricted ourselves to one example for every kind of program text which will occur. A kind such as a+b=c (a, b, c, natural numbers) is represented by the single example 4+2=6, though actually the program will contain twenty or more of this kind. If variables are introduced as a new kind of word, every text containing variables will be repeated with other variables, so that a sufficiently large stock of words meaning variables and interpreted as variables by the receiver becomes available. Program texts in which certain humans display a certain behaviour are to be followed by similar texts in which a similar behaviour is displayed by other humans, so that the receiver can understand that this is a general behaviour; and so on.

Written program texts are enclosed in pairs of signs * which do not belong to the program text itself.

- 0 23. For the convenience of the reader the program texts have been collected in paragraphs, and the paragraphs in chapters. The first chapters bear the titles
 - I. Mathematics, II. Time, III. Behaviour, IV. Space, Motion, Mass.

Further chapters entitled "Matter", "Earth", "Life", and a second chapter "Behaviour" have been planned for the second volume of this project. The reader should not attach too much importance to the division into chapters and to the titles of the chapters. Systematics is a passion not only of mathematicians. According to an ancient pattern, teaching a language is subdivided into chapters, running from the article and the

noun to conjunctions and interjections. This system is to be condemned not because it is illogical, but because it is a system. Systematics is an antididactic principle. To the same degree that it is a comfort to the teacher, it can be a discomfort to the pupil. There will be many people who are keen on systematics and who will judge the present work to be rather chaotic. My own feeling is rather that I have too often preferred systematics to didactics.

0 24. A short abstract will give an idea of the subjects dealt with in the present volume, and of the method which has been followed. All details will be omitted.

Chapter I. Mathematics.

Ostensive numerals, quickly superseded by algorithmic ones, written in the dyadic positional system. '+', '>', '=', '<' introduced by numerical examples. Introduction of variables. Means of punctuation. The logical connectives '-->' '\(---\)', '\(\)', '\(\)' are introduced as connectives between arithmetical propositions. Problems are set, using the interrogation mark. Zero and negative integers, further arithmetical operations and relations. Dyadic fractions, and especially periodic dyadic fractions, which give the opportunity to introduce the word "Etc" meaning "and so on". The first predicates (sets) are 'Num' and 'Int' (natural numbers, and integers). The connectives '€', '\varepsilon'. The quantors '\'and '\'arise as infinite conjunctors and disjunctors. More predicates and relations are defined, especially 'Pri' (prime numbers), 'Rat' (rationals), 'Rea' (reals), 'Com' (complex numbers), in order to prepare the more general notion 'Agg' (sets). Real numbers appear originally as infinite fractions; afterwards a more systematic theory is given. Set theory is dealt with in a rather unsophisticated way. Functions are introduced as we usually do in elementary courses; afterwards a formal definition is given. Algebraic notions such as group and field. Propositional calculus; the words 'Ver', 'Fal', '-,' (true, false, non).

Chapter II. Time.

Duration measured by seconds is introduced by means of "time signals". A clock is installed. Mentioning events by means of time data.

Chapter III. Behaviour.

Actors appear. The fundamental action, 'Inq', is "saying". Actions are valuated by the judgements 'Ben' (good) and 'Mal' (bad). Conversations start with mathematics. One actor sets a problem, another solves it, and the first states whether the solution was good or bad. (If the problem is to find the x with 6x=4, the solution 4/6 will be bad.) Typical behaviour

words, such as "Tan' and 'Sed', meaning "but". A great many interrogative pronouns and adverbs can be formed by means of the interrogation mark. Ordinals, multiplicative numerals, and words meaning "counting", "computing", "proving", "seaking", "finding", "analyzing", "responding", "changing", are introduced by showing the related actions. "Supposing that", "in the way of", "thus", "instead of", "omitting", "adding", "why", "because" appear. In common languages "why" and "because" are polysemantical. Provisionally we have adopted this use in Lincos. Later on it will be possible to distinguish the various meanings of "why" and "because" by means of indices.

Much care has to be bestowed on introducing words meaning "knowing", "perceiving", "understanding", "thinking".

An entirely new element appears with a word meaning "nearly". "Error", "much", "little", "too" (such as in "too large") are akin to it. "Now" is defined as "nearly the moment of pronouncing now". "Short time back", "long ago", and so on are akin to "now".

"Necessary" and "possible" are the first modalities that appear. Their meaning is borrowed from common language, not from modal logic.

The next topics are "age" (of humans), "existence" (actual, not logical), "birth", "death", "man", "animal", demographic statistics of actual mankind, average development of humans.

Modality is continued with words meaning "wishing", "allowing", "promising", "being obliged", "being allowed", "it is courteous", "it is convenient", "difficult".

Mathematical games are used as a means of showing human behaviour. Words meaning "giving", "penny", "fortune", law"", "game", "playing", "winning", "losing", "probability", "utility", "fair", "aiding", "ally" appear in this connection.

At the end the paradox of the liar is dealt with.

Chapter IV. Space, motion, mass.

Mechanics is not drawn up as an axiomatic system. The fundamental mechanical notions are introduced by acts of behaviour, approximatively in the way schoolboys usually learn them. Afterwards they are embedded into a mathematical system.

Difference of place is characterized by the delay of messages; distance is said to be proportional with the time of delay; the length unit is called centimeter, but its value is still unknown. Space is the union of all places. Special properties of Euclidian space are mentioned, and finally a description of Euclidian space as a three-dimensional vector space is given. Angle and volume are defined. Measures of humans and animals are mentioned. Movements are studied. Causing movements, carrying, throwing, mean

velocity, oscillation, vibration, communicating through undulating mediums, acoustical velocity, light velocity — this sequence of words leads to the numerical definition of the length-unit which is sharpened by communicating Rydberg's constant.

Mass is introduced as a property of things that determines the difficulty of moving the things. More precisely mass ratio is defined by the law of elastic collision. The mass unit remains unknown. Density, especially that of human bodies. Things, humans, and animals have bodies. The existence of the body of a human begins earlier and finishes later than the human itself. Origin and place of a human body before the existence of the human; notions "mother", "father", "female", "male".

Velocity, acceleration, free fall, Newton's attraction law; this fixes the mass unit. Planetary motions, revolution, rotation. Masses, sizes, and some movement elements of members of our solar system, and of a few stars in our neighbourhood.

Solid bodies (their shape is difficult to change), liquid bodies (their shape is easy to change, and their volume is difficult to change), gaseous bodies (their volume is easy to change). Measuring: length by means of measuring-staffs, time by means of a clock, angles by means of theodolites, mass by means of balances — the instruments are described in a schematic way. Measuring length by means of light signals leads to relativistic mechanics. The fundamental relativistic notions are introduced. This has been indispensable, because of the previous introduction of the velocity of light as a universal constant.

0 25. In the remainder of this introduction we shall continue the analysis of linguistic and semantic notions interrupted by the exposition of the main features of Lincos.

When designing Lincos the main problem I had to solve was how to get rid of Russell's imposing theory of "description". This analysis is particularly illustrative of Russell's work and, more generally, of formalist semantics. Its background is not linguistic in the usual sense of the word. No arguments are taken from language as a means of communication. Description as it originates from this analysis, is not a means of describing, but a mere technical device for reformulating and abbreviating certain logistic expressions. Description in the usual sense has been eliminated by Russell's analysis.

This is in accordance with the goal of formalist semantics. Its main problem is to find formal rules by which the meaningfulness or even the truth of linguistic expressions can be decided, independently of any context. This is an important problem, but it is not ours. So an analysis of sentences like "the present king of France is bald" cannot serve our

aims, firstly because they are too simple, but mainly because it is hard to find any context in which they could occur.

Even for Russell's standard example "the author of Waverley" I cannot imagine contexts in which it could be conceived as a description of Walter Scott. (Of course from Russell's standpoint this is not at all a disadvantage.) "The author of Waverley" is rather akin to "the Stagirite", "the laughing philosopher", "the king of beasts", "the Queen of the Adriatic", "the rock called the key of the Mediterranean", "the roof of the world", "the poet who is claimed by seven cities as their native", "the knight of the rueful countenance", "der Führer", "wo der Zimmermann das Loch gelassen hat", "l'Incorruptible", "l'Aiglon". These expressions might be called appositions with an omitted support. The name or appellative against which they are supposed to lean is omitted, because it is uniquely determined in a certain context and it is well-known. Appositions with an omitted support are used as a device of poetical or rhetorical expression. At the present stage we are allowed to disregard them. They can also be used as a means of imputation though expressions more akin to real descriptions may better serve this purpose.

Real descriptions can be found in the sentences

The gentleman living over the way is ill.

I have lost the pipe I bought at Smith's.

The first book of the Bible is Genesis.

I did not grasp what he said.

A variable (gentleman, pipe, book of the Bible) is bound by means of a restrictive clause (living over the way, I bought at Smith's) or an adjective (first). The uniqueness of the result is indicated by the definite article (the). Yet sometimes both the variable and the article can be hidden; in the last sentence the variable would be "the thing" and the restriction "that he said". Binding a variable with the definite article only can produce unambiguous descriptions of one thing. In a suitable context

He opened the door

may mean "to open one sharply defined door". Here the definite article is not only the symptom but also the instrument of binding. In another context the same result might be got even by the indefinite article:

I asked a policeman the way

does not point to the existence of policemen that I asked the way ("there is a policeman whom I asked the way"), but at an event of asking one well-defined policeman the way. I could also read the last sentence

I asked the way of the policeman I happened to meet but it would be as good an interpretation to assert that in this sentence a policeman

means

the policeman whom I asked the way.

Then the sentence would run

I asked the way of the policeman whom I asked the way which is not at all a tautology. On the contrary I think that generally in a sentence like

The x with the property A has the property B

the property B may play an essential part in the description that restricts the x; in other words that in many cases such a sentence should be read

The x with the properties A and B has the property B.

This is not surprising, because in factual assertions the explicit description will never be sufficient to bind the variable unambiguously. All the examples we have given can only be understood in a broader context. Yet we cannot admit the broader context as a part of the description, and at the same time reject the nearer context, i.e. postulate the sentence in which the variable occurs formally, shall not have any influence. So even the construction of binding by means of the indefinite article can only be wholly understood if due regard is paid to the context.

As to the use of the article as a means or a symptom of binding (whether the definite article or the indefinite article or no article shall be used) natural languages disagree. There are even languages without articles. People educated in a non-article language find it enormously difficult to learn the use of articles in foreign languages, and users of an article language are often embarassed by ambiguous readings caused by the lack of articles in non-article language text, especially in short contexts.¹) Non-article languages prove that natural languages can do without articles, though article substitutes may occur in non-article languages, such as the Greek article $\tau \delta$, sometimes used in medieval Latin (e.g. in $\tau \delta$ esse, the being) or Russian "to že x" which would probably be used to translate "the x" in

The x with the property A has the property B.

But such occurrences are rather unusual or artificial.

In non-article languages binding of a variable to one value is formalized by means of relative pronouns, adverbs, superlatives, ordinals, demonstrative pronouns, adjectives, and so on. Since in Lincos these categories are not available, we cannot dispense with a general device which formalizes this kind of binding. We have called it article, though this word might

¹⁾ I once translated a Russian book-title "Istoriya zakona" by "History of Law" instead of The history of a law (viz. the law of conservation of energy).

remind the reader too strongly of idiomatic peculiarities of some natural language. Perhaps neutral terms like description-connective and description-binding would have been better than article and article-binding. But this is relatively unimportant. My distinguishing definite and indefinite articles will perhaps evoke more serious criticism. I shall revert to this point.

According to Russell

The x with the property A has the property B is equivalent to

There is one and only one x with the property A, and all x with the property A have the property B.

By this equivalence article binding is reduced to quantification (\vee - and \wedge -binding). Special devices are needed in order to ascertain how this reduction is to be carried out when "the property B" is a complex expression containing negations, quantifications and further article-bound variables. It may be doubtful where the existential quantifier that arises from the reduction of the article is to be put. So according to Russell

The present king of France is bald is false, though

All present kings of France are bald is true, and

The present king of France is not bald is false if it is parallelized to

All present kings of France are not bald, which is true, and it is true if it is parallelized to

Not all present kings of France are bald, which is false.

Paradoxical results like this refute neither Russell nor English. Nevertheless Russell's procedure is misleading. As it has been often imitated, but never questioned, it is worthwhile to analyze it.

Russell's exposition is formulated as if it were an analysis of some phenomena of the English language. But in fact it is an explanation and justification of some habits adopted in *Principia Mathematica*. Here an extensive use is made of interlinear versions into English. Interlinear versions may be very useful when teaching foreign languages, but the part they play in Russell's exposition is not easy to understand. They seem to be used as arguments though it is not clear what they are supposed to prove, especially when they are accompanied by ontological arguments.

Nevertheless they are extremely instructive. Though most of his examples can hardly be put into a reasonable context, Russell does face the problem of lending them a definite meaning or a definite truth value. This is a consequence of formalization.

If problems of meaningfulness must be evaded as much as possible, the ideal solution will be to give some meaning to any expression whatsoever. Formulas clearly borrowed from mathematics like "since so far the expression such and such is meaningless, we are free to endow it with the meaning so-and-so" are not unusual in formalist semantics, though Russell prefers more dogmatic statements. Without denying the importance of eliminating problems of meaningfulness we may stress that our problem is just the inverse, i.e. not to invent some meaning for a given linguistic expression, but to invent linguistic means for sending certain messages.

Russell's analysis of description turns around the problem what

the x with the property A

means if there is no x or if there is more than one x with the property A. According to Russell, who gives ontological arguments, that expression is admissible, but meaningless, and a sentence in which it appears is (roughly speaking) false; according to Hilbert-Bernays both are forbidden; according to Quine it means the void class. All these solutions are possible, though Russell strongly advocates his.

I could not decide in favour of any of them. If a language is designed as a means of communication, it is better not to restrict its possibilities as long as there is no urgent need to do so. Furthermore it is not certain that void descriptions can be dispensed with in communication. Void descriptions are at least an important means of imputation, and since imputation is a not unusual form of social behaviour, we cannot but provide linguistic means to show it. In a certain context "the arrant rogue" may be an apposition with omitted support. In another context it might aim at a real description, though it is not one. In any case it is not clear how such expressions can be avoided, if human behaviour is to be shown. Of course one can imagine a language in which every question beginning with "whether" would have to be answered by "yes" or "no", and where it is not allowed to ask a husband whether he has stopped beating his wife, unless he has been convicted of a conduct like the imputed one in the past. Unfortunately such a language would not reflect human behaviour.

Yet this is not the main point. Closer inspection will show that we cannot apply any formalist theory of description. Such a theory would presuppose the possibility of formal descriptions, but at this stage of

Lincos formal descriptions are lacking. The best we can do is to try descriptions that can be understood by people who endeavour to understand what they are told. Formalist languages have always been dealt with by means of highly-developed non-formalist metalanguages. At a more advanced stage of Lincos we can try to do likewise, but at present it would be rather a waste of time.

The descriptions we shall give of concrete things will possibly be somewhat sharper than they are in natural languages, but they will remain extremely vague as compared with descriptions of mathematical entities. Every speaker will strongly rely on the context and on the good will of the listener.

Our rule for article-binding will not be purely syntactical. "The x", when honestly used, will mean one x; which one must appear from the formal description ("with the property A") together with the nearer and farther context. By article-binding the variable is intentionally, not formally bound to one instance, but the listener may always claim that the intentions of the speaker are clear enough. The result of article-binding must be as sharply defined as it has been in the examples at the beginning of this paragraph, put into a reasonable context. Of course article-binding interpreted in this way cannot be reduced to quantification, which is much more formal.

Mention was made above of a distinction between definite and indefinite article. I cannot deny that this distinction will be somewhat artificial. I shall prefer the indefinite to the definite article ("an x with the property A" to "the x with the property A") if the one and only one x that satisfies the context may be replaced by many formally equivalent other instances so that it does not matter which of them actually satisfies the context. In a suitable context this distinction between indefinite and definite article may be reflected by the pair of English sentences

I opened a door,

I opened the door of my bedroom,

or by the pair

He wrote with a pencil,

He wrote with the pencil I had lent him,

though indefinite articles like these will often be omitted in English

I went by car,

I went by the car of my friend.

It is a natural consequence of this policy that in the case of the definite article the description will be much more formal than in that of the indefinite article.

I may take the liberty of explaining how I arrived at this solution. In the first elaboration of this project I tended to avoid existential binding unless the existence had to be stressed.

I perceived a noise

was transformed into

There is a noise I perceived,

only if the existence of the noise was the main point. I made an extensive and sometimes reckless use of article-binding such as we do in the vernacular. Definite and indefinite article were not distinguished. Article binding is often a much less exact means of expression than quantification. When I looked the work over, I was assailed with doubts. For conscience sake, anxious about exactness of expression in Lincos, I undertook to eliminate the article in many occurrences, especially what could be called indefinite articles. Obviously I was still under the influence of Russell's theory of description. In any case I had not yet adequately clarified my ideas on this matter. Afterwards I happened to notice that a device for indefinite article-binding already existed in Lincos though I had failed to use it. Finally a number of texts have been transformed back by using this device. (Discoveries like this have been made more than once. Lincos proves as it were its independence. One suddenly sees that things not dreamt of beforehand can be expressed by devices created for other purposes. This is a strange sensation. Such a piece of good luck is not easily refused.)

Our rejection of Russell's theory of description is not absolute. Our dissenting from formalist semantics is only a consequence of our dealing with language as a means of communication. A language such as Russell's in which "Walter Scott" and "the author of Waverley" may be freely substituted for each other, cannot serve our purposes, because it does not allow to communicate the fact that Scott is the author of Waverley and to ask the question whether Scott is the author of Waverley. Nor can we do this with a language like Carnap's in which the incidental admissibility of this substitution is decided on the strength of formal criteria. The reader will have remarked a strong influence of intuitionism in our exposition. Against the formalists Brouwer was the first to stress the communicational character of language; but as mathematics can be formalized to a degree, examples drawn from it have failed to provide a convincing argument for who in daily life and in discussing mathematics use unformalized language as a means of communication though they wish to renounce this liberty, when formalization seems to be possible. I can only hope that extra mathematical examples for language as a means of communication will be more convincing.

0 26. The next notion we shall deal with is "convention". In the vernacular this has a fairly definite meaning. In some idioms, especially in that of analysis of science, it has undergone a development which is typical for many misused words. An initially aphoristic use has been interpreted as if it were seriously meant.

The original meaning of "convention" is "agreement" i.e. the act of agreeing as well as the result of agreeing, the thing agreed on. Convention may mean treaty (as in "Geneva convention") but mostly "convention" is less formal than "treaty". This informality is stressed, when "convention" is applied to forms of social behaviour which could have been the result of some explicit engagement (though actually they are the result of education). Fifty years ago it was a convention that women did not smoke in public. Convention now allows women to smoke in public. People more or less tacitly agreed to this change. Smoking or not smoking are not conventional, but smoking or not smoking under certain circumstances may be conventional. It is not conventional as long as it is not a form of social behaviour. There is no convention that people eat, because eating as such is biological, not social behaviour. But meal-times, manners of eating, dishes may be conventional. "Convention" seems always to refer to forms of social behaviour, but it is not true that all forms of social behaviour are conventional. That people do not kill each other is probably not a convention; rather is it true that many conventions have been adopted in order to ensure individual indemnity. Rousseau claimed that social organization was the result of a social contract. Of course this is not true, but it is even less true that the fact of social organization (though a symptom of social behaviour) is a convention. On the other hand many existing means of social organization may be highly conventional.

H. Poincaré in his La Science et l'Hypothèse (1st ed. 1902?) claimed that all general judgements of natural sciences were conventional. At least this proves to be the purport of his argumentation. He deals with different fields of physics, but constantly uses one pattern of argumentation which can meet many more general needs. So later conventionalists declared all general (not mathematical and not purely phenomenological) judgements to be conventional.

Conventionalists undoubtedly hold that their assertion "general judgements are conventional" is true, not merely conventional. Thus the first question we may raise is: "Why does the judgement 'general judgements are conventional' not apply to itself?" A provisional answer might be: This assertion is not conventional, because most people will reject it. This answer shows that the thesis of conventionalists should rather read: General (not mathematical and not purely phenomenological) judgements which are generally believed to be true, are merely conventional.

This gives rise to a second question: How about general judgements put forward by individuals and not generally approved? They are not conventional, but are they true? Was the Copernican world-system true as long as it was not generally admitted, did it become a convention afterwards, and did the Ptolemean system reach the status of truth as a consequence of Copernicus' triumph? We need not answer this question, but if we adopt the conventionalist view, we are obliged to say what kind of thing a convention, e.g. the theory of relativity, was before it was generally agreed on. Clearly conventionalists do not intend to hold that it was a truth. They have probably never considered this problem. Maybe they would answer that before being agreed on, it was a hypothesis (note that "hypothèse" occurs in the title of Poincaré's work, though it is not often used in the book itself). But this solution is impossible. Poincaré and other conventionalists continually speak about linguistic conventions; the counterpart of a linguistic convention would be a linguistic hypothesis, but it is evident that relativity might have been a hypothesis but not a linguistic hypothesis before it became a convention. Perhaps the predecessor of a convention could be called a proposal.

Poincaré and other conventionalists deal with "convention" and "linguistic convention" (convention de language) as if they were synonyms. This is unreasonable. Calling the metric unit of length "metre" was a linguistic convention. Adopting the metric system is not a linguistic, but a metrological convention. Adopting the name "metre" commits us to use the word "metre" under certain conditions, whereas adopting the metric system does not result primarily in certain linguistic behaviour, but in the behaviour of using metric measuring-staffs.

To elaborate: making the first metre-rod from platinum was a convention (a rheological convention); the choice of a solid as the material of the rod was not conventional, but implied by the aim, namely to get a measuring-staff. The actual depository of the standard metre in Sèvres is conventional; it is not a convention that the standard metre is held at a place where the temperature can be kept constant.

Poincaré's starting point was Euclidian geometry. He considered geometry as a physical science ("si donc il n'y avait pas de corps solides, il n'y aurait pas de géométrie"). Poincaré is far from Hilbert's axiomatic point of view. Interpreting Poincaré's work in the sense of axiomatics would be anachronistic. It appeared after Hilbert's Grundlagen der Geometrie, but it is a collection of essays published separately long before. Poincaré became acquainted with the idea of axiomatics only through Hilbert's Grundlagen. This is testified by his review of Hilbert's work where he was struck by Hilbert's questioning other axioms than that of parallels. He was not even familiar with the development of axiomatics

from Helmholtz to Hilbert. The sentence quoted above proves that he agreed with Helmholtz' erroneous criticism of Riemann's work; his ideas about geometry are mainly influenced by Helmholtz. In any case Poincaré's thesis that Euclidian geometry is conventional cannot be interpreted as an adhesion to axiomatics.

Poincaré argues that physical space could be described as well by non-Euclidian as by Euclidian geometry. This is correct as long as the curvature of geometric space is assumed equal to that of physical space (with due account to the precision of measurements). The argument still holds, if one agrees on notions of straight line, congruence and distance which differ from the usual notions. But then it ceases to be an argument. Any true theory will be replaced by another as soon as words used for one kind of things are used for a different kind. Poincaré's answer to this objection would run: In that case the physical notions of straight line, congruence and distance are conventional. To a certain extent this might be correct, but as a matter of fact the notions straight line, congruence and distance, and our physical methods of measuring distances and angles, are implied by our everyday experiences of space, in this sense that the physical notions are refinements of the everyday notions. In his discussion with Poincaré Russell pointed out that even if physical geometry is very loosely related to a very crude everyday geometry, it would not be possible to attach a curvature radius as small as 1 centimeter to the universe. Poincaré would answer once more that everyday geometry is conventional. Yet this is still more improbable. Everyday geometry is exhibited mainly by nonlinguistic means, by bodily movements and by handicraft. The straightline shape of arrows, of fold lines, of walking sticks, of paths and furrows, the circular shape of wheels and millstones, the rectangular shape of doors, the rotational symmetry of vases can hardly believed to be conventional. Of course the words denoting straight lines, circles and so on have been conventional, but as soon as a certain language is agreed upon, they lose this conventional character.

Other fields of physics have been analyzed by Poincaré with the same arguments as geometry. It is a common feature of all these analyses (but in the case of geometry it is less striking) that the rational character of physical theories is overstressed and their roots in everyday physics are rather neglected. He holds that for example the laws of mechanics are conventional definitions for the physical magnitudes related by means of these laws. Clearly such views are rooted in the same soil as definitions in French textbooks of rational mechanics like "mass is nothing but the ratio of force and acceleration". Of course "the number of pages of a sheet" may be said to be a definition of the number 16, but then the fact that such a clause defines a natural number, is still a law, and likewise

the constancy of the said ratio is a non-trivial law. Natural definitions of mechanical notions have to arise from notions of everyday mechanics by a process of refinement. For the mass of material points such a definition would run as follows:

- 1. The mass of a material point is a positive number connected with a unit symbol. It does not depend on the place and velocity of that point.
- 2. Two material points have equal masses if they are interchangeable in mechanical experiments (e.g. interchangeability in a symmetric arrangement).
- 3. The mass of the union of two material points is the sum of the masses of the two points.

By these principles (supplemented with some principle of continuity) mass is fixed, logically and operatively. In mechanical problems these principles are continually used, but textbooks of mechanics do not deign to state such evident truths. The term "axiomatic system" (of mechanics or of electromagnetism) is often used for a collection of fundamental mechanical laws or for Maxwell's equations respectively. If by physical axioms we understand something like the above-mentioned principles for the notion of mass, we cannot but state the lack of an axiomatic system for any field of physics whatsoever. This lack may be made responsible for the usual misunderstandings as to the logical and epistemological status of mechanics and other fields of physics. As long as a few mechanical laws or Maxwell's equations can be believed to provide a sufficient basis for mechanics or electromagnetism, it is also possible to believe that physics is linked to experience by convention. In the chapter "Mechanics" a more careful analysis has been attempted, but the solution there suggested is a rather artificial one owing to the artificial character of our general problem of cosmic communication. A more natural approach is still required.

Physics isolated from experience has been responsible for Poincaré's conventionalism; if physical notions and laws are not viewed as having arisen from everyday physics, it is not astonishing that the bond between both is thought to be conventional. Yet Poincaré's difficulties are still greater. He is puzzled by the phenomenon that no isolated physical law can yield a satisfactory definition of any physical magnitude. So all notions and all laws are declared to be conventional. Poincaré disregards the possibility that nevertheless the total system could fix these notions, logically and operatively. If we remember that Poincaré wrote La Science et l'Hypothèse before Hilbert's Grundlagen, we can easily understand his perplexity. The idea of implicit definition by means of an axiomatic system

dominates Hilbert's work, but even after the publication of the *Grundlagen* it is still rejected by many mathematicians; even Frege believes it to be wrong. So we can understand why Poincaré makes no use of the concept of implicit definition by means of an axiomatic system.

How conventionalism arises from isolating the notions or the facts of a field of knowledge from each other, may be shown by many examples. When isolating one single phenomenon, one is right to hold that asserting the existence of microbes in a drop of water on microscopic evidence is a linguistic convention, and that without denying any fact we could assert the microscope to have produced the spots called microbes in the water. If, however, the action of looking through the microscope and interpreting the image seen, is taken in its usual context of biology and optics, and if that assertion as a linguistic phenomenon is to be taken as plain English, then this verdict is not allowed. Of course the conventionalist may declare biology, optics and English to be conventional and replaceable by some biology*, optics*, and English*, so that in the asterisk context it is true that it is the microscope that has produced the spots called microbes in the drop. Yet this is sooner said than done. As long as nobody has tried building up a biology* and an optics*, calling on them as arguments is premature. English* is an easier case. We should agree on saying "B produces A" instead of "A becomes perceivable by means of B". Then we are committed to introducing a new word "produce*" in order to translate the original "produce" in all its occurrences. English* would be merely coded English. Of course this is a poor result. By the same method one could prove "two and two make four" to be conventional, for it becomes false as soon as I interchange the English words three and four.

Let us imagine somebody who holds that it is a convention to consider all dogs as one species and all cats as another; one could as well agree on all male dogs and male cats as one species and on all female dogs and female cats as another. As long as no other notions and facts are affected, this may be correct. In the context of systematic zoology that assertion is false because it conflicts with the usual sense of "species". The conventionalist can evade this by declaring the usual notion of species to be conventional. He could insist on using "species" in such a way that in the case of dogs and cats it differs from the usual "species". Such a new notion of species would conflict with the bulk of our scientific habits, but for the conventionalist our scientific habits are as conventional as the others. If science were isolated from life, this might be true. But to the extent that science is one of our vital functions and means of subsistence, it is as little conventional as eating and sleeping are.

Isolating some parts of knowledge or language from their context is not bad in itself. It may even be very useful for analysis, just as considering

artificially-isolated physical systems may be useful in physics. But in physics one tries to account for the exterior disturbances when applying the results to a non-isolated real system. It is not only in conventionalist analysis that disregard of context is a weakness. It is a danger that threatens all attempts at analysis.

So far we have not given a proper answer to the first question of his paragraph: Why does the conventionalist not apply his thesis "all general assertions are conventional" to the sentence included within quotation marks? In this form the question asks for a motivation of conventionalist behaviour, which in fact is only a special case of a more general behaviour. People who analyze languages or idions or theories is in order to discover features common to them, often state or postulate general rules which are belied by the language or the idiom they use or by the theory they develop when dealing with languages or idioms or theories (e.g. nominalists forbid general concepts while using a great many general concepts, especially the general concept "general concept"). More or less artificial idealizations, abstractions, dissections, isolations which are applied (perhaps because they cannot but be applied) to languages, idioms, theories as a subject matter, are lacking in the language, the idiom, the theory which is used as a working medium. We may disregard the communicating character of a subject language and the connection of a subject theory with reality, but in the meantime we still use the working language (metalanguage) as a means of communication, and we still believe that the working theory (metatheory) is related to some reality.

This explanation can probably not be applied to Poincaré's behaviour. In fact it has been Poincaré's merit to turn against two interpretations of geometry and mechanics, the aprioristic and the empiristic. He pointed out that geometry and mechanics were neither a priori nor empirical. Initially conventionalism was a merely aphoristic answer to the question what geometry and mechanics are, if they are neither apriori nor empirical. This aphorism has shared the fate of many aphorisms in being taken seriously and repeated more emphatically. While Poincaré discovered "conventions" and "linguistic conventions" everywhere, other conventionalists preferred the terms "arbitrary (linguistic) conventions" and "purely arbitrary (linguistic) conventions", but I doubt whether they sincerely meant that e.g. geometry is purely arbitrary; it seems more probable that they only wished to express their very strong conviction that geometry was conventional.

There is still another way of explaining conventionalist behaviour. In physics and other natural sciences it has often happened that in order to explain some observational facts investigators proposed two or more theories which seemed to fit equally well. Experiments were carried out

in order to decide for one of the competing theories. Gradually one of them became more and more probable, and finally the evidence for this one and against the other one became overwhelming. A certain theory which is now taken to be the true theory, might have been valued only as better or as not worse than other competing theories. So it would be wise to consider it not as a true theory, but as much the best theory now available. This remark is quite reasonable. It matches very well the actual behaviour of investigators who usually work with the best theory as if it were true. If you want to know the value of the elementary charge, a a physicist will answer you: the best value now available is such and such. (Usually the accuracy of a value will be measured by the observational error.) Such an answer does not mean: The value such and such is the present conventional value. It is possible that the only thing conventionalists wish to state is that theories are always questionable. This is correct, but in this respect observational facts are no better.

Valuations like "the best theory now available", though disregarded by analysts, are very important in science. Two correct mathematical theories in the same field need not be equally good; one of them may give more insight into the subject-matter. Among several definitions of the same mathematical object there may be one which is better than the others. Physicists might discuss the best method of arranging a certain experiment. People who work in analysis of scientific knowledge or of scientific language (idiom) are just trying out explanations which organize the subject matter somewhat better than other explanations do. Mostly there is no overwhelming evidence for one of them. So the explanation "a true theory is a theory which is much better than competing theories" is itself at most only much better than other explanations.

The reader might wonder why so much attention has been paid to the conventionalist thesis. If this thesis were correct, our Lincos project would be doomed to failure. If e.g. comparison of time intervals were purely conventional, we could not expect our messages containing a time signal together with the numerical value of its length to be interpreted as such. If Euclidean space as a good approximation of real space were conventional, we could not expect the receiver of our messages about space to interpret the intended object as space. In my opinion physical notions are not arbitrary. This belief is based on the history of science. In earlier times some discoveries might have been accidental. But theories which are current now have been the result of systematic attempts. Kepler and Planck tried a great many hypotheses before they succeeded with the one which has been accepted, not only by logical coercion, but also on the strength of overwhelming evidence. I suppose that these historical facts are little known.

The opinion that scientific theories are completely arbitrary is also held by irrationalists, as a consequence of their fundamental conviction that truth cannot be reached by rational methods, but only by intuition. I refrain from discussing this attitude.

0 27. Finally we shall try to justify our use of the word "meaning" and of related words.

We start with a few artificial dialogues in which the verb "to mean" occurs.

- 1. A. I saw a robot at the exhibition.
 - B. What is a robot?
 - A. "Robot" means something like "golem".
- 2. A. I saw a robot at the exhibition.
 - B. What is a robot?
 - A. "Robot" means a machine which behaves like a human.
- 3. A. What does "Zoll" mean?
 - B. Where did you read it?
 - A. At Hamburg airport.
 - B. There it means the Customs.
- 4. A. Miss Parker is ill.
 - B. Who is Miss Parker?
 - A. I mean our nurse.
- 5. A. She is a nice girl.
 - B. Do you mean the blonde or the brunette?
- 6. A. She is a nice girl.
 - B. Do you mean Miss Parker?
- 7. A. She is a nice girl.
 - B. Do you mean this girl?
- 8. A. This is a robot.
 - B. What do you mean by "robot"?
- 9. A. What did you pay for the car?
 - B. It was very cheap.
 - A. What do you mean by "cheap"?

I have avoided examples like "this cloud means rain", "I mean you no harm", "he meant what he said", in which we meet a semantically quite different verb "to mean".

There can be little doubt about the meaning of "to mean" in our examples. Words, sentences, and other linguistic phenomena may mean something, and what they mean can be shown, or said, by definition or

circumscription in more intelligible terms or by translation into a better-known language. Two slightly different syntactical forms are possible: "X means Y" and "by X somebody means Y". The first looks more objective, but in fact whenever the word "to mean" is used, it is bound to a special situation in a special way. In the 8th example "Do you mean..." refers in a more explicit way to the user of the word "robot" than "to mean" does in the 2nd or 3rd examples. If a more explicit reference to the actual situation is wanted in the explanation of the 2nd example, A could say: If "robot" is used in one context together with "exhibition", it means a machine which behaves like a human. Or: people speaking about an exhibition mean by "robot" a machine which behaves like a human. Or in the 3rd example: By "Zoll" people in Hamburg airport mean "customs".

(The double syntactic pattern of English "to mean" is neither reflected by French "signifier", nor by German "bedeuten". "Vouloir dire" and "bezeichnen" show the same feature as "to mean".)

So far "to mean" raises no special problems. Even the noun "meaning" would not cause difficulties if we used "the meaning of X is Y" as synonymous with "X means Y". Yet there is a strong tendency to assert that a thing which means something has a definite meaning, just as a thing which ends somewhere has a definite ending. "Meaning" (or "denotation" or "nominatum") is used as if it were a function from the set of linguistic entities to that of real entities — somehow this is the leading idea of any attempt of formalist semantics. I would not object to this use of "meaning" if it were suitably restricted and if it did not falsely suggest a much too simple description of the relation between language and reality.

I concede that linguistic expressions generally have some meaning, but I cannot imagine general rules of meaningfulness, and I do not know how to attach just one meaning to every meaningful linguistic expression. I think that in every not-too-poor language the meaning of expressions will strongly depend on the context in which they occur. The importance of the context is often disregarded as soon as language is not primarily considered as a means of communication.

The ancient belief that everything has a true name and that the system of the true names yields the true language, has given way to the insight that there is essentially a variety of languages, and that the meaning of a linguistic expression depends on the language in which it is met. Moreover nobody will seriously deny that in natural languages there are no general rules for stating whether some expression is meaningful and that the meaning of a meaningful expression may depend on its context. If someone asserts the contrary or if he acts as though the contrary were true, it is probable that he has in view not a natural language but some artificial

language and that he has tacitly asked permission to take his examples from a natural language pending the existence of that artificial language. If formalized mathematics is destined to play the part of that artificial language, one may ask why the examples have not been taken from mathematics. Formalized mathematics might be a suitable field to which the formalist semantical notions could be applied, but I doubt whether this is still true if a richer language with more essential relations to reality is to be discussed. I think Carnap is right when he confines his analysis to an artificial language given in advance and precisely described, though I have the impression that his languages are just not rich enough to disclose essential problems.

Formalist semantics is atomistic. It makes the assumption that some elementary expressions have a definite meaning, and that the meaning of composed expressions is derived by certain rules from that of the elementary expressions. If this could be realized, the meaning of any expression could be made independent of the context. In Lincos I have not pursued this aim, mainly because I do not believe in it. Nevertheless I admit that it might be important to fix one's attention on this aim as long as possible.

Proper names seem to be a strong argument in favour of atomistic semantics. But this argument cannot be maintained on closer inspection. Proper names with a definite meaning are extremely rare (I disregard mathematical notions). Socrates, Rembrandt, Walter Scott are good examples, but what about John Smith and Mary Young? William Shakespeare is never mentioned in this connection. In common English it will depend on the context whether "William Shakespeare" means some person known by certain portraits or a certain playwright or both. It is an advantage that in common English one can raise questions like that of the identity of Shakespeare even if they are whimsical. Formalist semantics aim at a language in which this freedom is not allowed.

A moment ago we tackled a well-known problem. Usually it is connected to the name of Walter Scott: Is the same meaning to be ascribed to "Walter Scott" and "the author of Waverley"? This case looks somewhat easier because Scott's authorship of Waverley is usually admitted to be objective truth. The so-called problem of intension in a language with formalist semantics is how to account for the fact that there might be people who are not acquainted with that objective truth or who do not believe it. In fact the case of Walter Scott is not essentially different from that of Shakespeare. In some epoch the identity of the author of Waverley was more obscure than that of the author of Hamlet (e.g. before Scott published Waverley). At some moment Scott's authorship became known. A language in which it will be possible to account for discoveries

of new facts cannot do with expressions which have a definite meaning. Formalist semantics will even not match mathematics, if mathematics is considered not as a stock of true propositions but as a field of invention.

The convention about the meaning of propositions is the strangest feature of formalist semantics. According to Frege a true proposition means truth, a false proposition means falsehood. As opposed to incomplete expressions such as subjects, predicates, and so on, propositions are syntactically independent. Therefore it is easier to put an isolated word "red" or "now" into a virtual context than to do the same with a proposition. The context of incomplete sentences like "it is raining" can also be guessed with relative ease. But what about propositions like "Socrates is a man" or "It rained at Utrecht on December 20th, 1957, at 15 o'clock"? The only context in which "Socrates is a man" can be imagined is the famous syllogism where it is the minor, and in this context the only thing that matters is just to know whether the minor is true. The second sentence might be met with in the report on a bet or something like this, and in this place its meaning might be truth or falsehood indeed. Such arguments may have been responsible for Frege's convention. Moreover if "meaning" is agreed to be a non-trivial function from the set of linguistic expressions which does not depend on the context, the truth function defined for propositions will be welcomed, because it seems to fit very well the requirements of the "meaning" function. In my opinion this does not justify infringing the terminology of the vernacular with respect to "meaning" as radically as Frege did. even if "meaning" is replaced by "denotation" or "nominatum". It is dangerous to create the impression that difficult semantic problems could be solved by a mere convention.

0 28. Attempts to construct symbolic languages have often been linked to some philosophy, though the linking usually proved to be less close afterwards than the author believed. Dealing with language as a means of communication has been the only philosophy I have espoused in the present project. Though devised for a special aim Lincos must be a general-purpose language in which all sayable things can be said. By formal linguistic means we cannot forbid any religious, ethical, philosophical, political or social faith or even any nonsense. So it is better to abstain from such attempts. Our method has been determined by linguistic considerations only. Our pragmatism has been that of an engineer, not of a philosopher. The use of both predicates and sets from the beginning does not prove any partiality in favour of nominalism or realism. The use of the word behaviour does not involve support for behaviorism. Introducing human bodies much later than humans themselves is not to be interpreted as a spiritualist attitude. Defining life as the possibility of

perceiving is not the consequence of any form of sensualism. Lincos words for "good", "bad", "allowed", "forbidden" do not reflect any ethical theory.

Lincos is moderately formalized, but we do not object to fully-formalized language. Foolproof languages have an importance of their own. I hope Lincos will be wise-proof. Formalized languages appeal to the malevolent reader. The present book is addressed: lectori benevolenti.

Utrecht, December 23rd, 1957.

MATHEMATICS

1 00 0. Pairs of signs * will enclose the printed image of a program text.

A (metatextual) "and so on" after a text indicates that this text is an exemplary extract from the factual program. When carrying out the program we will replace this text by a large number of texts similar to the text we have printed. If the number of examples is large enough, we may expect that the receiver can generalize the quasi-general definition or proposition that is intended by the program text.

Loose Lincos words are presented, without any context, in order to stress their individuality. So it will be somewhat easier for the receiver to recognize them when they occur in a certain context.

The bold-faced strokes mean pauses.

1 01 1. * · · · · > · · * and so on.

1 01 2. * · · · < · · · * and so on.

1 01 3. * · · · = · · · * and so on.

1 01 4. * · · · + · = · · · · * and so on.

1 01 5. * · · · · · · · · · · * and so on.

In these texts the Lincos phoneme that corresponds to the round dot is a short radio-signal (a peep). A Lincos word that consists of n successive phonemes of this kind, separated by short and equal intervals, is written as a group of n round dots. It both means and shows the natural number n. It is an ideophonetic word, which has the power of an image as well as that of a word. We also call it an ostensive numeral. The greater part of the Lincos vocabulary will be purely conventional; words may be permutated at pleasure.

This is not true of ideophonetic words. Their essential features must not be changed.

The Lincos words written >, <, =, +, -, and so on designate connectives with the usual meaning. The receiver should guess their meaning from the context. Therefore each of the first program texts contains one unknown word only.

1 02 1.
$$* \cdot = 1$$

 $\cdot \cdot = 10$
 $\cdot \cdot \cdot = 11$
 $\cdot \cdot \cdot \cdot = 100$
 $\cdot \cdot \cdot \cdot \cdot = 101$
 $\cdot \cdot \cdot \cdot \cdot \cdot \cdot = 1101 *$
and so on.

Ostensive numerals are superseded by algorithmic ones, composed of syllables written 0 and 1 according to the rules of the dyadic positional system. For the convenience of the terrestrial reader we shall sometimes use the decimal code, but as a matter of fact such occurrences should be translated into the dyadic code.

1022. Texts as those of 1011-1015 will be repeated, using algorithmic numerals instead of ostensive ones.

1031.
$$#111 - 110 + 1 - 101 + 10 = 100 + 11 = 11 + 100 = 10 + 101 = 1 + 110 # and so on.$$

1 03 2.
$$*111 + 11 > 11 + 101 > 1 + 100 = 101 * and so on.$$

In elementary arithmetics expressions such as "4+3" are often considered as problems, not as numbers; "4+3" is read as if it were "add four and three", "4+3=7" as if it were "if I add four and three, I get seven". Our program texts will prevent such interpretations. The receiver will learn the *relational symmetry* of the Lincos word written =, as well as the "static" character of expressions as '111+11'.

1 04 1.
$$*100 > 10$$

 $100 + 11 > 10 + 11$
 $100 + 1101 > 10 + 1101$
 $100 + 1 > 10 + 1$
 $100 + 110 > 10 + 110$
 $100 + 11111 > 10 + 11111$
 $100 + a > 10 + a *$
and so on.

By such texts we introduce Lincos words that mean variables. In Lincos speech words meaning variables will be distinguished from other words by certain phonetic characteristics.

104 2. * 100 + 111 = 111 + 100

$$100 + 1 = 1 + 100$$

 $100 + 1101 = 1101 + 100$
 $100 + 11 = 11 + 100$
 $100 + a = a + 100 *$
and so on.

1043.
$$*a + 1110 = 1110 + a$$

 $a + 11 = 11 + a$
 $a + 11011 = 11011 + a$
 $a + 1 = 1 + a$
 $a + b = b + a *$
and so on.

- 1 04 4. Texts as those of 1 04 1-1 04 3 will be repeated using other words that mean variables. A sufficiently large stock of such words should be provided for.
- 1050. So far no explicit attention has been paid to punctuation. In fact we can dispense with an exposition of the means and rules of punctuation. In speech, we mark punctuation mainly by pauses, and we will do so in Lincos speech as well. Pauses will be used as Lincos punctuations. It is understood that the longer pause will have the stronger power of separation, and the receiver will understand this tacitly imposed rule.

In the printed image spaces would be the equivalent of pauses, but such texts would annoy the reader. Russell-Whitehead square dots are much better as long as the number of dots in one punctuation mark does not exceed six. Unfortunately seven- and eight-point marks will not be unusual in our texts; we shall even meet with eleven- and twelve-points marks. So we have devised a variation on the Russell-Whitehead method:

A square dot will account for 1, 2, 4 punctuation units according as the level of its placing is the bottom, the middle, or the top of the letters b, d, h and so on. A short vertical bar will account for 8 units, a long vertical bar for 16 units, and a long horizontal bar for 32 units. Every group of such elements represents one punctuation mark. In each group the units are added, in order to get the punctuation value of the corresponding mark. The separative power of a mark will be a monotonously increasing function of its punctuation value.

We give a list of punctuations with their values:

We shall not use all these marks. In *algebraic* formulas we shall often stick to the classical procedure (brackets with the usual conventions), but this is rather a concession to the terrestrial reader.

In our system there is a great variety of punctuation marks. So there is no need to economise punctuations, as there is in the Russell-Whitehead system. We can also avoid another drawback of that system and of other systems: the use of logical connectives as punctuation marks, and artificial conventions regarding their separative power.

Let us emphasize once more that these considerations do not affect Lincos speech, but only written Lincos.

105 1.
$$*101+11.+111 \cdot = \cdot 101+.11+111=.101+11+111:$$
 $101+11.+1 \cdot = \cdot 101+.11+1 \cdot = \cdot 101+11+1:$
 $101+11.+101 \cdot = \cdot 101+.11+101 \cdot = \cdot 101+11+101:$
 $101+11.+c \cdot = \cdot 101+.11+c \cdot = \cdot 101+11+c:$
 $101+1001.+c \cdot = \cdot 101+.1001+c \cdot = \cdot 101+1001+c:$
 $101+b.+c \cdot = \cdot 101+.b+c \cdot = \cdot 101+b+c:$
 $1101+b.+c \cdot = \cdot 1101+.b+c \cdot = \cdot 1101+b+c:$
 $a+b+c \cdot = \cdot a+b+c \cdot = \cdot a+b+c \cdot =$
 $105 2. *a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$
 $\cdot a+b+c \cdot + d \cdot = \cdot a+b+c \cdot + d \cdot =$

 $a > 111 \rightarrow 10 < a$

and so on.

1064.
$$*a>b.\rightarrow a+1.>b:$$

 $a>b.\rightarrow a+10.>b:$
 $a>b.\rightarrow a+110.>b*$
and so on.

106 5.
$$*a < b \rightarrow a < b + 1$$
:
 $a < b \rightarrow a < b + 111$:
 $a < b \rightarrow a < b + 10$ *
and so on.

1 06 6.
$$*a=b \cdot \rightarrow \cdot b=a$$

 $a>b \cdot \rightarrow \cdot b
 aa :
 $a>b \cdot \rightarrow \cdot b<\cdot a+101$ *
and so on.$

1067.
$$*a+b.=c. \rightarrow a+b+1.>c:$$

 $a+b.=c. \rightarrow a+b+11.>c:$
 $a+b.=c. \rightarrow a+b.<.c+1:$
 $a+b.=c. \rightarrow a+b.<.c+110*$
and so on.

106 8.
$$*a+b = c \rightarrow a = c-b$$
:
 $a+b = c \rightarrow b = c-a$:
 $a=c-b \rightarrow a+b = c$ *
and so on.

At 1061 the first connective with a logical meaning has occurred. In order to avoid misinterpretations we have inserted texts in which the '--' could be replaced by '--'. We have deliberately not admitted texts like

$$10+1 = 11 \cdot \rightarrow \cdot 11+1 = 100$$
:
 $10+1 = 101 \cdot \rightarrow \cdot 11+1 = 100$:
 $10+1 = 101 \cdot \rightarrow \cdot 11+10 = 100$:
 $1+1 = 10 \cdot \rightarrow \cdot a+b = b+a$

and so on. Such texts are meaningless for any receiver who is not acquainted with the formal use of implication.

106 9.
$$*a>11.\rightarrow:b>10.\rightarrow:a+b.>110$$

 $a>c.\rightarrow:b>d.\rightarrow:a+b.>.c+d$
 $a>b.\rightarrow:b>c.\rightarrow:a>c$
 $a=c.\rightarrow:b=d.\rightarrow:a+b.=.c+d$
 $a=b.\rightarrow:b=c.\rightarrow:a=c$
and so on.

107 1. *
$$a > 100 . \leftrightarrow . a > 100$$
:
 $a > 100 . \leftrightarrow . 100 < a$:
 $a > b . \leftrightarrow . b < a$:
 $a > b . \leftrightarrow . a + 1 . > . b + 1$:
 $a > b . \leftrightarrow . a + c . > . b + c$:
 $a + b . = c . \leftrightarrow . a = . c - b$ *
and so on.

1072.
$$*a=b \rightarrow c=d \rightarrow c=d \rightarrow a=b \rightarrow a=b \rightarrow c=d$$

The tautology $p \rightarrow q \rightarrow q \rightarrow p \rightarrow p \rightarrow p$ has been stated in an exemplary way. As long as we have no words that designate propositions, this will be the way of formulating tautologies.

1 08 1. *
$$1 \neq 10.1 \neq 11.1 \neq 100.1 \neq 110111$$
.
 $10 \neq 1.10 \neq 11.10 \neq 1010$.
 $11 \neq 1.11 \neq 10.11 \neq 111$.
 $1101 \neq 11011.110 \neq 111.111010 \neq 1100$ * and so on.

1082.
$$*a > b \rightarrow a \neq b$$

 $a < b \rightarrow a \neq b$
 $a \neq b \rightarrow a + c \neq b + c$
and so on.

We try to avoid formulas that no longer hold true when the domain of the variables is extended.

1 09 1.
$$*a > 1 \cdot \lor \cdot a < 10 \cdot a > 10 \cdot \lor \cdot a < 11 \cdot a > 111 \cdot \lor \cdot a < 1000 \cdot a > 11 \cdot \lor \cdot a < 1000 \cdot a > 1 \cdot \lor \cdot a < 11 \cdot a > 10 \cdot \lor \cdot a < 100 \cdot a > 1110 \cdot \lor \cdot a < 110001 * and so on.$$

1 09 2. *
$$a \neq 101. \rightarrow `a < 101. \lor .a > 101:$$
 $a \neq b. \rightarrow `a < b. \lor .a > b:$
 $a = b. \lor .a < b. \lor .a > b:$
 $a = b. \lor .a < b. \lor .a > b:$
 $a = 10. \lor .a = 1001. \lor .a = 111. \rightarrow .a < 1011:$
 $a = 10. \lor .a = 1001. \lor .a = 111. \rightarrow .a > 1:$
 $a = 10. \rightarrow `a = 10. \lor .a = 111. \oplus .a > 1:$
and so on.

1093.
$$*a=b. \rightarrow :a=b. \lor. c=d:$$
 $a=b. \lor. c=d: \leftrightarrow :c=d. \lor. a=b:$
 $a=b. \lor. c=d: \leftrightarrow :a=b. \rightarrow :c=d:$
 $a=b. \lor. c=d: \lor. e=f: \leftrightarrow :a=b. \lor. c=d. \lor. e=f:$
 $a=b. \lor. c=d: \lor. e=f: \leftrightarrow :a=b. \lor. c=d. \lor. e=f*$
and so on.

Exemplary tautologies. See 1 07 2.

1 10 1.
$$*1 \le 1.1 \le 10.1 \le 11.1 \le 10101$$
.
 $10 \le 10.10 \le 11.10 \le 111$.
 $11 \le 11.11 \le 11011$.
 $1 \ge 1.10 \ge 1.11 \ge 1.1100 \ge 1$.
 $10 \ge 10.111 \ge 10$.
 $1110 \ge 101 *$
and so on.

1 10 2.
$$\# a \leq b . \Leftrightarrow `a = b . \lor . a < b :$$
 $a \geq b . \Leftrightarrow `a = b . \lor . a > b :$
 $a \leq b . \lor . a \geq b #$
and so on.

1 11 1.
$$*a = 111. \rightarrow .a \ge 10$$
:
 $a = 111. \rightarrow .a \le 11001$:
 $a = 111. \rightarrow .a \ge 10. \land .a \le 11001$:
 $a \ge 10. \land .a \le 111. \leftrightarrow .a = 10. \lor .a = 11.$
 $\lor .a = 100. \lor .a = 101. \lor .a = 110. \lor .a = 111.$
and so on

1 11 2.
$$*a=111. \leftrightarrow `a \le 111. \land .a \ge 111:$$

$$a=b. \leftrightarrow `a \le b. \land .b \le a *$$
and so on.

1 11 3. *
$$a \le 111 \cdot \land \cdot b \le 10 \cdot \rightarrow \cdot a + b \cdot \le 1001$$
:
 $a > 110 \cdot \land \cdot b \ge 10 \cdot \rightarrow \cdot a + b \cdot > 1000$:
 $a \le b \cdot \land \cdot c \le d \cdot \rightarrow \cdot a + c \cdot \le \cdot b + d$:
 $a > b \cdot \land \cdot c \ge d \cdot \rightarrow \cdot a + c \cdot > \cdot b + d$ *

1114.
$$*a=b \cdot \land \cdot c=d \cdot \rightarrow \cdot a=b$$
 $a=b \cdot \land \cdot c=d \cdot \rightarrow \cdot c=d \cdot \land \cdot a=b$ $a=b \cdot \land \cdot c=d \cdot \land \cdot e=f : \Leftrightarrow : a=b \cdot \land \cdot c=d \cdot \land e=f$ $a=b \cdot \land \cdot c=d \cdot \land \cdot e=f : \Leftrightarrow : a=b \cdot \land \cdot c=d \cdot \land \cdot e=f$ $a=b \cdot \rightarrow \cdot e=f \cdot \land \cdot c=d \cdot \rightarrow \cdot e=f : \Rightarrow : a=b \cdot \lor \cdot c=d \cdot \rightarrow \cdot e=f$ $a=b \cdot \rightarrow \cdot c=d \cdot \land \cdot c=d \cdot \rightarrow \cdot e=f : \Rightarrow : a=b \cdot \rightarrow \cdot e=f$ $a=b \cdot \rightarrow \cdot c=d \cdot \rightarrow \cdot e=f : \Leftrightarrow : a=b \cdot \land \cdot c=d \cdot \rightarrow \cdot e=f$ and so on.

1 12 1. *
$$?x \cdot x + 10 = 111 \cdot x + 10 = 111 \cdot x \cdot x = 101 \cdot ?x \cdot a < b \cdot \land \land x + a \cdot = b \cdot a < b \cdot \land \land x + a \cdot = b \cdot \Rightarrow \land x = b - a * and so on.$$

In symbolic logic little attention has been paid to sentential forms other than propositional ones. Problems and questions, though indispensable in mathematics and life, cannot be formalized by the usual symbolic methods. The philosophy that mathematics is rather a stock of true propositions than an art of discovering is responsible for this omission.

We have introduced a word written as an interrogation mark in order to formalize *questions*.

1 13 1.
$$\# ?x \cdot x + 101 = 11 \#$$

and so on. No answer is given.

1133.
$$*a-a = 0$$

 $a < b \rightarrow : a-b = -b-a$
and so on.

1 13 4.
$$*?x \cdot x + a \cdot = b$$
:
 $x + a \cdot = b \cdot \rightarrow \cdot x = \cdot b - a + a$
and so on.

1 13 5. *
$$10+11=101$$
:
 $1+11.=100$:
 $0+11.=11$:
 $-1.+11.=10$:
 $-10.+11.=1$:
 $-11.+11.=0$:
 $-100.+11.=-1$ *
and so on.

Of course different words must be used for the "minus" operating on one object (in -a) and the "minus" operating on a pair of objects (in a-b).

Addition and substraction laws for integers will be developed in detail by the method of quasi-induction. We omit these texts.

1 13 6. *
$$0 \neq 1:0 \neq 10:0 \neq 11:0 \neq 1101:$$

 $-1. \neq 0:-1. \neq 1:-1. \neq 10:-1. \neq 11101:$
 $-10. \neq 0:-10. \neq 101:$
 $-100. \neq .-1010 *$
and so on.

1 13 8.
$$*a < b \cdot \leftrightarrow -b \cdot > -a$$
:
 $a > 0 \cdot \leftrightarrow -a \cdot < 0$:
 $a > b \cdot \leftrightarrow a - b \cdot > 0$:
 $- \cdot -a \cdot = a *$
and so on.

1 13 9. Some earlier formulas will be repeated in order to suggest that variables now range through the system of all integers.

1 14 1.
$$*101 \times 1 = 101$$
.
 $101 \times 10 = 1010$.
 $101 \times 11 = 1111$.
 $101 \times 100 = 10100$.
 $101 \times 101 = 11001$.
 $101 \times 1010 = 110010 *$
and so on.

1 14 2. *
$$101 \times 0 = 0$$
 * $101 \times ... = 1 = ... = 101$ * $101 \times ... = 10 = ... = 1010$ * $101 \times ... = 1010 = ... = 110010$ * and so on.

1 14 3.
$$*a \times 0 = 0$$
:
 $a \times b + 1 = a \times b + a$:
 $a \times b - 1 = a \times b - a$
and so on.

1 14 4.
$$*a \times b = b \times a$$
:
 $a \times .b \times c \cdot = `a \times b . \times c = .a \times b \times c$:
 $a + b . \times c \cdot = `a \times c . + .b \times c$ *
and so on.

For the convenience of the reader of written Lincos we shall now adopt the habit of omitting the multiplication sign and of using brackets in algebraic contexts.

1 14 5. *
$$ab = 0$$
 \leftrightarrow ` $a = 0$ \lor ` $b = 0$ `
$$ab \neq 0 . \leftrightarrow$$
 ` $a \neq 0$ \land ` $b \neq 0$ `
$$ab > 0 . \leftrightarrow$$
 : $a > 0$ \land ` $b > 0$ \lor ` $a < 0$ \land ` $b < 0$ `
$$aa \ge 0 *$$
and so on.

1 15 1.
$$?x.11x=110$$
.
 $11x=110. \rightarrow .x=10$:
 $?x.11x=1100$.
 $11x=1100. \rightarrow .x=100$:
 $?x.11x=100$.
 $11x=100. \rightarrow .x=100/11$ # and so on.

1 15 2. * 11
$$x = 100 \cdot \Leftrightarrow .110 x = 1000 \cdot 100/11 = 1000/110 = 1100/1001 = 11000/10010 = (-100)/(-11) = (-1000)/(-110) * and so on.$$

1 15 3. #
$$?x \cdot a \neq 0$$
. $\land .ax = b$:
 $a \neq 0$. $\land .ax = b \cdot \rightarrow .x = b/a$ #

1 15 4. *
$$ax = b \cdot \rightarrow \cdot cax = cb \cdot c \neq 0 \cdot \rightarrow \cdot b/a = (cb)/(ca)$$
 *

1 15 5.
$$\# a/1 = a \#$$

1 15 6. * 110/101.>0.

$$110/(-101).<0.$$

 $ab>0. \Leftrightarrow a/b>0.$
 $ab<0. \Leftrightarrow a/b<0.$

55

1 15 7. In an analogous way arithmetics of fractions will be exposed in all necessary details.

```
1 16 1. *1100/10 = 110.

110/10 = 11.

11/10 = 1,1.

1,1/10 = 0,11 *
```

and so on.

1 16 2.
$$*1/10=0,1$$
.
 $1/100=0,01$.
 $1/100000=0,00001$ * and so on.

1 16 3. *1,101101 = 1 + 1/10 + 1/1000 + 1/10000 + 1/1000000 * and so on.

```
1 16 4. * 1/11 = 0.01010101 Etc.

1/101 = 0.0011001100110011 Etc.

1/110 = 0.001010101 Etc.

1/111 = 0.001001001001 Etc.

1/1001 = 0.000111000111000111 Etc.

10/11 = 0.1010101010 Etc.

1000/11 = 10.1010101010 Etc:

(1000/11) - (10/11) = \cdot 10.101010 Etc. - . 0.101010 Etc. = 10.0000 Etc = 10:

11/111 = 0.011011011011 Etc.

(11/111) - (1/1001) = 0.010100010100010100 Etc.

(11/111) + (1/1001) = 0.100010100010100010 Etc.

10(11/111) = 0.110110110110 Etc.
```

This is an important step. A Lincos word meaning "and so on" and written Etc is introduced. I doubt whether it could be done at an earlier stage, if we wish to avoid ambiguous interpretations. The child that computes its first periodic fraction faces *infinity*. By showing periodic fractions we get the most natural opportunity of introducing 'Etc'.

'Etc' is to correspond to the "..." of common mathematical language. Note that 'Etc' should not be confused with the metatextual "and so on".

A few examples will illustrate the use of 'Etc'.

1 16 5.
$$*1+10=0,1\times10\times11.$$

 $1+10+11=0,1\times11\times100.$
 $1+10+11+100=0,1\times100\times101.$
Etc.
 $1+10+11+\text{Etc}+(a-1)+a=0,1\times a\ (a+1)*$

1 16 6. *
$$a^0 = 1$$
.
 $a^1 = aa^0 = a$.
 $a^{10} = aa^1 = aa$.
 $a^{11} = aa^{10} = aaa$.
Etc.
 $a^{n+1} = aa^n = aa \times \text{Etc} \times a$:
 $a^{-n} = 1/a^n$ *

1 16 7. *
$$a \neq 1$$
. \rightarrow . $1 + a + a^{10} + a^{11} + \text{Etc} + a^n = (1 - a^{n+1})/(1 - a)$.
 $-1 < a < 1$. \rightarrow . $1 + a + a^{10} + a^{11} + \text{Etc} = 1/(1 - a)$.

1 17 1. * 1 ∈ Num.10 ∈ Num.11 ∈ Num.100 ∈ Num.101 ∈ Num. Etc. 0 ∈ Int.1 ∈ Int. − 1 ∈ Int.10 ∈ Int. − 10 ∈ Int.11 ∈ Int. Etc.*

The Lincos word written Num (fL^1) numerus) means "natural numbers". The Lincos word written Int (fL) integer) means "integers". Words designating sets have been introduced when at last two sets can be shown. As a rule we shall not name a particular kind of thing unless things not of this kind are known.

1 17 2. *
$$a \in \text{Num} . \Leftrightarrow `a = 1. \lor . a = 10. \lor . a = 11. \lor . a = 100.$$

$$\lor . a = 101. \lor \text{Etc}:$$

$$a \in \text{Int} . \Leftrightarrow `a = 0. \lor . a = 1. \lor . a = -1. \lor . a = 10.$$

$$\lor . a = -10. \lor . a = 11. \lor . \text{Etc} *$$

1 17 3.
$$*a \in \text{Num}. \Leftrightarrow `a \in \text{Int}. \land .a > 0$$
'
 $a \in \text{Num}. \Rightarrow `a + 1. \in \text{Num}$ '
 $a \in \text{Int}. \Leftrightarrow :a + 1. \in \text{Int}` \land `a - 1. \in \text{Int}$ '
 $a \in \text{Int}. \Leftrightarrow `-a. \in \text{Int}$ '
 $a \in \text{Num}. \land .b \in \text{Num}` \Rightarrow `a + b. \in \text{Num}$ '
 $a \in \text{Num}. \land .b \in \text{Num}: \Rightarrow :a - b. \in \text{Num}` \Leftrightarrow .a > b$ *

^{1) &}quot;fL" means "from Latin". Most written Lincos words are abbreviations of Latin words or word-groups.

```
1 18 1. \# x/1 \cdot \in \text{Int} : \Leftrightarrow x = 0 \cdot y \cdot x = 1 \cdot y \cdot x = -1 \cdot y \cdot \text{Etc}:
                                                                               x/10 \in \text{Int} \leftrightarrow x=0 \quad \forall x=10 \quad \forall x=-10 \quad \forall x \in \text{Etc}
                                                                               x/11 \in \text{Int} \leftrightarrow x=0 \cdot x=11 \cdot x=-11 \cdot x
                                                                               Etc.
                                                                               x/1 \in \operatorname{Int} \cdot \wedge \cdot x/10 \in \operatorname{Int} \cdot \wedge \cdot x/11 \in \operatorname{Int} :
                                                                                                  \leftrightarrow x=0 \lor x=110 \lor x=-110 \lor Etc
                                                                               x/1 \in \text{Int} \cdot \wedge \cdot x/10 \in \text{Int} \cdot \wedge \cdot x/11 \in \text{Int} \cdot \wedge \cdot x/100 \in \text{Int}:
                                                                                                   \Leftrightarrow x = 0, \forall x = 1100, \forall x = -1100, \forall Etc
                                                                              Etc.
                                                                            x/1 \cdot \in \text{Int} \cdot \wedge \cdot x/10 \cdot \in \text{Int} \cdot \wedge \cdot x/11 \cdot \in \text{Int} \wedge \cdot x/100 \cdot \in \text{Int} \cdot \wedge x/100 \cdot \in \text{Int} \cdot \wedge
                                                                                                x/101 \in \text{Int} \cdot \land \cdot \text{Etc} : \leftrightarrow \cdot x = 0.
                                                                            a=1. \rightarrow x/a. \in Int: \land : a=10. \rightarrow x/a. \in Int:
                                                                                                  \wedge : a = 11 . \rightarrow : x/a . \in Int : \wedge : Etc
                                                                                                                   \Leftrightarrow : a \in \text{Num} \rightarrow x/a \in \text{Int}:
                                                                              x/1 \in \text{Int} \cdot \wedge \cdot x/10 \in \text{Int} \cdot \wedge \cdot x/11 \in \text{Int} \cdot \wedge \cdot \text{Etc}:
                                                                                                  \leftrightarrow \land a : a \in \text{Num} \rightarrow x/a \in \text{Int}
                                                                               \land a: a \in \text{Num} \rightarrow x/a \in \text{Int} \leftrightarrow x = 0 *
```

A connective written \wedge and meaning "for every" has arisen from an infinite conjunction. This seems to be the most natural way of introduction. Some phonetic resemblance between ' \wedge ' and ' \wedge ' would be desirable.

Texts such as

*
$$a = 1 . \rightarrow . a^{10} > 0 \cdot \land \cdot a = 10 . \rightarrow . a^{10} > 10 \cdot \land \cdot a = 11 .$$

 $\rightarrow . a^{10} > 0 \cdot \land \cdot \text{Etc}:$
 $\land a \cdot a \in \text{Num}. \rightarrow . a^{10} > 0 *$

could hardly serve to introduce ' \wedge '. Remember we have always been considering

$$a \in \text{Num} \rightarrow a^{10} > 0$$

as a (true) proposition. This means that we have tacitly agreed to generalize over free variables. (It would be difficult to avoid this convention.) If we should start with the last program text, the receiver would not understand why we at once add the word written ' \wedge ', and he would be unable to guess its meaning. We are obliged to introduce the word written \wedge by means of a proposition, in which it is not syntactically superordinated to the rest of the proposition. Upon this point of view we choose our example. After this we are allowed to superordinate the connective written \wedge to a whole sentence too.

It is advisable to give more examples:

1 18 2. *
$$0 \notin \text{Num}: -1. \notin \text{Num} \cdot -10. \notin \text{Num} \cdot -11. \notin \text{Num} \cdot -110. \notin \text{Num} \cdot \text{Etc}$$

$$-100. \notin \text{Num} \cdot \text{Etc} \cdot 11/111. \notin \text{Num} \cdot -111/10. \notin \text{Num} \cdot \text{Etc}$$

$$a \notin \text{Num}. \leftrightarrow : a \neq 1. \land . a \neq 10. \land . a \neq 11. \land . a \neq 100. \land . \text{Etc}$$

$$a \notin \text{Num}. \to : x = 1. \to . a \neq x \cdot \land \cdot x = 10. \to . a \neq x$$

$$\land \cdot x = 11. \to . a \neq x \cdot \land \cdot x = 100. \to . a \neq x \cdot \land \cdot \text{Etc}$$

$$a \notin \text{Num}. \leftrightarrow : \land x \cdot x \in \text{Num}. \to . a \neq x *$$

1 18 3.
$$* \land x \cdot x \neq a \cdot \lor x \neq b : \rightarrow a \neq b *$$

1 18 5. *
$$a = 1 \cdot \land \cdot x = 10a \cdot \lor \cdot a = 10 \cdot \land \cdot x = 10a \cdot \lor \cdot a = 11 \cdot \land \cdot x = 10a \cdot \lor \cdot a = 100 \cdot \land \cdot x = 10a \cdot \lor \cdot Etc:$$

$$\Leftrightarrow \cdot a \in \text{Num} \cdot \land \cdot x = 10a \cdot \lor \cdot x = 10 \times 11 \cdot \lor \cdot x = 10 \times 100 \cdot \lor \cdot Etc:$$

$$\Leftrightarrow : \forall a \cdot a \in \text{Num} \cdot \land \cdot x = 10a^{+}$$

A connective written \vee and meaning "there is a" has arisen from an infinite disjunction. The case is entirely analogous to that of 1181.

1 18 6. *
$$y = 1 \cdot \land . x = y^{10} \cdot \lor \cdot y = 10 \cdot \land . x = y^{10} \cdot \lor \cdot y = 11 \cdot \land . x = y^{10} \cdot \lor \cdot y = 100 \cdot \land . x = y^{10} \cdot \lor \cdot \text{Etc} : \leftrightarrow \cdot y \in \text{Num} \cdot \land . x = y^{10} \cdot x = 1 \cdot \lor . x = 1000 \cdot \lor . x = 10000 \cdot \lor . \text{Etc} \cdot \leftrightarrow : \forall y \cdot y \in \text{Num} \cdot \land . x = y^{10} \cdot \lor \cdot y = 1 \cdot \land . x = y^{10} \cdot \lor \cdot y = -1 \cdot \land . x = y^{10} \cdot \lor \cdot y = 10 \cdot \land . x = y^{10} \cdot \lor \cdot \text{Etc} : \leftrightarrow \cdot y \in \text{Int} \cdot \land . x = y^{10} \cdot \lor \cdot x = 10000 \cdot \lor . \text{Etc} \cdot \leftrightarrow : \forall y \cdot y \in \text{Int} \cdot \land . x = y^{10} *$$

The last converted in the property of the property o

The last example is to prevent the misinterpretation that the 'V' would meand "there is one".

```
1 18 8. # \forall x \cdot \text{Etc.} x \cdot \text{Etc.} \rightarrow : \text{Etc.} 0 \cdot \text{Etc.} \lor \cdot \text{Etc.} 1 \cdot \text{Etc.} \lor \cdot \text{Etc.} 1

Etc. - 1 . Etc · \lor \cdot \text{Etc.} 101 \cdot \text{Etc.} \lor \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{Etc.} \bullet \cdot \text{Etc.} 10/111 \cdot \text{Etc.} \lor \cdot \text{Etc.} \bullet \cdot \text{E
```

1 19 1. $a \in \text{Int.} \land .b \in \text{Int.} \land : \land x \cdot x \in \text{Int.} \rightarrow .x^{10} + 10 a x + b > 0$ $\Rightarrow : z = 1. \Rightarrow .z^{10} + 10 a z + b > 0$ $\land \cdot z = 10. \Rightarrow .z^{10} + 10 a z + b > 0$ $\land \cdot \text{Etc.} \land \cdot z = -a. \Rightarrow .z^{10} + 10 a z + b > 0$ $a \in \text{Int.} \land .b \in \text{Int.} \land : \land x \cdot x \in \text{Int.} \rightarrow .x^{10} + 10 a x + b > 0$ $\Rightarrow .b > a^{10}$: $a \in \text{Int.} \land .b \in \text{Int.}$ $\Rightarrow : \land x \cdot x \in \text{Int.} \rightarrow .x^{10} + 10 a x + b > 0$ $\Rightarrow \cdot b > a^{10}$: $a \in \text{Int.} \land .b \in \text{Int.} \land .b > a^{10} \Rightarrow \land x \cdot x \in \text{Int.}$ $\Rightarrow \cdot x^{10} + 10 a x + b > x^{10} + 10 a x + a^{10} = (x + a)^{10} \ge a^{10} \ge 0$ $a \in \text{Int.} \land .b \in \text{Int.}$ $\Rightarrow : b > a^{10}. \Rightarrow \land x \cdot x \in \text{Int.} \Rightarrow \cdot x^{10} + 10 a x + b > 0$ $a \in \text{Int.} \land .b \in \text{Int.}$ $\Rightarrow : b > a^{10}. \Rightarrow \land x \cdot x \in \text{Int.} \Rightarrow \cdot x^{10} + 10 a x + b > 0$

- 1 19 2. # $a \in \text{Num} \rightarrow : \forall x \cdot x \in \text{Num} \land a \leq x^{10} \leq 10 a$ #
- 1 20 0. A few new words are introduced, in order to provide for a larger stock of sets.
- 1 20 1. * 10 Div 10.10 Div 100.10 Div 110. Etc.

 11 Div 11.11 Div 110.11 Div 1001. Etc.

 100 Div 100.100 Div 1000. Etc.

 1011 Div 1001101. Etc.

 10 Div $a \cdot \Leftrightarrow \cdot a/10 \cdot \in \text{Num}$:

 11 Div $a \cdot \Leftrightarrow \cdot a/11 \cdot \in \text{Num}$:

 Etc.: $b \text{ Div } a \cdot \Leftrightarrow \cdot b \in \text{Num} \cdot \land \cdot a/b \cdot \in \text{Num}$: $a \in \text{Num} \cdot \Rightarrow \cdot 1 \text{ Div } a \cdot \land \cdot a \text{ Div } a *$

The word written Div (fL dividit) means "divides".

```
1 20 2. \# 10 \in Pri. 11 \in Pri. 101 \in Pri. 111 \in Pri. 1011 \in Pri. Etc.
                   1 ∉ Pri. 100 ∉ Pri. 110 ∉ Pri. 1000 ∉ Pri. 1001 ∉ Pri. Etc
                  a \in \text{Pri.} \Leftrightarrow a = 10 \cdot \vee a = 11 \cdot \vee \cdot 101 \cdot \vee a = 111 \cdot \vee \cdot \text{Etc}
                  10 \neq a . \land . 10 \text{ Div } a \rightarrow . a \notin \text{Pri}:
                  11 \neq a \land \land 11 \text{ Div } a \rightarrow a \notin \text{ Pri}:
                  100 \neq a . \land . 100 \text{ Div } a \rightarrow .a \notin \text{Pri}:
                  Etc:
                  b \neq a . \land . b \neq 1 . \land . b \text{ Div } a \rightarrow . a \notin \text{Pri}
                  a \in \operatorname{Pri} \rightarrow b \operatorname{Div} a \rightarrow b = a \lor b = 1
                  a \in \operatorname{Pri}_{\bullet} \to b : b \operatorname{Div}_{\bullet} \to b = a \cdot \lor b = 1
                  a \in \text{Pri} \leftrightarrow a - 1 \in \text{Num} \land b : b \text{ Div } a \rightarrow b = a \land b = 1 
              The word written Pri means "prime numbers".
```

121 1. $\# a \in \text{Num} \rightarrow a \in \text{Int}$

Num CInt:

 $a \in \text{Pri.} \rightarrow .a \in \text{Num}$

Pri C Num:

 $a \in \text{Pri} \rightarrow a \in \text{Int}$

Pri CInt:

 $a \in \text{Num} \cdot \rightarrow \cdot a \in \text{Num} \cdot$

Num C Num:

 $a \in \text{Pri} \rightarrow a \in \text{Pri}$

Pri C Pri: Int CInt:

Num > Pri.

Int D Num.

Int $\supset Pri$

Pri C Num C Int.

Int D Num D Pri

Pri = Pri . Num = Num . Int = Int *

1 21 2. *
$$\land a : a \in A . \Leftrightarrow \cdot a = 11. \lor . a = 111. \lor . a = 10011 \to . A \subset Pri$$
:
$$\land a : a \in B . \Leftrightarrow \cdot a = 11. \lor . a = 111. \lor . a = 10011 \to . B \subset Pri$$
:
$$\land a : a \in A . \Leftrightarrow \cdot a = 11. \lor . a = 111. \lor . a = 10011 \to . A \subset Num$$
:
$$\land a : a \in A . \Leftrightarrow \cdot a = 11. \lor . a = 111. \lor . a = 10011 \to . A \subset Int$$
:
$$\land a : a \in A . \Leftrightarrow \cdot a = 100. \lor . a = 101. \lor . a = 10000 \to . A \subset Num$$
:
$$\land a : a \in A . \Leftrightarrow \cdot a = 100. \lor . a = 101. \lor . a = 10000 \to . A \subset Int$$
:
$$\land a : a \in A . \Leftrightarrow \cdot a = 100. \lor . a = 101. \lor . a = 0 \to . A \subset Int$$
and so on.

1213. *
$$\wedge a \cdot a \in A \cdot \rightarrow \cdot a \in \operatorname{Pri} : \leftrightarrow \cdot A \subset \operatorname{Pri} : \\ \wedge a \cdot a \in \operatorname{Pri} \cdot \rightarrow \cdot a \in A : \leftrightarrow \cdot \operatorname{Pri} \subset A : \\ \wedge a \cdot a \in \operatorname{Pri} \cdot \leftrightarrow \cdot a \in A : \leftrightarrow \cdot A = \operatorname{Pri} *$$
 and so on.

1 21 4.
$$*b \in \text{Num} . \rightarrow `\land a \cdot a \in A . \leftrightarrow . a \text{ Div } b : \rightarrow .A \subset \text{Int:}$$

$$b \in \text{Num} . \rightarrow `\land a \cdot a \in A . \leftrightarrow .b \text{ Div } a : \rightarrow .A \subset \text{Int:}$$

$$\land a : a \in A . \leftrightarrow `a \in \text{Int.} \land .a > 111 `\rightarrow .A \subset \text{Num} *$$
and so on.

1 21 6.
$$\# A \subset \text{Pri} \rightarrow A \subset \text{Num} \cdot \text{Etc} \#$$

1 22 1. #
$$?x \cdot x^{10} = 10 \cdot \land \cdot x > 0$$
:
 $x^{10} = 10 \cdot \land \cdot x > 0 \cdot \rightarrow \cdot x = \sqrt{10}$:
 $\sqrt{10} = 1,0110101000001$ Etc:
 $?x \cdot x^{10} = -1$:
 $x^{10} = -1 \cdot \leftrightarrow x = i, \forall \cdot x = -i$ #

1 22 2. *
$$\pi = 100 \times (1 - 1/11 + 1/101 - 1/111 + 1/1001 - \text{Etc})$$
. $e = 1 + 1/1 + 1/10 + 1/(10 \times 11) + 1/(10 \times 11 \times 100) + \text{Etc}$. $e^{10} = 1 + 10/1 + 100/10 + 1000/(10 \times 11) + 10000/(10 \times 11 \times 100) + \text{Etc}$ *

Of course i, e, π are constants with the usual meaning. In written Lincos we shall not try to distinguish such constants from variables by typographical characteristics. In Lincos speech they will be carefully distinguished by phonetic means.

1 22 3. # 0
$$\in$$
 Rat . 1 \in Rat . - 1 \in Rat . 10 \in Rat . Etc · $1/10 \in$ Rat . - $1/10 \in$ Rat . 101/11 \in Rat . - 1010/1011 \in Rat . Etc · $\sqrt{10} \notin$ Rat . $\pi \notin$ Rat . $e \notin$ Rat . $e^{10} \notin$ Rat · 111 $\sqrt{10} - 1$. \notin Rat · $\pi + e$. \notin Rat · 111 $\pi - 0$, 1 e . \notin Rat · 10 $\sqrt{10} - 0$, 01 $\pi + e^{10}$. \notin Rat · $i \notin$ Rat · $\pi + e^{10} \in$ Rat · 11 + 10 i . \notin Rat · Etc #

1 22 5.
$$*a \in \text{Rat} \cdot \land \cdot b \in \text{Rat} \cdot \rightarrow : a + b \cdot \in \text{Rat} \cdot \land \cdot a - b \cdot \in \text{Rat} \cdot \land \cdot a b \cdot \in \text{Rat}$$

$$a \in \text{Rat} \cdot \land \cdot b \in \text{Rat} \cdot \land \cdot b \neq 0 \cdot \rightarrow \cdot a/b \cdot \in \text{Rat} *$$

1 22 6.
$$\# x \in \text{Rat.} \land .x^{10} \in \text{Int.} \rightarrow .x \in \text{Int.}$$

 $x \in \text{Rat.} \land .x^{11} \in \text{Int.} \rightarrow .x \in \text{Int.}$

```
1 22 7. a \in \text{Int.} \rightarrow .a \in \text{Rat.}
Int \subset \text{Rat:}
a \in \text{Num.} \rightarrow a \in \text{Rat.}
Num \subset \text{Rat:}
Pri \subset \text{Rat:}
Rat \supset \text{Int.} \supset \text{Num.} \supset Pri:
A \subset \text{Int.} \rightarrow .A \subset \text{Rat.}
\land a \cdot a \in A . \leftrightarrow : \forall x \cdot x \in \text{Num.} \land .a = x/11. \rightarrow .A \subset \text{Rat.}
\land a \cdot a \in A . \leftrightarrow : \forall x \cdot x \in \text{Int.} \land .a = 1/x. \rightarrow .A \subset \text{Rat.}
A \subset \text{Rat.} \land . \land a \cdot a \in B. \leftrightarrow : \forall x \cdot x \in A . \land .x^{10} + x = a \cdot \rightarrow .B \subset \text{Rat.}
```

1 23 1. * $0 \in \text{Rea}: 1 \in \text{Rea}: -10. \in \text{Rea}: 111/11. \in \text{Rea}: \sqrt{10} \in \text{Rea}: \pi \in \text{Rea}: e \in \text{Rea}: e^{10} \in \text{Rea}: 111/10 - 101 + 0, 1\pi e. \in \text{Rea}: \text{Etc}: 110,101001110110001 \text{ Etc.} \in \text{Rea}: \text{Etc}: i \notin \text{Rea}: 111 - 11i. \notin \text{Rea}: \text{Etc}: a \in \text{Rea}. \land .b \in \text{Rea}. \land .b \neq 0 \cdot \rightarrow .a + bi. \notin \text{Rea}: a \in \text{Rea}. \land .b \in \text{Rea}. \rightarrow .a + bi. \in \text{Com}$

The word written 'Rea' (fL numerus realis) means "real numbers". The word written 'Com' means "complex numbers". Though arithmetics of real and complex numbers have not yet been dealt with, we may venture to give a list of examples in which sums and products of real numbers occur. In similar light-hearted fashion we write down infinite sums, though limits have not yet been defined. Every one of us knew real numbers and limits long before he learned an exact definition of these notions.

```
1 23 2. \# a \in \text{Rat}. \to .a \in \text{Rea}

Rat \subset \text{Rea}:

a \in \text{Rea}. \to .a \in \text{Com}

Rea \subset \text{Com}:

a \in \text{Pos}. \Leftrightarrow a \in \text{Rea}. \land .a > 0:

Pos \subset \text{Rea}:

Com \supset \text{Rea} \supset \text{Rat} \supset \text{Int} \supset \text{Num} \supset \text{Pri}:

A \subset \text{Rat}. \to .A \subset \text{Rea}:

\land a : a \in A. \Leftrightarrow a \in \text{Rea}. \land .a \notin \text{Rat} \to .A \subset \text{Rea} \#

The word written Pos means "positive real numbers".
```

1 23 3. * 0,1101010001 Etc +0,1011010010 Etc: =1,1000100 Etc: 1,00100111101 Etc +0,10011110010 Etc: =1,110001 Etc:

and so on.

A typical difficulty met with in computations with infinite fractions has been shown. This was quite unavoidable. In intuitionistic mathematics it is not true that every real number can be developed in an infinite fraction.

1 23 5. # $x \in \text{Rea.} \land .y \in \text{Rea.} \land .a \in \text{Rat.} \land .b \in \text{Rat.} \land .c \in \text{Rat.} \land .d \in \text{Rat$

1 23 6. *
$$x \in \text{Rea.} \land .y \in \text{Rea.} \rightarrow .x \leq y$$
.
 $\Leftrightarrow \land \alpha : \alpha \in \text{Rat.} \land .\alpha \leq x \rightarrow .\alpha \leq y$!
 $x \in \text{Rea.} \land .y \in \text{Rea.} \rightarrow .x = y . \Leftrightarrow .x \leq y \leq x$ *

By 1 23 4-1 23 6 real numbers as well as their sums and products (for non-negative factors) are formally defined. The definition, which is essentially Dedekind's, fits into intuitionistic mathematics also. It does not rely on infinite fractions.

The remaining portion of real arithmetics should be dealt with in an analogous way.

1 23 7. We do not dwell on complex numbers and their arithmetics as they do not cause any difficulty.

1 24 0. In accordance with our general policy sets and relations between sets (or classes) have preceded the appearance of the word "set". Of course we cannot formally tell what sets are. Even in axiomatics of set-theory we learn no more than how to expand a given stock of sets.

We shall teach the same thing, but we are bound to do it in a less formal way. The working of axiomatic systems has always been explained within highly developed meta-languages, whereas for the present we do not know any Lincos word that we could use if we were tempted to speak about the mathematics we are showing.

Set theory as such is not very important in our program, but as a tool it is indispensable. We have tackled this subject in traditional fashion. One point that should be mentioned, is our adoption of Quine's " $a \in a$ " for one-element-sets a.

In the Introduction we explained why the design of a language like Lincos should be so liberal as not to impose opinions and convictions by means of linguistic devices. Our attitude towards the foundations of set theory and of mathematics has been largely determined by this policy. No solution may seriously be excluded by peculiarities of Lincos. As a consequence we cannot yet tackle the set-theory paradoxes nor try to devise means of avoiding them. At a more advanced stage of Lincos we can start dialogues between representatives of different thoughts, in order to give an impartial account of the various solutions. Nevertheless we cannot avoid taking a few decisions even at this early stage. We are still free to abrogate them as soon as it becomes desirable.

We have adopted Quine's $a \in a$ for certain one-element-sets a, but this is a mere convention without serious consequences. We

have adopted a name (Agg) for "sets" (classes), though only after much hesitation. A name is useless if it is not subjected to syntactical rules, which may anticipate certain solutions and exclude other possibilities. The question may be raised whether the adherence of Num to Agg is of the same kind as that of 3 to Num, or whether it is so different as to justify the use of a new copula instead of '\(\in \)'. We have decided to write Num \(\in \) Agg provisionally.

Introducing 'Agg' in this way commits us to indicating a few entities not belonging to Agg, if we are to stick to our general principles. We could state $A \notin Agg$ for some sets A which give rise to the well-known set-theory paradoxes (e.g. for Agg itself). It should be noted that there is not much sense in this kind of statement unless it serves to kill the paradoxes. As far as I can see adopting $A \notin Agg$ for such sets A is as dangerous as admitting $A \in Agg$, if the paradoxes are not prevented by other means. If we wish to exclude $A \in Agg$, we should do it by syntactical means, i.e. by asserting that it is void of sense. At present means of this kind are not yet available.

We therefore decided to indicate predicates (relations) and connectives like 'Div' and ' \land ' as things not belonging to Agg. I believe this is a honest solution because the average mathematician, when asked for instances of non-sets, is likely to answer in the same sense. I admit it is rather unsatisfactory because it implies that expressions like $\text{Div} \in \text{Agg}$ are meaningful (though false). Of course no meaning may be attached to $x \in \text{Div}$ and so on; otherwise Div would be equal to the void set.

Some expressions void of sense will be mentioned in 1 36 3.

1 24 1. * Num
$$\in$$
 Agg . Int \in Agg . Pri \in Agg . Rat \in Agg . Com \in Agg : $101 \in \text{Agg} \cdot -10 \cdot \in \text{Agg} \cdot 11/111 \cdot \in \text{Agg} \cdot \sqrt{10} \in \text{Agg} \cdot i \in \text{Agg} \cdot \text{Etc} :$ Div \notin Agg . $+ \notin \text{Agg} \cdot > \notin \text{Agg} \cdot \land \notin \text{Agg} \cdot \text{Etc} \notin \text{Agg} \cdot \text{Etc} =$

The word written Agg (fL aggregatum) means "sets" ("classes").

1 24 2. *
$$\land a : a \in A . \Leftrightarrow `a \in \text{Com}. \land : \lor b : b \in \text{Num}. \land .a = 11 + bi : \\ \rightarrow .A \in \text{Agg I}$$

$$\land x ` x \in B . \Leftrightarrow : x \in \text{Rat}. \land : 111 x^{10} - 10 x + 0, 1 . \in \text{Num}: \\ \rightarrow .B \in \text{Agg I}$$

$$\land a ` a \in C . \Leftrightarrow : a \in \text{Pri}. \land : a + 10 . \in \text{Pri}: \rightarrow .C \in \text{Agg I}$$

$$\land n : n \in N . \Leftrightarrow : n \in \text{Num}. \land : \lor x ` \lor y : \lor z .$$

$$x \in \text{Num}. \land .y \in \text{Num}. \land .z \in \text{Num}.$$

$$\land .x^{n+10} + y^{n+10} = z^{n+10} : \rightarrow .N \in \text{Agg } \#$$

$$a \in \text{Com} \rightarrow x \in a \leftrightarrow x = a$$
:

 $a \in A \longrightarrow a \in Agg *$

1 24 4. * Num ≠ Num . ∧ . Num = Num · Etc:

Pri∉Num.Com∉Rat.Int∉Rea.Etc:

$$\land a: a \in A . \Leftrightarrow a = 11. \lor .a = 101$$

$$\land$$
 $\land a : a \in B . \Leftrightarrow \cdot a = 11 . \lor . a = 101 . \lor . a = 0$:

 \rightarrow 11 \in B. \land 101 \in B. \land . $A \notin B$. \land . $B \notin B$ *

1 24 5.
$$* \land a:a \in A \leftrightarrow a=0, \lor, a=1$$

$$\land$$
 $\land a : a \in B . \Leftrightarrow .a = 0 : \land$ $\land x : x \in C . \Leftrightarrow x = A . \lor .x = B$.

$$\rightarrow$$
 $A \neq B$. \land $B \neq C$. \land $A \neq C$. \land $A \notin A$.

$$\land .A \notin B . \land .A \in C . \land .B \in A . \land .B \in B . \land .B \in C$$
.

 $\land . C \notin A . \land . C \notin B . \land . C \notin C *$

1246.
$$* \land x : x \in C$$
. $\Leftrightarrow x = \text{Pri.v.} x = \text{Num.v.} x = \text{Int}$

 \rightarrow · Pri \in C · \land · Num \in C · \land · Int \in C · \land · C \neq Pri · \land · C \neq Num ·

 $\land . C \neq Int. \land . 1010 \notin C *$

1 24 7. $* \lor C \cdot C \in Agg . \land .Pri \in C$:

$$\land x \cdot x \in C \iff x = \text{Pri} : \Rightarrow C \in \text{Agg}$$

$$\land x \cdot x \in C . \Leftrightarrow .x = \text{Pri} : \Leftrightarrow .C = \lceil \text{Pri} \rceil$$

 $A \in Agg \rightarrow \land x \cdot x \in \lceil A \rceil \leftrightarrow x = A : \land \land \lceil A \rceil \in Agg$

 $x \in \lceil x \rceil$:

$$y \in \lceil x \rceil . \Leftrightarrow .x = y *$$

The connective a transforms a set into a one-element-set containing just the given set as an element.

1 25 1. *
$$A \subset B . \Leftrightarrow : \land x \cdot x \in A . \Rightarrow . x \in B$$

$$A \subset B \leftrightarrow B \supset A$$

$$A \subset B$$
, \land , $B \subset C$ \rightarrow , $A \subset C$

$$A = B \cdot \Leftrightarrow : \land x \cdot x \in A \cdot \Leftrightarrow \cdot x \in B$$

$$A = B \cdot \Leftrightarrow A \subset B \cdot \land B \subset A #$$

1 25 2. $* \land a : a \in A . \Leftrightarrow a \in Pri . \land . 10 Div a \rightarrow . A = 10$

$$\land a : a \in A . \Leftrightarrow `a \in \mathsf{Rat}. \land .a < / \overline{10} ` \land `$$

$$\land a: a \in B. \Leftrightarrow a \in \text{Rat}. \land a \leq \sqrt{10}: \Rightarrow A = B:$$

$$\land a : a \in A . \Leftrightarrow a \in \text{Rea.} \land . -1 < a < 1 \land .$$

$$\land a$$
 ' $a \in B$. \leftrightarrow : $\lor x$ ' $x \in A$. \land . $x^{11} = a$ ' \rightarrow . $A = B$ *

1 25 3. $A \subset \text{Num} \land 1 \in A \land \land \land x : x \in A \rightarrow x + 1 \in A \rightarrow A = \text{Num}$

$$A \subset \text{Int.} \land .1 \in A \cdot \land . \land x \cdot x \in A . \rightarrow : x + 1 . \in A \cdot \land \cdot x - 1 . \in A$$

$$\rightarrow A = Int *$$

1 25 4.
$$* \lceil x \rceil = \lceil y \rceil . \leftrightarrow . x = y$$
:
 $\lceil x \rceil = x . \leftrightarrow : \land y \cdot y \in x . \leftrightarrow . y = x *$

1 26 1.
$$* \land a : a \in A . \Leftrightarrow `a \in \text{Rea.} \land .a \ge 11$$
'
$$\land ` \land a : a \in B . \Leftrightarrow a \in \text{Rea.} \land .a < 101:$$

$$\Rightarrow : A \cup B . = \text{Rea.} \land ` \land a : a \in .A \cap B ` \Leftrightarrow .11 \le a < 101:$$

$$\land ` \land a : a \in A \setminus B ` \Leftrightarrow .a \ge 101 *$$

1 26 2.
$$*x \in A \cap B \cdot \Leftrightarrow `x \in A \land .x \in B `$$
 $x \in A \cup B \cdot \Leftrightarrow `x \in A \lor .x \in B `$
 $x \in A \setminus B \cdot \Leftrightarrow `x \in A \land .x \notin B `$
 $A \in Agg \land .B \in Agg \cdot \Rightarrow :A \cap B \cdot \in Agg \cdot \land .A \cup B \cdot \in Agg \cdot \land .A \setminus B \cdot \in Agg \cdot \Leftrightarrow .A \cap B \cdot \in Agg \cdot Agg \cdot \Leftrightarrow .A \cap B \cdot \in Agg \cdot Agg \cdot$

1 26 3. *
$$\land a: a \in A. \Leftrightarrow `a = 10. \lor. a = 101. \lor. a = -1`$$

 $\Leftrightarrow : A = `10 \cup 101 \cup .-1`$
 $A = .1 \cup 10 \cup 101 \cdot \land `B = .1 \cup 0 \cup 111:$
 $\Rightarrow : A \cup B. = .0 \cup 1 \cup 10 \cup 101 \cup 111 \cdot$
 $\land `A \cap B. = 1`$
 $\land `A \setminus B. = .10 \cup 101 *$

1 26 4.
$$*A \setminus A = \lceil \rceil *$$
The void set.
 $*x \neq y . \Leftrightarrow \lceil x \rceil \cap \lceil y \rceil . = \lceil \rceil *$

The word written Car means "cardinal number of . . . ".

1 27 2.
$$\# \lor m \cdot m \in \text{Int.} \land . m \ge 0 \cdot \land . \text{Car } A = m : \land . x \notin A$$

$$\rightarrow : \text{Car } A \cup \ulcorner x \urcorner \cdot = \cdot \text{Car } A \cdot + 1 :$$

$$\lor m \cdot m \in \text{Int.} \land . m \ge 0 \cdot \land . \text{Car } A = m :$$

$$\land : \lor n \cdot n \in \text{Int.} \land . n \ge 0 \cdot \land . \text{Car } B = n :$$

$$\land \cdot A \cap B \cdot = \ulcorner \urcorner \urcorner$$

$$\rightarrow : \text{Car } A \cup B \cdot = \cdot \text{Car } A \cdot + . \text{Car } B :$$

$$\lor m \cdot m \in \text{Int.} \land . m \ge 0 \cdot \land . \text{Car } A = m : \land . B \subset A \urcorner$$

$$\rightarrow : \text{Car } B \in \text{Int.} \land \cdot 0 \le . \text{Car } A = m :$$

$$\land : \lor m \cdot m \in \text{Int.} \land . m \ge 0 \cdot \land . \text{Car } A = m :$$

$$\land : \lor n \cdot n \in \text{Int.} \land . n \ge 0 \cdot \land . \text{Car } B = n \urcorner$$

$$\rightarrow \because \text{Car } A \cup B \cdot + \cdot \text{Car } A \cap B : = \cdot \text{Car } A \cdot + . \text{Car } B \#$$

1 27 3. *
$$n+1$$
. \in Num $\cdot \wedge \cdot \wedge x : x \in A$. $\leftrightarrow \cdot x \in$ Num $\cdot \wedge \cdot 1 \le x \le n$:
 $\rightarrow \cdot \operatorname{Car} A = n$:
 $n+1$. \in Num $\cdot \wedge \cdot \wedge x : x \in A$. $\leftrightarrow \cdot x \in$ Num $\cdot \wedge \cdot 111 \le x \le n+110$:
 $\rightarrow \cdot \operatorname{Car} A = n$:
 $n+1$. \in Num $\cdot \wedge \cdot \wedge x : x \in A$. $\leftrightarrow : \forall y \cdot y \in \operatorname{Num} \cdot \wedge \cdot 1 \le y \le n$. $\wedge \cdot x = y^{10}$:
 $\rightarrow \cdot \operatorname{Car} A = n$ *

- 1 27 4. * Car Pri = Car Num = Car Int = Car Rat = $\chi_0 \neq$ Car Rea = Car Com *
- 1 27 5. # Car ∉ Agg #
- 1 28 1. $* Car A = 11. \land . Car B = 10 \rightarrow . Car [A.B] = 110 *$

Pairs of sets, especially of one-element-sets.

「¬ and 「¬ must not be confused.

1 28 3.
$$\# \mbox{ } \mbox{$m:m+1. \in \mbox{Num} \cdot \land . \mbox{ } \mbox{Car } A=m$}$$

$$\mbox{$\wedge \cdot \lor n:n+1. \in \mbox{Num} \cdot \land . \mbox{ } \mbox{Car } B=n$}.$$

$$\mbox{$\to \cdot \mbox{ } \mbox{Car } \mbox{$\cap A$}. \mbox{B} \mbox{$\cap A$}. \mbox{$\cap B$} \mbox{$\cap A$}. \mbox{$\cap A$}. \mbox{$\cap A$}.$$

- 1 28 5. $*a \in Agg . \land .b \in Agg \cdot \rightarrow . \lceil a .b \rceil \in Agg :$ $x \in \lceil a .b \rceil . \leftrightarrow .x = \lceil a .b \rceil :$ $\lceil a .b \rceil = \lceil c .d \rceil . \leftrightarrow .a = b . \land .c = d *$

1 28 7. *
$$a \in Agg. \land .b \in Agg. \land .c \in Agg. \rightarrow . \ulcorner a.b.c \urcorner \in Agg. \mid x \in \ulcorner a.b.c \urcorner . \leftrightarrow .x = \ulcorner a.b.c \urcorner \mid$$

$$\lceil a.b.c \urcorner = \lceil u.v.w \urcorner . \leftrightarrow \cdot a = u. \land .b = v. \land .c = w \mid$$

$$A \in Agg. \land .B \in Agg. \land .C \in Agg. \rightarrow . \ulcorner A.B.C \urcorner \in Agg. \mid \land .Ax : x \in \ulcorner A.B.C \urcorner .$$

$$\land : \lor a : \lor b : \lor c \cdot a \in A. \land .b \in B. \land .c \in C. \land .x = \ulcorner a.b.c \urcorner \mid$$

$$1 + Car A . \in Num \cdot \land \cdot 1 + Car B. \in Num \cdot \land \cdot 1 + Car C. \in Num :$$

$$\rightarrow . Car \lceil A.B.C \rceil = Car A \times Car B \times Car C *$$

Triples of sets, especially of one-element-sets.

1 29 0. Notations may be simplified:

1 29 2. *
$$x = a \cdot \lor \cdot x = b \cdot \lor \cdot x = c \cdot \lor \cdot x = d \cdot \leftrightarrow \cdot x = .a \checkmark b \checkmark c \checkmark d$$
:
 $a \in X \cdot \land .b \in X \cdot \land .c \in X \cdot \land .d \in X \cdot \leftrightarrow \cdot a \land b \land c \land d \cdot \in X$:
 $a \in X \cdot \lor .b \in X \cdot \lor .c \in X \cdot \lor .d \in X \cdot \leftrightarrow \cdot a \checkmark b \checkmark c \checkmark d \cdot \in X$ *

This is to legalize a familiar use of "and" and "or" in common mathematical and natural language.

A set-forming connective, written ^ has been introduced. In an earlier version I put a circumflex (a small quantifier of generalization) over the letter representing the variable to be bound by set-forming. Typographical considerations led me to put the circumflex over a vertical stroke before the letter.

1 29 4.
$$\# a \in : ^\uparrow x \cdot \text{Etc.} x \cdot \text{Etc.} * x \cdot \text{Etc.} a \cdot \text{Etc.}$$

$$A \in \text{Agg.} \to : ^\uparrow x : x \in A \cdot \wedge \cdot \text{Etc.} x \cdot \text{Etc.} * x \cdot \text{Etc.}$$

1 29 5. *
$$1 + n \cdot \in \text{Num}$$
: $\rightarrow : \text{Car}$: $^{\uparrow}x \cdot x \in \text{Num}$. $\land . 1 \leq x \leq n = n$:

 $^{\uparrow}x \cdot x \in \text{Rat}$. $\land . x < \sqrt{10}$: $= ^{\uparrow}x \cdot x \in \text{Rat}$. $\land . x \leq \sqrt{10}$:

 $^{\uparrow}x \cdot x \in \text{Rea}$. $\land . x < 101$: $\cap : ^{\uparrow}x \cdot x \in \text{Rea}$. $\land . x > 11$:

 $= : ^{\uparrow}x \cdot x \in \text{Rea}$. $\land . 11 < x < 101$ *

70

```
~.1∪11.~.10∪11.~A:
                  \operatorname{Car} A = 11. \rightarrow \operatorname{Car} A := 1000:
                  A \in Agg. \land l + Car A. \in Num:
                      \rightarrow: ^{\uparrow}B. B \subset A := 10^{\operatorname{Car} A}:
                  ^B.BC「¬•≠「¬•
                  ^{\uparrow}B.BC^{\Gamma}=^{\Gamma}^{\Gamma}=^{\Pi}
1 30 2. # ^{+}x, 10 Div x: \cap ^{+}x, 11 Div x:
                      \cap \cdot \uparrow x, 100 Div x \cdot \cap \cdot \uparrow x, 101 Div x \cdot \cap \cdot Etc: = \lceil \rceil:
                  ^{\uparrow}x \cdot x \in \text{Rat}, \land x \leq 0.1 : \cap ; ^{\uparrow}x \cdot x \in \text{Rat}, \land x \leq 0.01 :
                      \cap : ^{\uparrow}x \cdot x \in \text{Rat.} \land : x \leq 0.001 : \cap : \text{Etc}' = "^{\uparrow}x \cdot x \in \text{Rat.} \land : x \leq 0 :
                  \land X \cdot X \in M \leftrightarrow X = A \leadsto B \leadsto C \leadsto \text{Etc}:
                      \rightarrow \cap M = A \cap B \cap C \cap \text{Etc}:
                  M \in Agg. \rightarrow \cdot \cap M. \in Agg: 
                      \land \land \land x `x \in . \cap M `\Leftrightarrow : \land A `A \in M . \Rightarrow .x \in A *
1 30 3. * \land A : A \in M . \Leftrightarrow : \forall n : n \in \text{Num} . \land : A = : \uparrow x . n \text{ Div } x :
                      \rightarrow: \bigcap M := ^{\uparrow}x: \land n \cdot n \in \text{Num} : \rightarrow \cdot n \text{ Div } x = \lceil \rceil \rfloor
                  \bigcap : ^{\uparrow}y : \forall n : n \in \text{Num} . \land : y = ^{\uparrow}x . n \text{ Div } x : = ^{\uparrow} \exists 1
                  c \in \text{Rea.} \rightarrow :: \bigcap : ^{\uparrow}y : \forall a : a \in \text{Rea.} \land . a < c.
                      \land : y' = : \uparrow x : x \in \text{Rea.} \land : x > a : : = : \uparrow x : x \in \text{Rea.} \land : x \ge c \text{ } 
                  A \subset B \rightarrow ^{\bullet} \cap ^{\bullet} C \cdot A \subset C \subset B := A *
                  ^{\uparrow}x \cdot x \in \text{Rea.} \land .x \leq 1 : \cup : ^{\uparrow}x \cdot x \in \text{Rea.} \land .x \leq 10:
1 30 4.
                      \cup: ^{\uparrow}x \cdot x \in \text{Rea.} \land .x \leq 11: \cup: \text{Etc} = \text{Rea}
                  \land X \cdot X \in M . \Leftrightarrow .X = A \sim B \sim C \sim \text{Etc}:
                      \rightarrow \cdot \bigcup M = A \cup B \cup C \cup \text{Etc}:
                  M \in Agg \rightarrow \bigcup M \in Agg:
                      \land : \land x : x \in A : A \in M . \land : x \in A^*
\wedge y = 1^n \operatorname{Car} x \cdot x \operatorname{Div} n := m = Num I
                  B \cdot A := A *
1 31 1. * x \cdot 11x = 10 \cdot = 10/11:
                  x \cdot x^{10} = 1001, x \cdot x > 0 := 11
                  \forall x \cdot x \in \text{Num}. \land 0 \leq \sqrt{10} - x < 1 := 1
                  x : x \text{ Div } 1100 . \land . x \text{ Div } 10010 .
                      \land ' \land y : y Div 1100. \land . y Div 10010 \rightarrow . y Div x := 110!
                  ^{\vee}x \cdot x^{10} = 1001 \cdot \land \cdot 0 \le x : \text{Div } 1001 :
                  a \in \text{Num} . \land . a > 1 . \land . a \notin \text{Pri} . \land `A = : `y \cdot y \text{ Div } a . \land . 1 < y :
                  \rightarrow : \forall x : x \in A . \land : \land y : y \in A . \rightarrow . x \leq y : \leq \forall a :
```

```
a \in A \cdot \wedge \cdot \operatorname{Car} A = 1 \cdot \rightarrow : ^{ } x \cdot x \in A \cdot = a :
^{ } x \cdot x \in A \cdot = a : \leftrightarrow \cdot A = ^{ } a ^{ } :
^{ } x \cdot \operatorname{Ete} x \operatorname{Ete} : = a : \leftrightarrow : \operatorname{Ete} a \operatorname{Ete} \cdot \wedge \cdot \operatorname{Car} \cdot ^{ } x \cdot \operatorname{Ete} x \operatorname{Ete} : = 1 :
\operatorname{Ete} \cdot ^{ } x \cdot x \in A \cdot \operatorname{Ete} : \leftrightarrow ^{ } \operatorname{Car} A = 1 \cdot \wedge : \wedge a \cdot a \in A \cdot \rightarrow \cdot \operatorname{Ete} a \operatorname{Ete} \overset{+}{}
```

The Lincos connective written as a "bird" on a vertical stroke, $\dot{}$, has been used as definite article. (Originally the bird was put over the letter representing the variable to be bound. See 1 29 3.)

These examples are insufficient. The factual use of the article can be learnt only from non-mathematical contexts. Our general ideas about description and the use of the article have been set forth in 0 25. When regularly used, ${}^{\checkmark}x \cdot x \in A$ cannot occur unless $\operatorname{Car} A = 1$. But this does not mean that ${}^{\checkmark}x \cdot x \in A$ occurs only if $\operatorname{Car} A = 1$. Irregular between occurences of ${}^{\checkmark}x \cdot x \in A$ are needed if only as examples of irregular occurence.

An indefinite article will be introduced later.

As a consequence of our use of the definite article, we cannot decide in a mechanical way or on purely syntactic grounds whether certain expressions are meaningful or not. But this is no disadvantage. Lincos has been designed for the purpose of being used by people who know what they say, and who endeavour to utter meaningful speech.

```
1 32 1. # \land x \cdot x \in \text{Rea.} \rightarrow (x = x^{10} + 11x + 1):
                        \rightarrow: t0 = 1. \land .t1 = 101. \land .t. - 1 = -1. \land .t0, 1 = 10, 11.
                        \wedge \cdot t \cdot -0.1 = -0.01 \cdot \wedge \cdot t \cdot 0.01 = 1.1101 \cdot \wedge \cdot t -0.01 = 0.0101.
                        \land x \cdot x \in \text{Rea.} \land a \leq x \leq b \cdot
                        \rightarrow: \forall y \cdot c \leq y \leq d. \land y = (101x - 1)/(x^{10} + 11x + 1):
                        \rightarrow: \forall x . (101 x - 1)/(x^{10} + 11 x + 1) \land x . a \leq x \leq b.
                             \in : ^{\uparrow}x \cdot a \leq x \leq b \cdot \curvearrowright \text{Rea}
                        \land : \land z \cdot a \leq z \leq b \cdot \land z \in \text{Rea}
                             \rightarrow: \forall x . (101x-1)/(x^{10}+11x+1) \uparrow x . a \le x \le b \cdot z
                              =(101z-1)/(z^{10}+11z+1)ı
                    \land x : x \in \text{Rea.} \land a \leq x \leq b
                         \rightarrow Car: ^{\uparrow}y \cdot c \leq y \leq d \cdot \land \cdot y^{10} = x^{10} + 11x + 1' = 1
                             \rightarrow: \forall t: t \in : ^{\uparrow}x. a \leq x \leq b. \frown ^{\downarrow}y. c \leq x \leq d.
                         \wedge: \land x \cdot a \leq x \leq b.
                         \rightarrow . /x = \sqrt{x^{10} + 11 x + 1} \sim -\sqrt{x^{10} + 11 x + 1}:
                         \land \vdots f = ``\forall x : `\forall y \cdot c \leq y \leq d \cdot \land \cdot y^{10} = x^{10} + 11 x + 1 \ | \uparrow x \cdot a \leq x \leq b 
                    f \in A \curvearrowright B \Leftrightarrow f \subset \lceil A \cdot B \rceil.
                         \wedge : \wedge x : x \in A \rightarrow `Car \cdot \uparrow y . \exists x . y \exists \in f := 1
                    f \in A \curvearrowright B \rightarrow x \in A \rightarrow fx = y \rightarrow x \cdot x \cdot y \in f
```

With the set-theory means now available we are allowed to consider functions as special sets, especially a function from A to B as a certain subset of $\lceil A \cdot B \rceil$ (a graphic). Yet this would not be the point to start from. When in an earlier period of our mathematical life we learned the notion of function and became familiar with it, a function was, as it were, a kinematical law, by which to every value of a variable, say x, corresponds one value of another variable, say y. Whenever we use functions, this idea predominates, even if we know some more formal definition. Psychologically and didactically the definition of a function as a special set would be inadequate. Nobody would grasp the sense and the importance of the concept of function. It is even unlikely that users of this concept will recognize any well-known idea in the formalized definition of function, if there is nobody to tell them what is really meant.

Perhaps it would have been better to delay the introduction of the concept of function, and to deal with it in the Chapter "Behaviour" (say after the notion "law"); for our purpose didactic arguments should be more decisive than systematic ones. Nevertheless I have tried to introduce the notion of function at this point. The exact definition has been prepared by a number of examples that might be multiplied ad lib. The examples are rather clumsy, but this can hardly be avoided. For instance a text such as

$$\land x : x \in \text{Rea.} \rightarrow : \text{Car.} \land y : y = x^{10} + 11 x + 1 := 1$$

 $\land . fx = x^{10} + 11 x + 1 := . f \in \text{Rea.} \land \text{Rea}$

would be of no use. Indeed, it has been our principle to avoid implications with an always true antecedent.

The class of functions from A to B is written $A \cap B$ instead of B^A , in order to avoid punctuations in exponents, and ambiguities if A and B are numbers. The "cup" upon the stroke is essentially Church's " λ ". (Its typographical history is analogous to those related in 1 29 3 and 1 31 1.) It transforms an expression into a function of the "cupped" variable. The "' is not a punctuation. It is a reminder of the vertical stroke as used in some versions of common mathematical language; after this sign the domain of the intended function is mentioned.

1 32 2.
$$*f \in A \curvearrowright \star \hookrightarrow : \lor B \cdot f \in A \curvearrowright B$$
'
 $f \in \star \star \curvearrowright B \cdot \hookrightarrow : \lor A \cdot f \in A \curvearrowright B$ '
 $f \in \star \star \curvearrowright \star \hookrightarrow : \lor \vdash A \cdot B \vdash f \in A \curvearrowright B = A \curvearrowright B$

1 32 4. *
$$\forall n : \forall y : y = 0 \sim 1 . \land : \forall x : x \in \text{Int} . \land . n = 2x + y : \text{Int} : = f :$$

$$\rightarrow f \in . \text{Int} \sim \text{Int} \land : f \in . \text{Int} \sim . 0 \cup 1 : \land \land \land f = . 0 \cup 1 :$$

$$\forall x : \forall y : y \in \text{Int} . \land . y \le x < y + 1 : \text{Rea} : \in . \text{Rea} \sim \text{Int} :$$

$$\forall f : x : y : x + y : f : \text{Rea} . \text{Rea} : = f : \rightarrow f \in . \text{Frea} . \text{Rea} : \text{Rea} :$$

$$\land : \land f : x : y : x \land y \in \text{Rea} . \rightarrow . f : f : x : y : x + y :$$

$$\forall x : f : x : y : x \land y \in \text{Rea} . \text{Rea} : \text{Rea} : \text{Rea} : \text{Rea} : \text{Rea} :$$

$$\forall x : f : x : y : x \land y \in \text{Rea} . \text{Rea} : \text{Rea} :$$

1 32 5. *
$$\omega_1 \lceil x \cdot y \rceil = x \cdot \omega_2 \lceil x \cdot y \rceil = y$$
:

$$\omega_1 \lceil x \cdot y \cdot z \rceil = x \cdot \omega_2 \lceil x \cdot y \cdot z \rceil = y \cdot \omega_3 \lceil x \cdot y \cdot z \rceil = z$$
:
Etc *

For the convenience of the terrestrial reader we use decadic numbers as indices. The words written ω_1 , ω_2 , ω_3 , and so on mean certain "universal" constants.

1 32 6. *
$$f = \cdot \forall x . x + 1$$
 Rea: $\land : g = \cdot \forall x . x^{10}$ Rea:
 $\rightarrow fg = \cdot \forall x . x^{10} + 1$ Rea: $\land : gf = \cdot \forall x . (x+1)^{10}$ Rea:
 $f \in A \curvearrowright B \cdot \land \cdot g \in B \curvearrowright C$:
 $\rightarrow gf \in A \curvearrowright C \cdot \land : \land x \cdot x \in A . \Rightarrow gf \cdot x \cdot g \cdot fx$:
 $f = \cdot \forall x . 1/(x+1)$ Rea\Rat: $\land : g = \cdot \forall x . (1/x) - 1$ Rea\Rat:
 $\rightarrow fg = \cdot \forall x . x$ Rea\Rat: $\land gf = \forall x . x$ Rea\Rat.

1 32 8.
$$\# 1 + \operatorname{Car} A :\in \operatorname{Num} : \wedge : 1 + \operatorname{Car} B :\in \operatorname{Num} :$$

 $\Rightarrow : \operatorname{Car} A \curvearrowright B := \operatorname{Car} B^{\operatorname{Car}} A \#$

1 32 9. #
$$f = {}^{\checkmark}x \cdot x^{10}$$
 | Com: $\wedge : g = {}^{\checkmark}x \cdot x^{10}$ | Rea ${}^{\checkmark} \rightarrow : g = f$ | Rea ${}^{\checkmark}f \in A \curvearrowright C \cdot \wedge : B \subseteq A$:

 $\Rightarrow {}^{\checkmark}g = f$ | $B \cdot \Leftrightarrow {}^{\checkmark}g \in : B \curvearrowright C \cdot \wedge : g = {}^{\checkmark}x \cdot fx$ | $B : B \subseteq A \cdot \Rightarrow : {}^{\checkmark}x \cdot \text{Etc } x \cdot \text{Etc } A \cdot B : = {}^{\checkmark}x \cdot \text{Etc } x \cdot \text{Etc } B = {}^{*}$

1 33 I. *
$$A = \cdot \uparrow a \cdot a > -11 : \rightarrow \inf A = -11 :$$

$$A = \cdot \uparrow a \cdot a \ge -11 : \rightarrow \inf A = -11 :$$

$$A = : \uparrow a \cdot a \in \operatorname{Rat} \cdot \land \cdot a^{10} < 10 \xrightarrow{\bullet} \inf A = -\sqrt{10} :$$

$$A \subset \operatorname{Rea} \cdot \land \cdot B \subset \operatorname{Rea} \cdot \land \cdot A \ne \lceil \neg \cdot \land \cdot B \ne \lceil \neg \cdot \rceil$$

$$\rightarrow : A \le B \cdot \leftrightarrow \lceil x \cdot y \rceil : x \in A \cdot \land \cdot y \in B \xrightarrow{\bullet} \cdot x \le y :$$

$$A \subset \operatorname{Rea} \cdot \land : \uparrow b \cdot A \ge b \xrightarrow{\bullet} \lceil \neg \rceil$$

$$\Rightarrow \inf A \in \operatorname{Rea} \cdot \land : A \ge \inf A \ge \cdot \uparrow b \cdot A \ge b \xrightarrow{*}$$

In an analogous way $\sup A$ will be defined.

The word written Agg Ord (fL aggregatum ordinatum) means "ordered sets". The word written Agg $\frac{1}{2}$ Ord means "partially ordered sets".

1 34 2. *
$$\lceil A \cdot B \rceil \in \operatorname{Agg} \operatorname{Ord} \operatorname{Ded} \cdot \Leftrightarrow \cdot \lceil A \cdot B \rceil \in \operatorname{Agg} \operatorname{Ord} \cdot \land \cdot C : C \subseteq A \cdot \land \cdot C \neq A \cdot \land \cdot C \neq \Gamma \rceil \cdot \land \land \land \land x \cdot y \rceil : y \in C \cdot \land \cdot \lceil x \cdot y \rceil \in B \cdot \Rightarrow \cdot x \in C : \Rightarrow : \forall c \cdot C \setminus c \cdot = \uparrow x \cdot \lceil x \cdot c \rceil \in B \mid B = : \uparrow z \cdot z = \lceil x \cdot y \rceil \cdot \land \cdot x \land y \in \operatorname{Rea} \cdot \land \cdot x < y \cdot \Rightarrow \cdot \lceil \operatorname{Rea} \cdot B \rceil \in \operatorname{Agg} \operatorname{Ord} \operatorname{Ded} *$$

The definition of Dedekind-ordered sets.

The word written Gru means "groups".

1 34 4. $\# \lceil A \cdot f \rceil \in \text{Gru Abe.} \Leftrightarrow : \lceil A \cdot f \rceil \in \text{Gru.} \land \lceil A \cdot y \rceil : \lceil x \cdot y \rceil \in \lceil A \cdot A \rceil.$ $\Rightarrow f \lceil x \cdot y \rceil := f \lceil y \cdot x \rceil #$

Abelian groups.

The word written Cam (fL "campus") means "fields".

1 34 6. *
$$\lceil A \cdot B \cdot f \cdot g \rceil \in \operatorname{Cam} \operatorname{Ord}$$
.
 $\Leftrightarrow : \lceil A \cdot f \cdot g \rceil \in \operatorname{Cam} \cdot \wedge \cdot \lceil A \cdot B \rceil \in \operatorname{Agg} \operatorname{Ord}$.
 $\land : \land : \lceil x \cdot y \cdot z \cdot u \cdot a \rceil : \lceil x \cdot y \rceil \in B \cdot \wedge \cdot \lceil z \cdot u \rceil \in \lceil A \cdot A \rceil \cdot \wedge \cdot f \lceil a \cdot a \rceil \cdot = a \cdot \wedge \cdot \lceil a \cdot u \rceil \in B$:
 $\Rightarrow \cdot \lceil f \lceil x \cdot z \rceil \cdot f \lceil y \cdot z \rceil \rceil \in B \cdot \wedge \cdot \lceil g \lceil x \cdot u \rceil \cdot g \lceil y \cdot u \rceil \rceil \in B$.
 $\lceil A \cdot B \cdot f \cdot g \rceil \in \operatorname{Cam} \operatorname{Ord} \operatorname{Ded}$.
 $\Leftrightarrow \cdot \lceil A \cdot B \cdot f \cdot g \rceil \in \operatorname{Cam} \operatorname{Ord} \cdot \wedge \cdot \lceil A \cdot B \rceil \in \operatorname{Agg} \operatorname{Ord} \operatorname{Ded}$.
 $B = : \uparrow z \cdot z = \lceil x \cdot y \rceil \cdot \wedge \cdot x \Leftrightarrow y \in \operatorname{Rea} \cdot \wedge \cdot x < y$.
 $\land : \land \lceil x \cdot y \rceil : \lceil x \cdot y \rceil \in \lceil \operatorname{Rea} \cdot \operatorname{Rea} \rceil$.
 $\Rightarrow : f \lceil x \cdot y \rceil : = x + y \cdot \wedge \cdot g \lceil x \cdot y \rceil : = xy$:
 $\Rightarrow \cdot \lceil \operatorname{Rea} \cdot B \cdot f \cdot g \rceil \in \operatorname{Cam} \operatorname{Ord} \operatorname{Ded}$

1 35 1. * Car $\in \star \sim \star$. $\sup \in \star \sim \star$. $\inf \in \star \sim \star$. $\omega_i \in \star \sim \star$. Etc *

1 35 2. * $n \in \text{Num} . \land `A = : ^*x \cdot x \in \text{Num} . \land . x \le n :$ $\rightarrow `\Sigma . x \in A . x \cdot = n (n+1)/10$ $: \land : \Sigma . x \in A . x^{10} \cdot = n^{11}/11 + n^{10}/10 + n/110 :$ $\text{Car } A = 11 . \land . A = \lceil a \rceil \cup \lceil b \rceil \cup \lceil c \rceil . \land \cdot f \in . A \curvearrowright \text{Com} :$ $\rightarrow : \Sigma . x \in A . fx \cdot = fa + fb + fc :$ $f \in . \bigstar \curvearrowright \text{Com} \cdot \land . A = \lceil \rceil : \Rightarrow : \Sigma . x \in A . fx \cdot = 0 :$ $1 + \text{Car } B . \in \text{Num} \cdot \land \cdot f \in . B \curvearrowright \text{Com} :$ $\rightarrow `\Sigma . x \in A . fx \cdot = : \Sigma \cdot x \in . A \cap B \cdot fx :$

```
1 + \operatorname{Car} C \cdot \in \operatorname{Num} \cdot \wedge \cdot f \in C \subset \operatorname{Com} \cdot \\ \wedge \cdot A \subset B \subset C \cdot \wedge \cdot \lceil a \rceil = \cdot B \setminus A : \\ \rightarrow : \Sigma \cdot x \in B \cdot fx = \cdot fa + \Sigma \cdot x \in A \cdot fx *
```

The use of the sum symbol (for finite sums only).

- 1 35 3. Finite and infinite sequences will be considered and written as functions. For the convenience of the reader we shall often use the customary index notation.
- 1 35 4. The definitive program will contain an exposition of Calculus. Though we shall use Calculus in the next chapters, we have omitted this exposition, because it cannot be our task to publish a text-book of Calculus in Lincos. But this argument cannot completely justify our silence in this matter. We should at least have dealt with notations. It is a fact that the usual system of notations of Calculus is inconsistent to a degree, especially in the field of functions of two and more variables. It is easy to remove these defects, but a notational system which preserves the obvious advantages of Leibniz's notation is still lacking. I have tried several methods but have been so far unable to decide for any one. I have therefore preferred to postpone the decision.
- 1 36 0. We now introduce a few words that will enable the user of Lincos to make statements in Lincos about some kinds of Lincos expressions. They are the words written

```
Ver (fL verum = "true") meaning "true",
Fal (fL falsum = "false") meaning "false",
Prp (fL propositio = "proposition") meaning "propositions".
Qus (fL quaestio = "question") meaning "questions".
Iud (fL iudicium = "judgement") meaning "truth value".
```

We shall also meet with the Lincos connective of negation, written —.

```
1 36 1. * 1+1=10. \in \text{Ver}:

1+1=11. \in \text{Fal}:

a+b=b+a. \in \text{Ver}:

\land \lceil a.b \rceil . a+b=b+a. \in \text{Ver}:

a+b \neq b+a. \in \text{Fal}:

1 \in \text{Int} . \in \text{Ver}:

1 \notin \text{Int} . \in \text{Fal}:

\sqrt{10} \in \text{Rea} . \in \text{Ver}:

\sqrt{10} \in \text{Rat} . \in \text{Fal}:

\text{Rat} \subset \text{Rea} . \in \text{Ver}:

\text{Rea} \subset \text{Rat} . \in \text{Fal}:

\text{Etc} *
```

```
1 36 2. * a \in A . \in \text{Ver} : \leftrightarrow : a \notin A . \in \text{Fal}:
a \in A . \in \text{Fal} : \leftrightarrow : a \notin A . \in \text{Ver}:
\land n : n \in \text{Num} . \rightarrow : \lor p : p \in \text{Pri} . \land . p \ge n : \in \text{Ver}:
\lor n : n \in \text{Num} . \land : \land p : p \in \text{Pri} . \rightarrow . p \ge n : \in \text{Fal}:
Etc. *
```

Some people might feel shocked by the use of the ' \in ' in such sentences, though I think it to be relatively harmless. Propositions have not so far occurred as elements of sets. We are therefore free to use the ' \in ' after a proposition as here. It would be more risky to admit 'Ver' and 'Fal' as elements of sets (especially 'Ver \in Agg' and 'Fal \in Agg').

We need not be afraid of the paradox of the liar and of other paradoxes. We have not yet dealt with human behaviour. We have not spoken about acts, nor even about the acts of pronouncing, proving and refuting propositions. In due course we shall learn that by behaviour rules some kinds of human acts are forbidden. One of these forbidden acts might be the famous act of the liar.

Note that our lexicological and syntactical apparatus does not suffice for a treatment of Lincos syntaxis. If we wish to pronounce both "x is an integer", and "x is a variable" in a common context, we are obliged to invent some means of distinguishing either between the two "x"-s or between the two copulas. (The first way out has become classical.)

```
1 36 3. p \in \operatorname{Ver} \rightarrow p \in \operatorname{Prp}
                  p \in \text{Fal.} \rightarrow p \in \text{Prp}
                  10 \rightarrow 1001 \cdot \notin \text{Prp}
                 a \in \text{Div.} \notin \text{Prp}
                  x \notin = . \notin \operatorname{Prp}
                  + \wedge \in \rightarrow . \notin Prp
                  \vee Num x \in 10 . \notin Prp:
                 Car A \neq 1. \rightarrow \ \ x \cdot x \in A \in B : \notin Prp.
                 Car A \neq 1 \rightarrow x \cdot x \in A Div b \neq Prp:
                 Car A \neq 1. \rightarrow \forall x \cdot x \in A = 11 : \notin Prp:
                  f \in A \cap B. \land a \notin A \rightarrow : Etc. f a \cdot  Etc. \notin Prp
                  Etc '
                  Ver \in Int . \notin Prp
                  Ver \notin Fal. \notin Prp
                  Etc #
1 36 4. # ? \cdot x \in \text{Rea.} \land .x^{10} = 10 : \notin \text{Prp}
                  ? \cdot x \in \text{Rea.} \land .x^{10} = 10 : \in \text{Qus}
                  ? x \in \text{Rea.} \land x^{10} = 10 : \in \text{Qus} \in \text{Ver}:
                  Etc +
```

- 1 36 5. $p \in \text{Ver} : \Leftrightarrow \text{. Iud } p = \text{Ver} : p \in \text{Fal} : \Leftrightarrow \text{. Iud } p = \text{Iud } p =$
- 136 6. $p \leftrightarrow p \in Ver$ $p \leftrightarrow p \notin Fal$ $\neg p \leftrightarrow p \in Fal$ $\neg p \leftrightarrow p \notin Ver$ $\neg l + l = ll$ $\neg Rea \subset Rat:$ $a \in A \leftrightarrow \neg a \notin A:$ $a \notin A \leftrightarrow \neg a \in A *$
- 136 7. * $p \land q . \in \text{Ver} \cdot \leftrightarrow \cdot p \in \text{Ver} . \land . p \in \text{Ver} :$ $p \land q . \in \text{Fal} \cdot \leftrightarrow \cdot p \in \text{Fal} . \lor . q \in \text{Fal} :$ $p \lor q . \in \text{Ver} \cdot \leftrightarrow \cdot p \in \text{Ver} . \lor . q \in \text{Ver} :$ $p \lor q . \in \text{Fal} \cdot \leftrightarrow \cdot p \in \text{Fal} . \land . q \in \text{Fal} :$ $p \to q . \in \text{Ver} \cdot \leftrightarrow \cdot p \in \text{Fal} . \lor . q \in \text{Ver} :$ $p \to q . \in \text{Fal} \cdot \leftrightarrow \cdot p \in \text{Ver} . \land . q \in \text{Fal} #$
- 1 36 8. # $\neg \cdot p \land q \cdot \leftrightarrow \cdot \neg p \cdot \lor \cdot \neg q :$ $\neg \cdot p \lor q \cdot \leftrightarrow \cdot \neg p \cdot \lor \cdot \neg q :$ $p \rightarrow q \cdot \leftrightarrow \cdot \neg p \cdot \lor q :$ $p \leftrightarrow \cdot \neg \neg p *$
- 1 36 9. * \neg : $\forall n \cdot n \in \text{Num}$. $\land : \land x \cdot x \in \text{Pri}$. $\rightarrow : x \leq n :$ $\land n \cdot n \in \text{Num}$. $\rightarrow \cdot \neg : \land x \cdot x \in \text{Pri}$. $\rightarrow : x \leq n :$ $\neg : \land x \cdot \text{Etc} x \cdot \text{Etc} : \leftrightarrow : \forall x \cdot \neg : \text{Etc} x \cdot \text{Etc} :$ $\neg : \forall x \cdot \text{Etc} x \cdot \text{Etc} : \leftrightarrow : \land x \cdot \neg : \text{Etc} x \cdot \text{Etc} *$

CHAPTER II

TIME

- 2000. Again we start with ostensive, ideophonetic signs, the so-called time-signals. They are even more ideophonetic than the peeps (written as dots) we used in Chapter I to introduce the natural numbers. While the peeps showed and meant arithmetical units, the new signs will not mean anything but themselves. So they can hardly be called words.
- 2010. The new signs are radio-signals time-signals of various duration and wave-length. They are written as horizontal lines.
- 2 01 1. * Dur = Sec a * and so on.

The 'a' as it stands does not belong to our program text in the proper sense. It is a meta-text variable, used as a substitute for a Lincos constant. Generally *Gothic letters* will represent such substitutes.

Eventually this 'a' should be replaced by a Lincos word meaning a positive real number a such that the sentence

"The duration of the factual time-signal indicated by the horizontal line is of a seconds"

is true.

The Lincos word written Dur (fL duratio=duration) means "duration". Syntactically it is to be handled as a function to the set of durations. The domain of this function is not exactly the set of time-signals. It is much broader, but at this stage it would be unwise and even impossible to circumscribe it in a too definite way. This is symptomatic of many functions we shall deal with.

The Lincos word written Sec means the time unit second. Syntactically it behaves as a function from 'Pos' to the set of durations. So it is a paradigm of Lincos syntaxis for physical units. 'Cmt' and 'Gra' (centimeter and gramme) will occur as symbols for functions from 'Pos' to the set of lengths and the set of masses respectively.

2012. The program text of 2011 will be repeated on different wavelengths and with different durations of their time-signals until the receiver may be expected to remark that the numbers indicated by a in the different texts are proportional to the durations he observes, independently of their wave-length. So he will guess that the numbers indicated by a are to mean the durations of those time-signals, and that the Lincos word 'Dur' is to mean "duration". He will learn our time-unit as well.

Note that his knowledge of our time-unit may be corrupted by a factor $1 - \frac{v}{c}$ (unrelativistically) as long as he does not know the relative velocity v of sender and receiver.

If, e.g., he is moving away from us, the waves of the end of a signal will travel a longer way to the sender than those of the beginning, and so the signal will be stretched. If he is moving towards us, it will be shortened (the "Olav Römer" effect). The receiver may know v from direct observation of our planet or of our solar system, but we cannot rely on this possibility. In due course we shall learn how to free time-reckoning from that unknown factor.

We could have used time-signals in order to introduce 'Rea'. But then we would have deprived ourselves of a means of showing duration.

It is important to send time-signals on different wave-lengths. Otherwise the number indicated by a could be understood as something relating to the number of waves or oscillations contained in the time-signal.

Now the a is a substitute for the positive number that counts the set of ether oscillations contained in the time-signal of the lefthand member of the text-equation.

The Lincos word written 'Nos' means "number of oscillations". It behaves as a function (from the set of time-signals or a larger set) to 'Pos'.

As the number indicated by a is different for time-signals of equal duration, but different wave-length, 'Nos' cannot be interpreted as "duration".

Intelligent beings who can receive electro-magnetic waves while ignorant of any theory of electromagnetic vibrations (e.g. beings who receive radio-signals by their senses as we do with light and sound) will not immediately understand the program texts of 2 01 3. But if they are not colour-blind (i.e. if they can distinguish wave-lenghts), they will perhaps learn from our information, and they will build some theory of radio-signals.

Note that, independently of the state of motion the number of oscillations in a signal given in our text will be equal to that observed by the receiver himself. within the limits of observation errors. All oscillations, from the beginning to the end of the signal, must reach the receiver, however distorted the signal may be.

The word written Fre (fL frequentia) means "frequency" (of an oscillation). 'Fre' and 'Sec⁻¹' ("Hertz") are treated as 'Dur' and 'Sec'. Again $\mathfrak a$ is a substitute for some real number that gives the frequency in Hertz of the time-signal in question.

2015. * Dur
$$x = \operatorname{Sec} a . \rightarrow .a > 0$$

Fre $x = \operatorname{Sec}^{-1} b . \rightarrow .b > 0$
Nos $x = c . \rightarrow .c > 0$:
Dur $x = \operatorname{Sec} a . \wedge .$ Fre $x = \operatorname{Sec}^{-1} b . \wedge .$ Nos $x = c . \rightarrow .ab = c$:
Sec $a \times \operatorname{Sec}^{-1} b = ab$.
 $c/\operatorname{Sec} a = \operatorname{Sec}^{-1} c/a$.
 $c/\operatorname{Sec}^{-1} b = \operatorname{Sec} c/b$ *

In a quite natural way we have switched over to variables as arguments of the new functions.

An example of arithmetics of concrete numbers has been given.

When the first time-signal finishes the second begins. The same fact is stated by the text within the "picture". The words written Fin (fL finis=end) and Ini (fL initium=beginning) mean "end" and "beginning" respectively. The two-time signals are sent on different wave-lengths. The receiver can therefore separate them and observe that the end of the one coincides with the beginning of the other. The text words are sent on still other wave-lengths. So the word group '=Ini' and the first signal will not disturb each other.

The difference of wave-length has been depicted by a difference in height.

2 01 7. * Ecc ____ :
$$\forall \neg x.y \neg \cdot \text{Fre } x = \mathfrak{a}. \land . \text{Fre } y = \mathfrak{b}. \land . \text{Ini } y = \text{Fin } x \text{ *}$$
 and so on.

The audio-pictorial part of 2015 has been repeated, but now together with a verbal description of the picture outside the picture. Of couse a and b are substitutes for such frequencies that

the text becomes true when related to the picture.

The word written $\operatorname{Ecc} (fL \text{ ecce} = \text{see})$ will always announce an audio-picture the legend of which is to be found outside the picture. Former texts may be repeated in order to explain this pattern:

2 01 8. # Ecc $\overline{}$: $\forall x \cdot \text{Dur } x = \mathfrak{a} \cdot \wedge \cdot \text{Nos } x = \mathfrak{b} \cdot \wedge \cdot \text{Fre } x = \mathfrak{c} + \mathfrak{c}$

and so on.

2 02 1. # Ecc ___ $\vee \lceil x \cdot y \rceil$ • Fre $x = \mathfrak{a} \cdot \wedge$ • Fre $y = \mathfrak{b} \cdot \wedge \cdot x$ Ant y # and so on.

The word written Ant (fL ante=before) means "before". The time-signals are separated by a pause. In the text they are named by their frequencies.

2 02 2. * $x \operatorname{Ant} y \cdot \leftrightarrow \cdot y \operatorname{Pst} x$ *

The word written Pst ($fL \operatorname{post} = \operatorname{after}$) means "after".

2 02 3.
$$\# x \operatorname{Pre} y . \Leftrightarrow \operatorname{Fre} x . = \mathfrak{w} \cdot \wedge \operatorname{Fre} y . = \mathfrak{w} \cdot \wedge . x \operatorname{Ant} y .$$

 $\wedge \cdot -_1 \vee z : \operatorname{Fre} z . = \mathfrak{w} \cdot \wedge . x \operatorname{Ant} z . \wedge . z \operatorname{Ant} y \#$

The word written Pre (fL precedit = precedes) means the relation of immediate succession, restricted to the case of time-signals of frequency w, where w is a meta-textual constant (the same in all Lincos texts).

On this frequency a clock will be installed. It will tick, one peep a second, through all Lincos programs from now on. The receiver will be instructed on how to read the clock:

2 02 4.
$$\# \text{Ini } x .= \text{Mom } n : \land .x \text{ Pre } y : \rightarrow : \text{Ini } y \cdot = \text{Mom } (n+1)^*$$

$$\text{Ini} = \text{Mom } 0 \#$$

The word written Mom means moment.

The moment Mom 0 is shown.

The rule will also be interpolated and extrapolated.

CH. II] TIME 83

> We may not use Sec n as a denotation of the moment Mom n, because Sec n means a duration.

> Time references will be used in order to mention past (or even future) events.

> The Lincos program clock will also possess hands for longer intervals (e.g. for 28 sec. and 216 sec.).

> By means of clock-reading, events like Ini x and Fin x may be considered as elements of the set of all moments Mom a (a real, even negative).

```
2025. \# \operatorname{Ini} x \in \operatorname{Tem} \cdot \operatorname{Fin} x \in \operatorname{Tem} \cdot
                        a \in \text{Rea.} \rightarrow .\text{Mom } a \in \text{Tem}
                        a 
leftharpoonup b \in \text{Rea.} \rightarrow : \text{Mom } a \cdot \text{Ant. Mom } b \cdot \leftrightarrow \cdot a < b
                        \operatorname{Mom} a. Ant. \operatorname{Mom} b \hookrightarrow \operatorname{Mom} a < \operatorname{Mom} b:
                        \operatorname{Mom} a = \operatorname{Mom} b : \Leftrightarrow a = b
                        \operatorname{Mom} a < \operatorname{Mom} b \rightarrow : \operatorname{Mom} b - \operatorname{Mom} a = : \operatorname{Sec} b - a
                        a \in \text{Tem} . \land . b \in \text{Tem} . \land . a < b \rightarrow . b - a \in \text{Dur}
                        Fin x - \text{Ini } x = \text{Dur } x:
                        \operatorname{Fin} x = \operatorname{Ini} y \longrightarrow \operatorname{Fin} y - \operatorname{Ini} x = \operatorname{Dur} x + \operatorname{Dur} y:
                        Sec. a+b=\cdot \operatorname{Sec} a+\operatorname{Sec} b:
                        Sec. a-b=·Sec a-Sec b:
```

 $a < b \rightarrow \operatorname{Sec} a < \operatorname{Sec} b$: $a = b \longrightarrow \operatorname{Sec} a = \operatorname{Sec} b *$

The word written Tem (fL tempus - time) means "time" (the time-axis). For elements of Tem a subtraction has been defined, but the differences are elements of 'Dur, not of Tem.

2.02.6. $\# Sec \in Pos \curvearrowright ^Dur$. $Dur \in \star \curvearrowright ^Dur$. $Nos \in \star \curvearrowright Pos$. $Fre \in \star \curvearrowright \hat{F}re$. Sec⁻¹ \in Pos \curvearrowright 'Fre. Fin∈★~Tem. $\operatorname{Ini} \in \star \curvearrowright \operatorname{Tem}$. $Mom \in \star \curvearrowright Tem #$

2031. Beside the verbal text the program may contain events. In broadcasting, the weather-forecast, the newsreel, reports and lectures are verbal program texts. A musical performance is an event, a radio play is a sequence of events most of which (but not all) consist of pronouncing a verbal text.

We will be able to mention past (and even future) events. Such

a mention is a proposition (but mentioning is again an event). We now exercise mentioning past events.

 $\mathfrak p$ is a metatextual substitute for an event, e.g. for a noise — i.e. a wild complex of radio signals that cannot be interpreted as a verbal text. The t_1 and t_2 are metatextual substitutes for Lincos words that designate some moments. The t_1 and t_2 above the line, surrounding the first $\mathfrak p$, do not relate to any part of the eventual program text, but to the activity of the clock. They should tell the terrestrial reader at which moments the event written $\mathfrak p$ has begun and has ended.

The word written Usq (fL usque ad = till) means "till". The word written Fit (fL fit = happens) means "happens". After the 'Fit' a copy of the first \mathfrak{p} occurs.

The first part of the text is an event that happened from t_1 to t_2 , namely the noise written $\mathfrak p$. The second part is a report that should be translated: From t_1 to t_2 $\mathfrak p$ happened. We shall see that reports need not be literal.

- 2 03 2. $\sharp t_1 \mathfrak{p} t_2 : \forall x.t_1 \operatorname{Usq} t_2 \operatorname{Fit} x : \forall x.t_1 \operatorname{Usq} t_2 \operatorname{Fit} x := \mathfrak{p} \sharp$
- 2 03 3. $\#^{t_1} p^{t_2} q^{t_3} t_1 Usq t_2 Fit p. \land .t_2 Usq t_3 Fit q. \land :t_1 Usq t_3 Fit \cdot p. q \#$
- 2034. $*t_1 \operatorname{Usq} t_2 \operatorname{Fit} \mathfrak{p}. \wedge .t_2 \operatorname{Usq} t_3 \operatorname{Fit} \mathfrak{q} \cdot \rightarrow : t_1 \operatorname{Usq} t_3 \operatorname{Fit} \cdot \mathfrak{p}. \mathfrak{q}$ *

 Note that the last t_1, t_2, t_3 are textual (written images of Lincos variables).
- 2 03 5. # Ecc t_1 —— t_2 Dur (x_1, t_1) Usq t_2 Fit $x_1 = t_2 t_1$ Fre (x_1, t_1) Usq (x_2, t_2) Fit (x_1, t_2) Fit (x_2, t_2) Fit (x_2, t_2) Fit (x_1, t_2) Fit (x_2, t_2) Fit (x_2, t_2) Fit (x_1, t_2) Fit (x_2, t_2) Fit (x_1, t_2) Fit (x_2, t_2) Fit (x_1, t_2) Fit (x_2, t_2) Fit (x_2, t_2) Fit (x_1, t_2) Fit (x_2, t_2) Fit (x_2, t_2) Fit (x_1, t_2) Fit $(x_2, t$

The horizontal line is again the picture of a time-signal. a is a substitute for its frequency.

As a rule propositions about events will mention the date of the event. An event does not contain its date; the date of an event can be read on the clock. For the terrestrial reader it is mentioned above the line.

For the sake of brevity only we shall often omit the 'Usq' in the date of an event which is mentioned in a proposition about that event.

¹⁾ Verbs are always considered to be used in the 3rd person singular. (Note that "verb" or "noun" are not Lincos syntactical or semantical categories. They refer to usual translations of Lincos words into English.)

2037. The rule that propositions about events should specify their date was not observed when we introduced the functions Dur and so on, and even later on, when we used the words written 'Ant' and 'Post'. We now give an explicit definition of 'Ant' that conforms to that rule:

*
$$x \operatorname{Ant} y . \Leftrightarrow : \forall \lceil t . t' . t'' . t''' \rceil \cdot t < t' < t'' < t''' .$$

$$\wedge . tt' \operatorname{Fit} x . \wedge . t'' t''' \operatorname{Fit} y *$$

2039. We shall also say

*
$$x \operatorname{Ant} y \cdot \leftrightarrow \cdot \operatorname{Ant} y \cdot \operatorname{Fit} x$$

 $x \cdot h \operatorname{Ant} \cdot y \cdot \leftrightarrow \cdot h \operatorname{Ant} \cdot y \cdot \operatorname{Fit} x$ *

so using 'Ant y' and 'h Ant y', as it were, as a date dating x.

*
$$y \operatorname{Pst} x . \leftrightarrow \cdot \operatorname{Pst} x . \operatorname{Fit} y$$
 $x . h \operatorname{Ant} . y \leftrightarrow \cdot y \cdot h \operatorname{Pst} . x$
 $y . h \operatorname{Pst} . x \cdot \leftrightarrow \cdot h . \operatorname{Pst} x \cdot \operatorname{Fit} y$ *

- 2040. In order to speak about an event ^{t₁} p ^{t₂}, we have used a copy of the event, whereas in formalist semantics it is agreed that things should not be used as words until they have been earmarked in a suitable way, e.g. by putting them between quotation marks. I concede our procedure is risky, but I do not believe that it is more so than any other method.
- 2041. Let us first analyze what quotation marks can do.

In due course we shall meet with formulas such as

$$t_1 t_2 \operatorname{Fit} \mathfrak{p} \longrightarrow (\forall x \cdot t_1 t_2 \operatorname{Fit} x).$$
 (**)

Now, there may be events that look like x. Then it could happen that

so that

$$\mathbf{t_1} \, \mathbf{t_2} \, \mathbf{Fit} \, x$$

would be true, as well as

$$\forall x \cdot t_1 t_2 \operatorname{Fit} x$$
,

though we should wish that no more than one of these formulas were propositions.

This can be avoided by using an earmarked p as a name of p.

Instead of quotation marks we may use a dot in the "lefthand ear" of \mathfrak{p} . When

we will speak about the p while using • p. So

The dot would be considered as belonging to the 'Fit' and it could even be incorporated into the 'Fit', if there were no formulas without a dot after the 'Fit'. Yet our formula (\star) will now read

$$t_1 t_2 \operatorname{Fit} {}^{\bullet} \mathfrak{p} . \leftrightarrow {}^{\bullet} \lor x . t_1 t_2 \operatorname{Fit} x$$
.

As x is an (ambiguous) name for the event in question, we have observed the rule that demands the use of names after 'Fit'. If now by chance $t_1 x_2 t_4$.

.

we have the opportunity to state

$$\mathbf{t_1} \, \mathbf{t_2} \, \mathrm{Fit} \, {}^{\bullet} \, x \, ,$$

which is not inconsistent with

$$\forall x \cdot t_1 t_2 \text{ Fit } x$$
.

It is a pity that in almost all cases the situation is much more complicated.

2 04 2. Let us suppose that there is a kind of events that may be characterized by some partial similarity, for instance events in which one can discover the algorithmic Lincos words for natural numbers as a part. These events may look like

where p and q are substitutes for some constant parts, whereas instead of 1010 we can meet with 101 or 11011 or some other numeral. Such events should be named

but if we restricted ourselves to this kind of name, speech about those events would become unusually clumsy. For instance we should be confined to a purely arbitrary name for the class of all these events, instead of an algorithmic name which tells us immediately which are the members of that class.

It is evident what we are to do. We shall name events looking like

by ambiguous names

$$ullet$$
 p $x\mathfrak{q}$

where x is a variable (not only: looks like a variable).

So, if

we are entitled to state

$$\forall x \cdot x \in \text{Num} . \land . t_1 t_2 \text{ Fit } \cdot p x q.$$

Unfortunately this is not completely correct.

$$\cdot pxq$$

is really the name of some

whereas we intended to name events such as

So it will be better not to use quotation, but quasi-quotation (as does Quine). The earmark • is to be changed into a quasi-earmark, let us say °. So

$$\forall x \cdot x \in \text{Num} . \land . t_1 t_2 \text{ Fit } \circ p x q.$$

The quasi-earmark will tell us that some parts of the text covered by it mean variables (not only look like variables), but it will not tell us which parts are exempted from its working. It will be merely a warning signal, which is due at almost all appearances of 'Fit'. So it is superfluous and unsatisfactory.

2 04 3. The only thing that matters is how to know which parts of '...' in 'Fit ...' are genuine, and which are substitutes. The latter should be anti-earmarked. As the substitutes both look like variables and are variables, one is tempted to use an extra-ordinary kind of variable in places like this. For instance such variables would be written as Gothic letters, whereas the genuine parts of the '...' in 'Fit ...' should be in italic. But this is a cumbersome procedure. We cannot prevent Gothic written words from appearing in authentic events such as

So we should be burdened with hypergothics and so on.

Moreover there is a much simpler device which works without earmarks and anti-earmarks.

2 04 4. Most of our program texts will be well poised speech. Parts of these texts may look like a hotchpotch if we try to read them as it were well poised speech. This is not important as long as we know that this hotchpotch is genuine. Fortunately we can easily recognize whether this is true or not. As soon as in a hotchpotch controlled by a 'Fit' substitutions have been made, so that

the text controlled by 'Fit' cannot be considered as the genuine name of some event, the text controlled by the 'Fit' will no longer look like a hotchpotch, at least not with respect to its relations with the surrounding context. A single non-genuine piece will mean a variable that does not truly belong to the text inside the working sphere of the 'Fit', but is imposed from outside. So quantification over such variables will take place outside. Conversely, if a variable suffers quantification outside the working sphere of the 'Fit', it is certain that this variable is not genuine inside. Moreover people from outside the 'Fit' who impose some substitution inside, will be well advised not to choose a variable x, if inside there is something that looks like x.

So, under our former assumptions

reports on an event

because there is no quantification over x, and consequently x is genuine.

$$t_1 t_2 \operatorname{Fit} \cdot \forall x \cdot p \ x q$$

reports an event

$$^{t_1} \lor x . p x q^{t_2}$$

because quantification took place inside the 'Fit' (it does not matter whether such an event tells us something or is merely a hotchpotch). In

$$\forall x \cdot x \in \text{Num} \cdot \land \cdot t_1 t_2 \text{ Fit } p x q$$
,

the x after the 'Fit' is not genuine because quantification has taken place outside.

The part under the influence of 'Fit' may be arbitrary. For instance

$$\forall x \cdot x \in \text{Num} . \land : t_1 t_2 \text{ Fit} \cdot \forall x . x$$

is a meaningful text, though the event mentioned in the text may be merely a noise.

2 04 5. I do not claim that this solves all problems that arise from mentioning an event, nor do I endeavour to solve them in future. The only thing I have tried is to get a working pattern.

The exposition has been somewhat abstract. Perhaps the reader will understand it better, if he re-reads it in connection with our analysis of oblique speech. Oblique speech in natural languages is a striking example of the occurrence of binding variables outside the speech. (See 3 01 5.)

- 2 04 6. Nothing would need to be added to the preceding exposition, if we had not developed the bad habit of dropping the 'A' quantification word in a large numbers of its occurrences. Speech would become more clumsy, if we were to abandon this habit, but fortunately we need not. The mere appearance of the variable x outside the working sphere of a 'Fit' betrays the non-genuine character of every piece inside that looks like x. (To add a remark that may seem unintelligible at the moment: that appearance need not be explicit.)
- 2050. Autonymity is the basis of our system of description. As a name of an event $\mathfrak p$ we use a copy of $\mathfrak p$. But the name $\mathfrak p$ will not unambiguously name the event $\mathfrak p$, unless $\mathfrak p$ has happened (and will happen) only once. This ambiguity cannot always be removed by mentioning the event $\mathfrak p$ that has happened from $\mathfrak t_1$ to $\mathfrak t_2$ by means of

(for instance if the date of the event p is to be concealed or asked for). Nevertheless p can be used as an unambiguous name of some event p as soon as it is clear from the context which event p is meant. This conforms to our general principles. In the sentence

we meet with the name "off" of an event of pronouncing "off", which in a certain context can be denoted unambiguously by this name. Under analogous conditions we shall use p as name of an event p in Lincos.

We can also mention an event through description, e.g. the event p that has happened from t_1 to t_2 by means of

$$\forall x . t_1 t_2$$
 Fit x .

From a formalist point of view such a description would be ambiguous, or rather, we should not be allowed to use the article because in every interval more than one event will happen. This formalist objection would even hold true if we restricted the use of 'Fit' to program events. For program events will be perceived, and perceiving a program event can also be a program event. This will yield at least two contemporary events. According to what is said in 025 we can use the expression

$$\mathbf{\hat{x}}$$
. $\mathbf{t_1}$ $\mathbf{t_2}$ Fit x

as an unambiguous description of that event, which is singled out by the date t_1 t_2 together with the nearer and farther context of the above expression. As long as we do not pursue formalization but communication, we can rest satisfied with this procedure.

CHAPTER III

BEHAVIOUR

3 00 1. For the time being it would be premature to try to describe human behaviour by a system of general rules like the mathematical and chronometric rules of the preceding chapters and some of the mechanical laws of the next chapter. Instead we shall show behaviour by quasi-general examples, from which the receiver may derive as many general behaviour rules as he pleases.

The things to be shown are events. The program of showing events could be transformed into a program of reporting on events.

Of course the receiver will understand that the events we are showing cannot have happened literally in the way we use to show them. He will understand that normal terrestrial events are not likely to be groups of radio-signals of such an intensity as to reach far-remote worlds. Nevertheless we shall behave as though this part of the program consists of events that are not only shown, but also happen, and we expect the receiver to do likewise.

Such fictions are not unusual. Spectators are expected to take it for granted that the events shown on the stage really happen. For this purpose it is essential that the events shown are not literally the same as those supposed to happen; for instance whispering is played in such a way that the discourse is heard in the house, though some people on the stage do not perceive it.

The aforesaid fiction is an essential tool in our program. The things shown are events that really happen though they do not happen literally in this way. This is not too unrealistic a view. Even in daily life the literal happening of an event is a fiction. All we know about an event is that we remember its happening, and though we do not remember it twice in precisely the same way, we state that it is the same event whenever we remember it.

3002. As the program events are to display behaviour, it is necessary for at least a part of them to be acts, i.e. caused by persons. Our Lincos vocabulary is still far from sufficient for introducing the bodies of the acting persons. So the only kind of act that can be

displayed immediately is the act of speaking. The Lincos word that designates this activity, is written Inq (fL inquit=says).

91

The terrestrial reader should guard against a too narrow interpretation of this 'Inq'. In the present chapter the physical background of the Inq-events, whether it be accoustical or optical or tactile or anything else, will remain undiscernible. 'Inq' will aim at communication acts in the broadest sense. After all we have said about literal happenings, it is clear that the same event can be played in different ways, but at the moment there is no sense in asking the question in which way it "really" happens.

- 3 00 3. The names of the dramatis personae will be written Ha, Hb, Hc, and so on. In due course we will state that these persons are members of the set called Hom (fL homo = man) in written Lincos.
- 3004. Our theatre is still incomplete. Besides persons and acts a third thing is needed. We have been able to build a vocabulary of mathematics without valuating our propositions. We had only to confine ourselves to true propositions. The falsehood of 1=2 could be formulated as $1 \neq 2$. Yet we cannot show behaviour by good actions only. We must stage bad ones too, if we wish to condemn them. We have to create a vocabulary that contains words meaning "good" and "bad" and intermediate valuations.

To begin with, we shall stick to two values. Of course they cannot be 'Ver' and 'Fal', which are values of propositions. What we wish to valuate are acts, not propositions. (We are not here considering value-judgements of esthetics.)

Our valuating words will be written Ben (fL bene=well) and Mal (fL male=badly). They mean "good" and "bad" respectively. These words can be introduced in a surprisingly simple manner. (This may be a striking refutation of a prejudice which is very common among representatives of Geisteswissenschaften: Werturteile other than "true" and "false" are said to be beyond the reach of mathematical methods.)

As we are neither inclined nor able to settle behaviour rules, we shall not pronounce valuating judgements otherwise than by the mouth of an acting person. 'Ben' and 'Mal' will be no more than expressions used by some people in order to approve or disapprove of some events. Occasionally pronouncing a true proposition may be judged to be a bad action, and pronouncing a false proposition a good one.

3 00 5. The elementary events are written

$Ha \operatorname{Inq} Hb \mathfrak{p}$.

This reads "Ha tells (says) Hb the p". Here p is again a metatextual substitute for something that can be said. This may also be a question, an incomplete answer, an order, a nonsensical group of words, a noise, and so on; it might even be an air sung by Ha.

- 3 00 6. The last point to be discussed is to know which kind of saying is the best to start with. As far as I have tried it, mathematics seems to provide the most appropriate texts. So we shall construct dialogues about mathematical topics, as it were classes in mathematics.
- 3 00 7. We have not endeavoured to give every acting person a constant character. The persons named Ha act independently in the different talks. This is a serious disadvantage. In a final redaction of our project the person named Ha (and likewise the others) should really be one and the same person during the whole program.
- 3 01 1. $*Ha \notin \text{Num.} Ha \notin \text{Com.} Ha \in Ha$. $Hb \notin \text{Num.} Hb \notin \text{Com.} Hb \in Hb$. Etc *Inq $\notin \text{Agg } *$
- 3 01 2. * $Ha \operatorname{Inq} Hb \cdot ? x \cdot 10 x = 101$: ${}^{t_1}Hb \operatorname{Inq} Ha \cdot 101/10^{t_2}$: $Ha \operatorname{Inq} Hb \operatorname{Ben} *$ and so on.

The incompleteness of answers is intentional. By means of this device the connexion with the preceding part of the text is stressed. The receiver will more easily understand that it is a dialogue. Yet one will also send:

```
* {}^{t_1}Ha \operatorname{Inq} Hb \cdot ? x \cdot 10 x = 101^{t_2}:

Hb \operatorname{Inq} Ha \cdot 10 x = 101 \cdot \rightarrow \cdot x = 101/10^{t_3}:

Ha \operatorname{Inq} Hb \operatorname{Ben}^{t_4} *

* Ha \operatorname{Inq} Hb \cdot ? x \cdot 10 x = 101:

{}^{t_1}Hb \operatorname{Inq} Ha \cdot x = 101/10^{t_2} ·

Ha \operatorname{Inq} Hb : {}^{\vee}x \cdot t_1 t_2 \operatorname{Fit} x \cdot \in \operatorname{Ben}^{t_3} *
```

Ha's last answer seems to be still incomplete. When " $x \cdot t_1 t_2$ Fit x" is replaced by " $Hb \operatorname{Inq} Ha \cdot x = 101/10$ " the answer will run " $Hb \operatorname{Inq} Ha \cdot x = 101/10$ " $\in \operatorname{Ben}$ ", and this is in fact incomplete, for it

mentions an event without the date that links the event to its context. This remark, however, is not correct. The aforesaid substitution is illegal. It is a fact that Ha said what he said, and not what we have got from it by some substitution. Quoting freely is permitted as long as it does not affect the meaning.

We have treated this question because it is the forerunner of more serious ones, not because the text could be a source of misinterpretations. Nevertheless we may add some comment to the talks under this number in order to ascertain which kind of substitutions are legal. We shall revert to this point in 301 s.

It is in our intention that the receiver should understand these sequences of questions and answers as dialogues. Yet we must also enable him to separate and to recognize the three components of 'Ha Inq Hb'. Therefore all talks will be repeated with other actors, so that 'Ha Inq Hb' cannot be read as a constant complex.

```
3 01 3. \# Ha \operatorname{Inq} Hb \cdot ?x \cdot 100x = 1010:
                    Hb \text{ Inq } Ha. 1010/100:
                    Ha Inq Hb Mal:
                    Hb \operatorname{Inq} Ha. 1/10:
                    Ha Inq Hb Mal:
                    Hb \operatorname{Inq} Ha . 101/10:
                    Ha \operatorname{Inq} Hb \operatorname{Ben} *
                * Ha \operatorname{Inq} Hb \cdot ?x \cdot x = 10 + 10:
                    Hb \operatorname{Inq} Ha \cdot 10 + 10:
                    Ha Inq Hb Mal:
                    Hb \operatorname{Inq} Ha 100:
                    Ha \operatorname{Inq} Hb \operatorname{Ben} *
                * Ha \operatorname{Inq} Hb \cdot ?x \cdot x^{10} = 11001.
                    Hb \operatorname{Inq} Ha \cdot 101 \times 101 = 11001:
                    Ha Inq Hb Mal:
                    Hb \operatorname{Inq} Ha \cdot 101 \times 101 = 11001 \cdot \in \operatorname{Ver}:
                    Ha \operatorname{Inq} Hb \operatorname{Ver} \operatorname{Tan} \operatorname{Mal} : \neg \cdot x^{10} = 11001 \rightarrow x = 101
                    Hb \operatorname{Inq} Ha \cdot 101 \sim -101:
                    Ha \operatorname{Inq} Hb \operatorname{Ben} *
                # Ha \operatorname{Inq} Hb \cdot ? \Upsilon x \cdot x^{10} = 11001:
                     Hb \operatorname{Inq} Ha 101:
                     Ha\operatorname{Inq} Hb\operatorname{Ben} *
                # Ha \operatorname{Inq} Hb \cdot ? \Upsilon x \cdot x^{10} = 11001
                     Hb \operatorname{Inq} Ha: {}^{\checkmark}x \cdot x^{10} = 11001 \cdot {}^{\checkmark}e \cdot {}^{\checkmark}x \cdot x^{10} = 11001 \cdot {}^{\checkmark}
                     Ha \operatorname{Inq} Hb \operatorname{Mal} #
```

? $\dot{}$ asks for an x (not for all x) satisfying the context. This notation is rooted in an older version of Lincos (compare 025). Though it does not fit in our present notational system, we do not feel an overwhelming need to change it.

Note the appearance of a word written Tan (fL tamen = nevertheless) meaning "but". Logical analysis seems to prove it to be superfluous. In the talk, as it runs, it could be replaced by ' \wedge '. Yet as an expression of behaviour it can hardly be dispensed with. What "but" means in behaviour might be illustrated by the two sentences

"He is a mathematician, but he is dull" and

"He is a mathematician, but he is clever".

Both sentences contain a thought that is not explicitly stated, "mathematicians are usually clever" in the first case, "mathematicians are usually dull" in the second.

It will be advantageous to translate the adversative "but" in a somewhat more circumstantial way. We shall also say instead of

- # Ver Tan Mal #
- * Qqm Ver Tan Mal *,

with a redundant 'Qqm' (fL quamquam=though). Then the last sentence reads "Ver, it it true, but Mal", or though Ver still Mal. Compare German "zwar" (zwar Ver, aber Mal). In the last pattern we may interchange the two components:

Tan Mal Qqm Ver #,

which reads "still Mal though Ver"; or more shortly

Mal Qqm Ver # .

All these patterns will be introduced in the same context, suitably repeated, as the original 'Tan'.

Here I must warn the reader against confusing 'Tan' with 'Sed', which will be introduced in 3 04 6. Compare the "but" of

He is not a clever man, but (Tan) he is very kind, and

He is not clever, but (Sed) dull, or of

He did not work hard, but (Tan) he passed the examination very well,

and

He did not work hard, but (Sed) wasted his time, or of

He is 16 years old, but (Tan) he looks like 18, and

He is not 16 years old, but (Sed) 18.

Though in many natural languages these two buts are usually identical, there is no reason whatsoever to stick to this idiomatic oddity in an artificial language. The *true adversative* 'Sed' and the *pseudo-adversative* 'Tan' are so well distinguished that one cannot expect these two notions to be identified in the thought of a being not influenced by our natural languages.

In German the difference between 'Tan' and 'Sed' is well-known. In many cases (not in all) it can be expressed by "aber (Tan)", and "sondern" (Sed). In written Dutch a similar distinction existed, but it fell into disuse.

3 01 4. # $t_1 t_2 Ha \operatorname{Inq} Hb p \leftrightarrow t_1 t_2 \operatorname{Fit} Ha \operatorname{Inq} Hb p #$

An abbreviation has been introduced. Do not confuse

 $t_1 Ha \operatorname{Inq} Hb \mathfrak{p}^{t_2}$ with $t_1 t_2 Ha \operatorname{Inq} Hb \mathfrak{p}$.

In the first case an event has been sent, in the second a report on this event.

3 01 5. #
$$Hc \operatorname{Inq} Hc : \land \lceil t_1 \cdot t_2 \cdot t_3 \rceil$$
 : $\lor t_4 \cdot t_1 t_2 Ha \operatorname{Inq} Hb \cdot ?x \cdot 100 x = 1010$:
 $\rightarrow t_2 t_3 Hb \operatorname{Inq} Ha \cdot 1010/100 \cdot \rightarrow t_3 t_4 Ha \operatorname{Inq} Hb \operatorname{Mal}$:
 $\land : t_2 t_3 Hb \operatorname{Inq} Ha \cdot 101/10 \cdot \rightarrow t_3 t_4 Ha \operatorname{Inq} Hb \operatorname{Ben}$ #

The person named Hc states a behaviour rule — perhaps he witnessed the dialogue of 3 01 3. It does not matter whether this rule is true or not. In any case Hc could be advised to suppose something about the lengths of the time-intervals within which the questions and answers should be comprised.

This is only a paraphrase of the preceding program text. Nevertheless it has an importance of its own, because it shows how to use variables in a context like this. We may not use one variable s instead of u and v, and we are not allowed to write

 $t_2 t_3 Hb \operatorname{Inq} Ha s$ instead of $t_2 t_3 Hb \operatorname{Inq} Ha u/v$,

and s = 1010/100 instead of $u = 1010 \cdot \land \cdot v = 100$, and s = 101/10 instead of $u = 101 \cdot \land \cdot v = 10$, for e.g. s = 101/10 would be true if s = 1010/100, and $t_2 t_3 Hb \ln Ha 101/10$ if s = 1010/100. In our actual formulation $t_2 t_3 Hb \ln Ha 101/10$ is only true if u = 101 and v = 10.

In this way more general behaviour rules may be stated:

$Hc \operatorname{Inq} Hc : t_1 t_2 Ha \operatorname{Inq} Hb : ? x \cdot a x = b :$

We face the problem of oblique speech here. Direct and oblique speech differ syntactically, not semantically, even in natural languages. In any case the semantic difference between

John said to me: "I am ill"

John told me that he was ill

is insignificant if compared with the syntactical transformation of "I" into "he" and "am" into "was". It is only exceptionally that a judge insists on a direct version of some witnessed statement talk, but even then he will not urge that the version be literal. It would be difficult to maintain that direct speech claims to be more literal than oblique speech.

In the oblique version of our example a variable (I) that belongs to John's discourse has been replaced by a variable (he) belonging to the discourse of the person who reports on John's speech.

In the examples

John said to me: "Mary is ill"

and

and

John told me that you were ill

(if said to Mary), it is a constant (Mary) that is replaced by a variable (you).

The "was" or "were" instead of "am" or "is" reflects the change from a (non-explicit) time variable "now" to another (non-explicit) time variable "recently".

It is an interlinguistic feature of oblique speech that inside the speech we find variables that are not genuine, but that belong to the discourse of the reporter. By synonymy this phenomenon may be obscured. Nevertheless it is clear that in

I bought a car, but John said that it was too expensive the "it" belongs to my discourse and not to John's.

The pattern of oblique speech is a special case of that of mentioning events (in an oblique way), namely the case where the event itself is speech. The last two program texts are good examples for the working pattern I have proposed in 2 04 0-2 04 6. In the last but one talk of Hc the variables u and v occurring in $t_2 t_3 Hb$ Inq Ha u' and $t_3 t_4 Hb$ Inq Ha v' cannot be variables used by Hb and Ha respectively, because they have been used by Hc previously ("previously" is not to be taken in the sense of temporal succession, but in that of logical structure; the ' \rightarrow ' has coordinating power). We could add $\wedge u$ to the discourse of Hc, but even if we abstain from doing so, it is evident that binding u must take place in Hc's discourse, at least if Hc is considered to be making a meaningful statement.

In the last talk there are still more variables in the oblique speeches of Ha and Hb that belong to the discourse of Hc, namely a, b, e, f and v. The subject of Hc's talk is not one question $? x \cdot ax = b$ (for all numerals a and b) imputed to Ha, but any question of the kind $? x \cdot ax = b$ (where a and b are numerals), e.g. $? x \cdot 100 x = 1000$, $? x \cdot 10 x = 11010$, and so on. The $? x \cdot ax = b$ is only a common name for all these questions, but it is not the question that is aimed at by Hc. And it is the same with Hb's answer. All these variables can be bound in Hc's discourse only.

3 01 6. There are more criteria that permit a distinction between 'Ben' and 'Mal':

```
# Ha Inq Hb · ? x · 100 x = 1010:

Hc Inq Ha · 101/10:

Ha Inq Hc Mal:

Hb Inq Ha 101/10:

Ha Inq Hb Ben #
```

Some comment of a fourth person on Ha's behaviour might be added.

3 01 7. The talk of 3 01 6 is perhaps the first from which the receiver may conclude that 'Ha Inq Hb ...' means "Ha says to Hb ..." and not "Hb says to Ha ...". But discourses like

Ha Inq Hb. t_1t_2 Fit $\mathfrak p$: Hb Inq Hc. t_1t_2 Fit $\mathfrak p$: Hc Inq Hd. t_1t_2 Fit $\mathfrak p$

* tha Inq Hb.t1t2 Fit pt.: Hb Inq Hc.t3t4 Ha Inq Hb.t1t2 Fit pt.:

Hc Inq Hd.t4t5 Hb Inq Hc:t3t4 Ha Inq Hb.t1t2 Fit p

are better suited for this purpose. We shall send a great many texts from which this will become still more evident.

- 3018. Comments on some texts may be very useful. E.g. on the first talk of 3012:
 - * $Hc \operatorname{Inq} Hd : t_1 t_2 Hb \operatorname{Inq} Ha : 10 x = 101 . \rightarrow .x = 101/10$ * $Hd \operatorname{Inq} Hc \operatorname{Ben} *$

A comment on the second talk of 3012:

+ $Hc \operatorname{Inq} Hd$ ' $t_1 t_2 Ha \operatorname{Inq} Hb$? $x \cdot 10 x = 101$: $\wedge : t_2 t_3 Hb \operatorname{Inq} Ha \cdot 101/10$: $Hd \operatorname{Inq} Hc \operatorname{Ben} +$

Another comment on the second talk of 3 01 2:

* $Hc \operatorname{Inq} Hd$ t₃ t₄ $Ha \operatorname{Inq} Hb$: $\forall x \cdot t_2 t_3 \operatorname{Fit} x \cdot \in \operatorname{Ben}$: $Hd \operatorname{Inq} Hc \operatorname{Ben}$ *

A comment on the third talk of 3012:

* $Hc \operatorname{Inq} Hd$ `t₂ t₃ $Ha \operatorname{Inq} Hb$: $Hb \operatorname{Inq} Ha$ 101/10 . $\in \operatorname{Ben}$: $Hd \operatorname{Inq} Hc$: Fal `t₂ t₃ $Ha \operatorname{Inq} Hb$: $\forall x$. $\operatorname{t_1}$ t₂ $\operatorname{Fit} x$ · $\in \operatorname{Ben}$: - `t₂ t₃ $Ha \operatorname{Inq} Hb$: $Hb \operatorname{Inq} Ha$ 101/10 . $\in \operatorname{Ben}$ ` \leftrightarrow `t₂ t₃ $Ha \operatorname{Inq} Hb$: $\forall x$. $\operatorname{t_1}$ t₂ $\operatorname{Fit} x$ · $\in \operatorname{Ben}$ *

From these comments the receiver will learn what liberties a person may take when quoting other people. One could add a hypercomment put into the mouth of still other persons and containing behaviour rules on honest quoting.

In a former version of Lincos we distinguished between literal and free quotations by means of a special notation which was dropped later on. Literal quotation is a rather unimportant limit position. We shall develop a means of comparing the exactness of quotations ('Err', 3 09 1, 3 19 1). This will prove to be more useful. If needed, literal quotations may be characterized by 'Err...=0'.

- 3 02 0. We shall here treat interrogative sentences:
- 3 02 1. Many interrogative pronouns and adverbs can be treated in the following manner:

```
+ t_1 Ha \operatorname{Inq} Hb \cdot ? x \cdot 100 x = 1010^{t_1}:

Hb \operatorname{Inq} Hc \cdot ? y : t_1 t_2 y \operatorname{Inq} Hb \cdot ? x \cdot 100 x = 1010:

Hc \operatorname{Inq} Hb Ha +
```

```
# t_1 Ha Inq Hb \cdot ? x \cdot 100 x = 1010^{t_2}
         Hb \operatorname{Inq} Hc ? y : t_1 t_2 Ha \operatorname{Inq} y : x \cdot 100 x = 1010:
          Hc \operatorname{Inq} Hb Hb #
 * t_1 Ha \text{ Inq } Hb \cdot ? x \cdot 100 x = 1010^{t_2} 
         Hb \operatorname{Inq} Hc: y \cdot t_1 t_2 Ha \operatorname{Inq} Hb y:
         Hc \operatorname{Inq} Hb \cdot ?x \cdot 100 x = 1010 *
+ t_1 Ha \text{ Inq } Hb ? x \cdot 100 x = 1010^{t_2}
         Hb \operatorname{Inq} Hc ? y : y Ha \operatorname{Inq} Hb ? x . 100 x = 1010:
         Hc \operatorname{Inq} Hb \cdot t_1 t_2 *
            These are the questions "who?", "to whom?", "what?",
"when?". In the same way one can express the factual question:
* t_1 Ha Ing Hb ? x . 100 x = 1010^{t_2}
         Hb \operatorname{Inq} Hc : y \cdot y = \operatorname{Iud} : t_1 t_2 Ha \operatorname{Inq} Hb : x \cdot 100 x = 1010
          Hc \operatorname{Inq} Hb \operatorname{Ver} *
 * t_1 Ha \text{ Inq } Hb : ? x \cdot 100 x = 1010^{t_2} 
         Hb \operatorname{Inq} Hc : y \cdot y = \operatorname{Iud} : t_1 t_2 Ha \operatorname{Inq} Hb : x \cdot 10 x = 101
         Hc \operatorname{Inq} Hb \operatorname{Fal} *
            The question "whether?" has been formulated. For the sake of
brevity we introduce a single word Utr (fL utrum = whether)
meaning "whether", by a behaviour rule:
# Hd \operatorname{Inq} Hd \cdot tt' Hb \operatorname{Inq} Hc \cdot i y \cdot y = \cdot \operatorname{Ind} \operatorname{Etc}
                     \leftrightarrow tt' Hb \operatorname{Inq} Hc \cdot \operatorname{Utr} \operatorname{Ete} *
            Likewise we can introduce another abbreviation:
                                   #?=Etc # instead of #?y.y=Etc #
and so on.
            With the means now available we can stress once more that in
 "Ha Inq Hb..." Ha is the person who is speaking:
* Ha \operatorname{Inq} Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x:
          Hb\operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2\operatorname{Fit} x:
          Hc \operatorname{Inq} Hd \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:
          Hd \operatorname{Inq} Hc \cdot t_1 t_2 \operatorname{Fit} \mathfrak{p}:
          Hc \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} \mathfrak{p}:
          Hb \operatorname{Inq} Ha \cdot t_1 t_2 \operatorname{Fit} \mathfrak{p}:
          Ha \operatorname{Inq} Hb . \operatorname{Ben} *
* Ha \operatorname{Inq} Hb \cdot \operatorname{Utr.} 100011111 \in \operatorname{Pri}^*
          Hb \operatorname{Ing} Hb : \neg \cdot 10 \operatorname{Div} \mathbf{Etc} \cdot \neg 11 \operatorname{Div} \mathbf{Etc} \cdot \neg \cdot 101 \operatorname{Div} \mathbf{Etc} \cdot \neg \cdot 100 \operatorname{Div} \mathbf{
                      - 111 Div Etc · 1011 Div Etc ·
```

 $Hb \operatorname{Inq} Ha . \notin \operatorname{Pri}^*$ $Ha \operatorname{Inq} Hb \operatorname{Ben} *$

3 02 2. It is somewhat more difficult to treat the identity question, e.g. questions such as

"who is Caesar?", "what is an integer?".

It would be of no use replacing them by

"what does the name Caesar mean?"

"what does the word integer mean?"

At this stage we are not able to introduce a word that means "means", and moreover, it has been our policy not to define, but to describe.

Yet the difficulty is not as great as one might believe. A text such as

```
# Ha \operatorname{Inq} Hb ? x \cdot x = Hc:

^{t_1}Hb \operatorname{Inq} HcHc^{t_2}:

Hb \operatorname{Inq} Ha \cdot ^{\vee}x \cdot t_1 t_2 Hb \operatorname{Inq} x Hc:

Ha \operatorname{Inq} Hb \operatorname{Ben} #
```

is clear enough. Caesar (Hc) has been introduced to Ha in accordance with etiquette. If Caesar is absent, the dialogue runs as follows:

```
# Ha \operatorname{Inq} Hb \cdot ? x \cdot x = Hc:

Hb \operatorname{Inq} Ha \cdot Hc = : \forall x \cdot t_1 t_2 x \operatorname{Inq} x \cdot Alea \text{ est iacta} :

Ha \operatorname{Inq} Hb \operatorname{Ben} *
```

Here t_1 t_2 is a substitute for some time-interval just before the crossing of the Rubicon. As we cannot name any other action but speaking, we have had to identify Caesar by one of his sayings — at the moment all his sayings are meaningless, so it does not matter which we choose.

The same pattern may be used for the question "what is an integer?"

```
# Ha \operatorname{Inq} Hb \cdot ?x \cdot x = \operatorname{Int} :
Hb \operatorname{Inq} Ha \cdot \operatorname{Exg} \cdot 0 \in \operatorname{Int} \cdot 1 \in \operatorname{Int} \cdot 1011 \in \operatorname{Int} \cdot -101 \in \operatorname{Int} \cdot \operatorname{Etc} :
11/101 \notin \operatorname{Int} \cdot \sqrt{10} \notin \operatorname{Int} \cdot \operatorname{Etc} :
x \in \operatorname{Int} \cdot \Leftrightarrow \cdot x = 0 \checkmark 1 \checkmark -1 \checkmark 10 \checkmark -10 \checkmark \operatorname{Etc} :
0 \in A \cdot \land \land x \cdot x \in A \cdot \Rightarrow : x + 1 \cdot \cup \cdot x - 1 \cdot \subset A : \Rightarrow \cdot \operatorname{Int} \subset A :
Ha \operatorname{Inq} Hb \operatorname{Ben} *
```

Exg (fL exempli gratia) means "for instance".

3 02 3. Our solution cannot be invalidated by the remark that the person named Hb could have answered 'Int=Int'. For then Ha should have replied 'Mal' as he would have done if ?=10+10 had been answered 10+10=10+10.

Another remark is more serious: How to formulate the identity question with respect to a thing Z for which the equation x=Z is void of sense? We may not argue that such events are impossible. It is true that if Lincos is endowed with a rational phonetic system, there will be a clear distinction between words Z for which the equation x=Z has some meaning, and words Z for which it is meaningless. So one should argue that the person named Ha would never ask e.g.

"
$$x \cdot x = \in$$
", but rather " $x \cdot a \cdot A \cdot a \cdot A \cdot a \cdot A$ ".

Yet this is not correct. We cannot forbid anyone asking meaningless questions. Moreover, the question " $x \cdot x = \in$ " is not at all meaningless, though it contains a syntactically bad part. Questions like this can be the adequate expression of some behaviour, e.g. they can have some didactical importance. By means of a syntactically wrong question Ha will be able to examine whether Hb has understood Lincos syntaxis.

```
* Ha \operatorname{Inq} Hb . ? = \in :
Hb \operatorname{Inq} Ha : \neg : \forall x \cdot x = \in . \in \operatorname{Prp} : \forall \lceil a . A \rceil . a \in A :
\operatorname{Exg} : 1 \in \operatorname{Num} . \sqrt{10} \in \operatorname{Rea} . \operatorname{Dur} \in \star \curvearrowright \star . \operatorname{Etc} :
Ha \operatorname{Inq} Hb \operatorname{Ben} : *
```

If Hb had restricted himself to the negative part of his answer, Ha would have shown his dissatisfaction by saying 'Mal'. For clearly Ha had asked for definite information about the use of ' \in ', and such information is not given in the first part of the answer.

- 3024. Other interrogative sentential forms are treated later on.
- 3 03 1. We send four separate time-signals of length $\mathfrak h$ on the frequency $\mathfrak a$ within the period from $\mathfrak t_0$ to $\mathfrak t_1$. The frequency $\mathfrak a$ is different from any frequency used for verbal texts (and from $\mathfrak w$). As usual the written image of a time-signal is a horizontal line.

```
# {}^{t_0}Ha \operatorname{Inq} Ha \operatorname{1um} \cdot Ha \operatorname{Inq} Ha \operatorname{10um} \cdot Ha \operatorname{Inq} Ha \operatorname{11um} \cdot Ha \operatorname{Inq} Ha \operatorname{100um} {}^{t_1} :

Hb \operatorname{Inq} Ha \cdot ? = {}^{\circ} \operatorname{Car} {}^{\circ} u :

\operatorname{Fre} u = \mathfrak{a} \cdot \wedge {}^{\circ} \operatorname{Ini} u \cdot \operatorname{Pst} {}^{t_0} \cdot \wedge {}^{\circ} \operatorname{Fin} u \cdot \operatorname{Ant} {}^{t_1} :

Ha \operatorname{Inq} Hb \operatorname{100} :

Hb \operatorname{Inq} Hb \cdot {}^{\circ} t_0 t_1 Ha \operatorname{Enu} :

? = {}^{\circ} \operatorname{Car} {}^{\circ} u : \operatorname{Fre} u = \mathfrak{a} \cdot \wedge {}^{\circ} \operatorname{Ini} u \cdot \operatorname{Pst} {}^{t_0} \cdot \wedge {}^{\circ} \operatorname{Fin} u \cdot \operatorname{Ant} {}^{t_1} :

Hb \operatorname{Inq} Ha \operatorname{Ben} #
```

The word written Enu (fL enumerat = counts) means "counts". We shall introduce many words that mean special kinds of speak-

ing. In some respects they will behave syntactically like 'Inq', though some of them will not have a personal object. The direct object of 'Enu' (the thing counted) will have to have the form of an interrogative sentence. Later on we shall adduce the reasons of this rule.

Ha uses ordinal numbers when counting. This seems to be the simplest way of introducing ordinal numbers. The termination of the ordinals is written um, in analogy with the Latin ordinals.

```
3 03 2. * Ha \operatorname{Inq} Hb^{!}?= \operatorname{`Car}: ^{\lambda}x \cdot x \in \operatorname{Pri}. \wedge .x \leq 100000!

^{t_0}Hb \operatorname{Inq} Hb^{!}: 10.1^{um}: 11.10^{um}: 101.11^{um}: 111.100^{um}:

1011.101^{um}: 1101.110^{um}: 10001.111^{um}: 10011.1000^{um}:

10111.1001^{um}: 11101.1010^{um}: 11111.1011^{um}: ^{t_1}

Hb \operatorname{Inq} Ha \operatorname{1011}:

Ha \operatorname{Inq} Ha^{!}: _{t_1}Ha \operatorname{Enu}: ?= \operatorname{`Car}: ^{\lambda}x \cdot x \in \operatorname{Pri}. \wedge .x \leq 100000:

Ha \operatorname{Inq} Hb \operatorname{Ben} *

3 03 3. * Ha \operatorname{Inq} Hb : ?= \operatorname{`Car}: ^{\lambda}x \cdot x \in \operatorname{Pri}. \wedge .x \leq 100000:

^{t_1}Hb \operatorname{Enu}. ?= \operatorname{Etc}^{t_2}:

Hb \operatorname{Inq} Hb \operatorname{1011}:

Ha \operatorname{Inq} Ha \cdot t_1 t_2 Hb \operatorname{Enu}. ?= \operatorname{Etc}:

Ha \operatorname{Inq} Hb \operatorname{Ben} *
```

The last program text is not simply a paraphrase of the last but one. It would be, if the whole program is considered as a report spoken by an anonymous person ("the sender"). If it is felt as a play (and the receiver should do this), the second text is rather a "parapraxis" of the first. In the first play the actor plays counting by pronouncing a sequence of numerals following the rules of counting. In the second play he shows another (not specified) behaviour that is characteristic of counting. For instance he whispers, he moves his fingers or a pencil, and so on. In the first paragraphs of the present chapter we announced that sometimes one and the same play will be performed in different versions.

The last text serves to free ' t_1 t_2 Hb Enu.? = ...' from the ' t_1 t_2 ' that makes it a proposition. We now dispose of 'Hb Enu.? = ...' which is an event like 'Ha Inq Hb ...'. This transformation can be stressed by saying

```
# t_1 t_2 Hb \operatorname{Enu}. ? = Etc. \Leftrightarrow: t_1 t_2 \operatorname{Fit} · Hb \operatorname{Enu}. ? = Etc. #
```

Numerous verbs will be treated in the same way.

By means of 'Etc' parts of the program text have been abbreviated. Gradually we shall form habits of omitting parts that are more or less evident:

```
* Ha \operatorname{Inq} Hb!? = \operatorname{`Car} : ^{x} : x \in \operatorname{Pri} . \land . x \leq 100000!

^{t_{1}}Hb \operatorname{Enu}^{t_{2}}!

Hb \operatorname{Inq} Ha 1011!

Ha \operatorname{Inq} Ha . t_{1} t_{2} Hb \operatorname{Enu}!

Ha \operatorname{Inq} Hb \operatorname{Ben} *
```

3 03 4. We send time-signals as in 3 03 1, but of different lengths, \mathfrak{h}_1 , \mathfrak{h}_2 , \mathfrak{h}_3 , \mathfrak{h}_4 .

```
***t**•______t***

Ha \operatorname{Inq} Hb : ? = \cdot^{\forall} u \cdot \operatorname{Dur} u :
^{\uparrow} u \cdot \operatorname{Fre} u = \mathfrak{a} \cdot \wedge \cdot u \operatorname{Pst} t_0 \cdot \wedge \cdot u \operatorname{Ant} t_1 :
Hb \operatorname{Inq} Ha : A = :^{\uparrow} u \cdot \operatorname{Fre} u = \mathfrak{a} \cdot \wedge \cdot u \operatorname{Pst} t_0 \cdot \wedge \cdot u \operatorname{Ant} t_1 :
\rightarrow : \operatorname{Dur} \cdot \operatorname{1um}^{\forall} u \cdot u \in A := \mathfrak{h}_1 : \wedge^{\circ} \operatorname{Dur} \cdot \operatorname{10um}^{\forall} u \cdot u \in A := \mathfrak{h}_2 :
\wedge^{\circ} \operatorname{Dur} \cdot \operatorname{11um}^{\forall} u \cdot u \in A := \mathfrak{h}_3 : \wedge^{\circ} \operatorname{Dur} \cdot \operatorname{Ultum}^{\forall} u \cdot u \in A := \mathfrak{h}_4 :
Ha \operatorname{Inq} Hb \cdot ? = i^{\operatorname{um}} :
Hb \operatorname{Inq} Ha :
i^{\operatorname{um}} u \cdot u \in A \cdot \operatorname{Ant} \cdot (i+1)^{\operatorname{um}} u \cdot u \in A : \wedge^{\circ} i^{\operatorname{um}} u \cdot u \in A :
= : \operatorname{Ultum}^{\forall} u \cdot u \in A \simeq : \operatorname{Ant} \cdot \operatorname{Ultum}^{\forall} u \cdot u \in A =
```

The syntaxis of ' n^{um} ' has been stated in a more definite way¹). Ult^{um} (fL ultimum = last) meaning "last", is a fairly important word. Its importance does not yet appear from this text.

The text may be repeated in abbreviated form:

3 04 1. * $Ha \operatorname{Inq} Hb \cdot ? = 101101 + 11101$: ${}^{t_1}Hb \operatorname{Inq} Hb \cdot 1 + 1 = 10 \cdot 1 + 0 + 0 = 1 \cdot 1 + 1 = 10 \cdot 1 + 1 + 1 = 10 \cdot 1 + 1 + 1 = 10 \cdot 1 + 1 + 0 = 10^{t_2}$: $Hb \operatorname{Inq} Ha \cdot 1001010$: $Ha \operatorname{Inq} Ha \cdot t_1 t_2 Hb \operatorname{Cpu} \cdot ? = 101101 + 11101$: $Ha \operatorname{Inq} Hb \operatorname{Ben} *$

¹⁾ Later on we mostly write 'i'm u' instead of 'i'm $^{\vee}u$ '. The second notation reflects an earlier version of Lincos. On a few places the notation has not been changed.

The word written Cpu (fL computat = computes) means "computes". Like 'Enu' it is constructed with an interrogative subsentence.

```
3 04 2. # Ha \operatorname{Inq} Hb \cdot ? = 11010 \times 1011 \cdot t_1 Hb \operatorname{Inq} Hb \cdot 11010 + 110100 + 11010000 = 100011110 \cdot t_2 Hb \operatorname{Inq} Ha \cdot 100011110 \cdot Ha \operatorname{Inq} Ha \cdot t_1 t_2 Hb \operatorname{Cpu} \cdot Ha \operatorname{Inq} Hb \operatorname{Ben} #
3 04 3. # Ha \operatorname{Inq} Hb \cdot ?x \cdot x^{10} - 1011x + 11110 = 0:
```

3 04 4. # $Ha \operatorname{Inq} Hb \cdot ?x \cdot x^{10} - 1011x + 11110 = 0$: $^{t_1}Hb \operatorname{Cpu} \cdot ?x \cdot x^{10} - 1011x + 11110 = 0^{t_2}$: $Hb \operatorname{Inq} Ha \cdot 101 \sim 110$: $Ha \operatorname{Inq} Ha \cdot t_1 t_2 Hb \operatorname{Cpu}$: $Ha \operatorname{Inq} Hb \operatorname{Ben} #$

 $Ha \operatorname{Inq} Hb \operatorname{Ben} #$

The 't₁ t₂ Hb Cpu...' has been freed from the 't₁ t₂'.

- 3 04 5. In 3 04 3 a word written Erg (fL ergo = thus) and meaning "thus" has been used. It is an important word that assimilates Lincos to common mathematical language. A person named Hc might give comment on the behaviour of Hb with respect to 'Erg'.
- 3 04 6. During, say, 8 seconds, from t_1 to t_2 , four short time signals a second will be given on the frequency a; $t_2 = t_1 + \text{Sec } 1000$.

```
# Ha \operatorname{Inq} Hb: ? = 'Car: ^u · Fre u = a \cdot \wedge \cdot t_1 Ant u Ant t_2:

Ecc:

t_1

Hb \operatorname{Inq} Hb: \operatorname{Car}: ^u · Fre u = a.

\wedge \cdot t_1 \leq t Ant u Ant (t + \operatorname{Sec} 1) \leq t_2' = 100:

\wedge \cdot t_2 - t_1 = \operatorname{Sec} 1000 \cdot \wedge \cdot 100 \times 1000 = 100000' Erg. 100000^{t_2}

Hb \operatorname{Inq} Ha \cdot 100000

Ha \operatorname{Inq} Ha \cdot \neg \cdot t_1 t_3 Hb \cdot \operatorname{Enu}. ? = . \operatorname{Car} \operatorname{Etc}:

Sed t_1 t_3 Hb \cdot \operatorname{Cpu} ? = . \operatorname{Car} \operatorname{Etc}:
```

'Sed' (fL sed = but) means "but", but in another sense than 'Tan' (see 3 01 3).

'Sed' is used when one proposition is contrasted with another that is contradictory to the first one. In a more specific way it can serve to single out that part of the two propositions that embodies the opposition between them:

$Ha \operatorname{Inq} Ha - t_1 t_3 Hb \operatorname{Enu} \cdot \operatorname{Sed} t_1 t_3 Hb \operatorname{Cpu} ? = \operatorname{Car} \operatorname{Etc} *$ or even

$$Ha\operatorname{Inq} Ha$$
 t_1t_3Hb \longrightarrow Enu $\operatorname{Sed}\operatorname{Cpu}$? = $\operatorname{Car}\operatorname{Etc}$ #

The original sentence has been transformed as though we had applied a law of distributivity. Of course word groups such as — Enu may only be used in contexts like the preceding program text. It makes no sense at all to interpret '— Enu' as the negation of an *event*. Negation can only be applied to *propositions*.

The aforesaid transformations can be exhibited by repeated talks and confirmed by behaviour rules.

```
3 05 1. * Ha \operatorname{Inq} Hb: Utr \land n \cdot n \in \operatorname{Num} \cdot \rightarrow : \forall p \cdot p \in Pri \cdot \land \cdot p > n:

t_1Hb \operatorname{Inq} Ha:

n \in \operatorname{Num} \cdot \rightarrow : \forall a \cdot \operatorname{Put} m \cdot a = 10 \times 11 \times 10 \times 101 \times \operatorname{Etc} \times n + 1:

\forall a \cdot \operatorname{Put} p \cdot a \in \operatorname{Pri} \cdot \land \cdot a \operatorname{Div} m:

\land x : x \leq n \cdot \rightarrow \cdot \neg \cdot x \operatorname{Div} m:

\operatorname{Sed} \cdot p \operatorname{Div} m \cdot \operatorname{Erg} : \land x \cdot x \leq n \cdot \rightarrow \cdot p \neq x:

\operatorname{Erg} \cdot p > n:

\operatorname{Erg} \cdot p \in \operatorname{Pri} \cdot \land \cdot p > n^{t_1}:

Ha \operatorname{Inq} Ha:

t_1 t_2 Hb \operatorname{Cpu} ? ? f : f \in \operatorname{Num} \curvearrowright \operatorname{Pri} \cdot \land \cdot \land n \cdot fn > n:

t_1 t_2 Hb \operatorname{Dem} : \land n \cdot n \in \operatorname{Num} \rightarrow : \forall p \cdot p \in \operatorname{Pri} \cdot \land \cdot p > n + n
```

The word written Dem (fL demonstrat = proves) means "proves". It behaves like 'Inq', though the indirect object is lacking. The direct object is a proposition.

3 05 2. The word written Put (fL puta = namely) means "namely". The use of 'Put' like that of 'Erg' tends to assimilate Lincos to colloquial mathematics. In mathematical texts it is even not unusual to designate the instance for an existential assertion with the same letter as that designating the existentially bound variable, without any comment. So a proof of the inverse of Cauchy's convergence criterion may run as follows:

"lim $a_n = a$. Thus given an $\varepsilon > 0$, there is such an n_0 that $|a - a_n| < \frac{1}{2}\varepsilon$ for all $n > n_0$.

Hence for all $m > n_0$, and $n > n_0$

$$|a_m-a_n|^2 \leq |a-a_m|+|a-a_n|<\frac{1}{2}\varepsilon+\frac{1}{2}\varepsilon=\varepsilon$$
.

In the second sentence of this proof n_0 was a bound variable, in the third sentence it is something like a constant.

By means of 'Put' we try to legalize this habit.

The syntactic features of 'Put' will become clear, when new variables occur inside the working sphere of the 'V' to which 'Put' belongs. In

 $\forall x$. Put a. Etc

or

 $\forall x . \text{Put } x . \text{Etc}$

the variable x is bound to a constant value. So it will not depend on not-explicitly-bound variables that might occur inside the working sphere of ' $\forall x$ '. This means that these variables must be supposed to be bound by generalizing inside, not outside. Of course x may depend on data which have occurred before.

- 3 05 3. The same goal can be reached by the use of a word written 'Pon' (L pone = put), meaning "put". One can introduce it by repeating the last program text with
 - * Pon m = Etc * instead of * $\forall a . \text{Put } m . a = \text{Etc}$ *

Later on 'Pon' will be used in a generalized sense.

3 06 1. # Ha Inq Hb · ? x · t₁ t₂ x Inq Ha p:

Hb Inq Ha · · · · · t₁ t₂ Hb Inq Ha p:

*Hb Inq Hc · Utr · t₁ t₂ Hc Inq Ha p:

Hc Inq Hb Fal:

Hb Inq Hd · Utr · t₁ t₂ Hd Inq Ha p:

Hd Inq Hb Fal:

Hb Inq He · Utr · t₁ t₂ He Inq Ha p:

He Inq Hb Fal:

Hb Inq Hf · Utr · t₁ t₂ Hf Inq Ha p:

tHf Inq Hb Ver:

Hb Inq Ha · t₁ t₂ Hf Inq Ha p:

Ha Inq Ha · t₃ t₅ Hb Pet · ? x · t₁ t₂ x Inq Ha p:

ha Inq Hb Ben #

The words written Pet and Rep (fL petit=seeks and reperit=finds) mean "seeks" and "finds" respectively. Syntactically they behave as 'Enu' and 'Cpu'.

3 06 2. # $Ha \operatorname{Inq} Hb : ? \cdot \operatorname{lum} \forall x \cdot x \in \operatorname{Pri} \cdot \wedge \cdot x > 10^{1001}$

```
^{t_1}Hb \operatorname{Ing} Hb \cdot 11 \operatorname{Div} (10^{1001} + 1) \cdot 10 \operatorname{Div} (10^{1001} + 10).
                       101 \operatorname{Div} (10^{1001} + 11) \cdot 10 \operatorname{Div} (10^{1001} + 100).
                       1011 \operatorname{Div} (10^{1001} + 101) \cdot 10 \operatorname{Div} (10^{1001} + 110).
                       11 Div (10^{1001} + 111) \cdot 10 Div (10^{1001} + 1000)^{t_2}.
                       (10^{1001} + 1001) \in \text{Pri}^{t_0}
                   Hb \operatorname{Ing} Ha . 1000001001 \in \operatorname{Pri}^{\bullet}
                  Ha \operatorname{Inq} Ha : t_1 t_3 \operatorname{Hb} \operatorname{Pet} \wedge t_3 t_4 Hb \operatorname{Rep} \cdot ? \cdot 1^{\operatorname{um}} x \operatorname{Etc}
                  Ha \operatorname{Inq} Hb \operatorname{Ben} *
3 06 3. # Ha \operatorname{Ing} Hb : ? \cdot 1^{\operatorname{um}} x \cdot x \in \operatorname{Pri} \cdot x \cdot x > 10^{1001}
                  ^{t_1}Hb \operatorname{Pet} ^{t_2}Hb \operatorname{Rep} ^{t_3}Hb \operatorname{Ing} Ha 10^{1001} + 1001
                  Ha \operatorname{Ing} Ha \cdot t_1 t_2 Hb \operatorname{Pet} \cdot \wedge \cdot t_2 t_3 Hb \operatorname{Rep}
                  Ha \operatorname{Inq} Hb \operatorname{Ben} *
                    This is now a well-known pattern of transformation (see 3 03 2-
               3 03 3, and 3 04 3-3 04 4).
                    Examples of 'Pet ^ . - Rep' can easily be added.
3 07 1. p, q, r will be metatextual substitutes for three different noises.
               # t1p.q.r:t2
                   Ha \operatorname{Inq} Hb \cdot i u \cdot t_1 t_2 \operatorname{Fit} u:
```

Ho Inq Ha'

p Par u . ^ . q Par u . ^ . r Par u . ^ : u Uni · p . q . r :

Ha Inq Ha : t₃ t₄ Hb Anl · <sup> \ \ \ \ u . t₁ t₂ Fit u :

Ha Inq Hb : ? x · x . 1 um Par . u :

t₅ Hb Inq Ha : p Ant q Ant r · Erg : p . 1 um Par . u · ^ · q . 10 um Par u ·

^ r . Ultum Par . u t₅ :

Ha Inq Ha : t₃ t₄ ~ t₅ t₆ . Hb Anl · <sup> \ \ \ \ u . t₁ t₂ Fit u :

Ha Inq Hb Ben *</sup></sup>

The word written Anl (fL analyzat) has a somewhat complex meaning. It is used for the act of structurally describing a structure of events (later on also with respect to spatial structures). Other translations might be "analyzes", "describes", "sketches".

The word written Par (fL pars = part) means "part". It should not be confused with the set-theory "inclusion", 'C'.

The word written Uni (fL unio = union) means "union". It should not be confused with the set-theory "union", 'U'.

We use these words if structural parts and their structural union are considered.

More contexts are needed for a better understanding of these new words.

Later on we shall also use a slightly different syntax of Uni.

We shall say

u Uni A

instead of

 $u \operatorname{Uni} z_1 z_2 \operatorname{Etc} z_n$,

if

$$A = \lceil z_1 \rceil \cup \lceil z_2 \rceil \cup \text{Etc} \cup \lceil z_n \rceil$$
.

This can be stated by a behaviour rule.

3 07 2. # t₁ p.p.p^{t₂}:

 $Ha\operatorname{Inq} Hb \cdot ?u \cdot t_1 t_2 \operatorname{Fit} u$:

t, Hb Inq Ha · t1 t2 11es Fit pt.

 $Ha \operatorname{Inq} Ha t_3 t_4 Hb \operatorname{Anl} \Upsilon u t_1 t_2 \operatorname{Fit} u$:

p.11es Par.u: A.u Uni.11es p.

 $Ha \operatorname{Inq} Hb \operatorname{Ben} *$

Multiplicative numerals ("n times") are formed with a termination es (from the Latin termination of almost all multiplicative numerals).

3 07 3. # t₁p.q.p.r.p.p.r.q.r.^{t₂}

 $Ha \operatorname{Inq} Hb \cdot ? u \cdot t_1 t_2 \operatorname{Fit} u$:

taHb Inq Ha . t1 t2. 100es Fit p ~ 10es Fit q ~ 11es Fit r. t4

 $Ha \operatorname{Inq} Ha : t_3 t_4 Hb \operatorname{Anl} \cdot \dot{u} : t_1 t_2 \operatorname{Fit} u$:

 $Ha \operatorname{Inq} Hb \operatorname{Utr} : t_1 t_2 \operatorname{Fit} \cdot \mathfrak{p} \cdot \mathfrak{p} \cdot \mathfrak{p} \cdot \mathfrak{q} \cdot \mathfrak{q} \cdot \mathfrak{r} \cdot \mathfrak{r} : \mathfrak{r}$

Hb Inq Ha Fal: Sed . p.q.p.r.p.p.r.p.r.

 $Ha \operatorname{Inq} Hb \operatorname{Ben} *$

3 07 4. t₁Ha Inq Ha 1um. Ha Inq Ha 10um.

Ha Inq Ha 11um, Ha Inq Ha 100um : t.

 $Hb \operatorname{Inq} Ha : ?x \cdot t_1 t_2 \cdot x^{es} \operatorname{Fit}^{-}$

Ha Inq Hb 100es:

 $Hc \operatorname{Inq} Hb \cdot ? y \cdot t_1 t_2 Ha \operatorname{Inq} Ha y$:

to Hb Inq Hc t1 t2 Ha Enu: ? x t1 t2. xes Fit - t4

 $Hc \operatorname{Inq} Hc: t_3 t_4 Hb \operatorname{Anl} \cdot \forall z \cdot t_1 t_2 Ha \operatorname{Inq} Ha z$:

 $Hc \operatorname{Inq} Hb$: ? $u t_1 t_2 Ha \operatorname{Enu}$: ? $x t_1 t_2 . x^{es} \operatorname{Fit}$: $\operatorname{Mod} u$:

Hb Inq Hc: t1 t2 Ha Enu 'Etc'

Mod: Ha Inq Ha · 1um . 10um . 11um . 100um:

Hc Inq Hb Ben #

According to this dialogue designating an act by 'Enu' may be a structural description of this act.

Furthermore we have learned the word written Mod (fL modo = in the manner). It is to conjoin a verb with a sentence that states in what manner the action described by the verb takes place (a

"modal subsentence"). We have also learned how to translate the modal interrogative adverb "how"? by "u. Etc. Mod u".

3 07 5. # t.Ha Inq Ha 100um. Ha Inq Ha 1000um.

Ha Inq Ha 1100um ta.

Hb Inq Ha: ? x · t1 t2 · xes Fit

Ha Inq Hb 1100es.

ta Hb Inq Hc: t1 t2 Ha Enu · ? x · t1 t2 xes Fit

Hc Inq Hb: ? ` u · t1 t2 Ha Enu · Etc · Mod u.

ta Hb Inq Hc: Mod · Inq 100um · 1000um · 1100um ta.

Hc Inq Hc: t3 t4 Hb Anl ` v · t1 t2 Ha Inq Ha v.

^ t5 t6 Hb Anl ` v · t1 t2 Ha Inq Ha v.

3 07 6. # $Ha \operatorname{Inq} Hb$. ? = 11010 × 1011 ${}^{t_1}Hb \operatorname{Inq} Ha$. 11010 + 110100 + 11010000 = 100011110 t_2 $Ha \operatorname{Inq} Ha$ · ${}^{t_1}t_2 Hb \operatorname{Cpu}$. ? = 11010 × 1011 t_2 $Hc \operatorname{Inq} Ha$: ${}^{t_2}t_3 Ha \operatorname{Anl}$ · ${}^{\checkmark}u$ · ${}^{t_1}t_2 Hb \operatorname{Inq} Ha u$ #

3 07 7. * $Ha \operatorname{Inq} Hb \cdot ? = 1111111^{10}$ ${}^{t_1}Hb \operatorname{Cpu} \cdot ? = 1111111^{10} \cdot {}^{t_1}$ $Hb \operatorname{Inq} Ha \cdot 111110000001$ $Ha \operatorname{Inq} Hb \cdot ? \overset{\checkmark}{} u \cdot Hb \operatorname{Cpu} \operatorname{Mod} u$ ${}^{t_1}Hb \operatorname{Inq} Ha :$ $\operatorname{Mod} \cdot (x+y)^{10} = x^{10} + 10 xy + y^{10} \cdot \wedge \cdot x = 10^{110} \cdot \wedge \cdot y = -1 \cdot {}^{t_1}$ $Hc \operatorname{Inq} Ha : {}^{t_2} t_3 Ha \operatorname{Anl} \cdot \overset{\checkmark}{} u \cdot {}^{t_1} t_2 Hb \operatorname{Inq} Ha u \overset{*}{}^{*}$

3 07 8. $*Ha \operatorname{Inq} Hb : \operatorname{Utr} : \land n : n \in \operatorname{Num} : \rightarrow : \forall p : p \in \operatorname{Pri} : \land : p > n :$

 $^{t_1}Hb\operatorname{Dem}: \wedge n \cdot n \in \operatorname{Num}. \rightarrow : \vee p \cdot p \in \operatorname{Pri}. \wedge . p > n^{t_2}:$ $Hc\operatorname{Inq} Ha: ? \vee u \cdot ^{t_1t_2}Hb\operatorname{Dem}. \operatorname{Etc}. \operatorname{Mod} u:$ $^{t_3}Ha\operatorname{Inq} Hc:$ $\operatorname{Mod} \cdot Hb\operatorname{Pet}: ? \vee p \cdot p \in \operatorname{Pri}. \wedge . p\operatorname{Div}(10 \times 11 \times 100 \times \operatorname{Etc} \times n + 1)^{t_2}:$ $Hc\operatorname{Inq} Hc: t_3t_4Ha\operatorname{Anl} \cdot \vee v \cdot t_1t_2Hb\operatorname{Inq} Hbv =$ Likewise one should exhibit the use of 'Anl' with respect to 'Pet' and 'Rep'.

3 07 9. * t_1 p.q.r t_1 : $Ha \operatorname{Inq} Hb \cdot ?u \cdot t_1 t_2 \operatorname{Fit} u$: $^{t_2}Hb \operatorname{Inq} Ha$ $p \operatorname{Par} u \cdot \wedge \cdot q \operatorname{Par} u \cdot \wedge \cdot r \operatorname{Par} u \cdot \wedge \cdot u \operatorname{Uni} \cdot p \cdot q \cdot r$: $Ha \operatorname{Inq} Hc : t_3 t_4 Hb \operatorname{Anl} \cdot \forall u \cdot t_1 t_2 \operatorname{Fit} u$: $Hc \operatorname{Inq} Ha \cdot ? \forall v : t_3 t_4 Hb \operatorname{Anl} \cdot \forall u \cdot t_1 t_2 \operatorname{Fit} u \cdot \operatorname{Mod} v$: $^{t_2}Ha \operatorname{Inq} Hc \cdot \operatorname{Mod} Hb \operatorname{Pet} : ? = \cdot \uparrow_{z \cdot z} \operatorname{Par} u \cdot \iota_{z \cdot z}$ $Hc \operatorname{Inq} Hc : t_5 t_6 Ha \operatorname{Anl} \cdot \forall w \cdot t_3 t_4 Hb \operatorname{Inq} Ha w + \iota_{z \cdot z}$

3 08 1. We now build nouns (nomina actionis) from verbal stems by means of a termination ¹⁰ (a Latin termination for building nomina actionis).

Remark: No importance should be attached to the words verb and noun in this connection. They do not describe Lincos categories whether syntactical or semantical. They refer rather to conventional translations of Lincos texts into common English. It would be possible to translate Inq, Enu, Cpu, Dem, Pet, Rep, Anl not by verbal forms, but by nomina agentis, e.g. $t_1 t_2 Ha$ Cpu x by "Ha is the computer of x during $t_1 t_2$ ".

```
3 08 3. # Ha \operatorname{Inq} Hb \cdot ? = \operatorname{Cpu}^{io}:

Hb \operatorname{Inq} Ha \cdot \operatorname{Exg}^{\bullet} tt' Hc \operatorname{Inq} Hd \cdot ? \operatorname{Ete} \cdot \wedge \cdot t' t'' Hd \operatorname{Cpu} \cdot ? \operatorname{Ete} :

\rightarrow : \forall u \cdot t' t'' Hd \operatorname{Inq} Hd u \cdot \in \cdot \operatorname{Cpu}^{io} \cdot ? \operatorname{Ete} #
```

Likewise one should introduce Demio, Petio, Repio, Anlio.

```
3 09 1.  ^{\sharp t_1}\mathfrak{p} \cdot \mathfrak{q} \cdot \mathfrak{r}^{t_2} \vdots 
Ha \operatorname{Inq} Hb \cdot ?x \cdot \mathfrak{t}_1 \mathfrak{t}_2 \operatorname{Fit} x \vdots 
 ^{t_3}Hb \operatorname{Inq} Ha \colon \operatorname{Car} \dot{} ^{\uparrow}y : y \cdot \operatorname{Par} \dot{} ^{\vee}x \operatorname{Etc} \vdots = 11^{\mathfrak{t}_4} \vdots 
Ha \operatorname{Inq} Hb \overset{\circ}{} ^{\vee}z \cdot \mathfrak{t}_3 \mathfrak{t}_4 Hb \operatorname{Inq} Ha \cdot \in \operatorname{Anlio} \dot{} ^{\vee}x \operatorname{Etc} : 
\operatorname{Tan} \dot{} ^{?}y \overset{\circ}{} u \in \operatorname{Anlio} \dot{} ^{\vee}x \operatorname{Etc} \wedge : \operatorname{Err} u < \cdot \operatorname{Err} \dot{} ^{\vee}z \operatorname{Etc} \vdots 
 ^{t_3}Hb \operatorname{Inq} Ha \overset{\circ}{} \mathfrak{p} \operatorname{Par} \dot{} ^{\vee}x \wedge \cdot \mathfrak{q} \operatorname{Par} \dot{} ^{\vee}x \wedge \cdot \mathfrak{r} \operatorname{Par} \dot{} ^{\vee}x 
 \wedge : \overset{\circ}{} x \operatorname{Etc} \cdot \operatorname{Uni} \overset{\circ}{} \mathfrak{p} \cdot \mathfrak{q} \cdot \mathfrak{r} \overset{\circ}{} \overset{\circ}{} ^{t_3} \operatorname{Etc} : 
 \wedge : \operatorname{Err} \dot{} ^{\vee}u \operatorname{Etc} \cdot \mathsf{c} \cdot \operatorname{Err} \dot{} ^{\vee}z \operatorname{Etc} : 
 \operatorname{Anlio} \dot{} ^{\vee}x \operatorname{Etc} : \wedge : \operatorname{Err} v \cdot \mathsf{c} \cdot \operatorname{Err} \dot{} ^{\vee}u \operatorname{Etc} : 
 \operatorname{Ha} \operatorname{Inq} Ha \overset{\circ}{} \mathfrak{p} \operatorname{Par} \dot{} ^{\vee}x \wedge \cdot \mathfrak{q} \operatorname{Par} \dot{} ^{\vee}x \wedge \cdot \mathfrak{r} \operatorname{Par} \dot{} ^{\vee}x 
 \wedge : \overset{\circ}{} x \operatorname{Etc} \cdot \operatorname{Uni} \overset{\circ}{} \mathfrak{p} \cdot \mathfrak{q} \cdot \mathfrak{r} : \wedge \cdot \mathfrak{p} \operatorname{Ant} \mathfrak{q} \operatorname{Ant} \mathfrak{r} : 
 \operatorname{Ha} \operatorname{Inq} Hb \operatorname{Ben} \overset{\sharp}{} \overset{\sharp}{} \overset{\sharp}{} \operatorname{Ben} \overset{\sharp}{} \overset{
```

The word written Err (fL error = error) is used when the exactness of analyses (Anlio) or other statements are to be compared. The three analyses to be found in the program text go in the

direction of increasing exactness and decreasing inexactness (Err). Later on we shall meet with 'Err' in more numerical contexts.

Of course the use of Err does not at all imply the existence of a simple ordering of statements with respect to exactness.

One should repeat former analyses in order to compare their exactness by means of 'Err'.

3 10 1. * ${}^{t_1}Ha \operatorname{Inq} Hb \cdot ?x \cdot 1010x = 101^{t_2}$! $Hb \operatorname{Inq} Ha \cdot 1/10^{t_2}$! $Hc \operatorname{Inq} Hd \cdot t_2 t_3 Hb \operatorname{Rsp} Ha$! $Hd \operatorname{Inq} Hc \cdot ? = \operatorname{Rsp}$! $Hc \operatorname{Inq} Hd \cdot \operatorname{Exg} \cdot tt' Ha \operatorname{Inq} Hb \cdot ?x \cdot \operatorname{Etc} x \operatorname{Etc} :$ $\wedge \cdot t' t'' Hb \operatorname{Inq} Ha \cdot \operatorname{Etc} 1/10 \operatorname{Etc} \rightarrow \cdot t' t'' Hb \operatorname{Rsp} Ha$ *

The word written Rsp (fL respondit = answers) means (the verb) "answers".

- 3 10 2. * ${}^{t_1}Ha \operatorname{Inq} Hb \cdot ? x \cdot 1010 x = 101^{t_2}$: $Hb \operatorname{Inq} Ha \cdot 101/1010^{t_2}$: $Hc \operatorname{Inq} Hd \cdot t_2 t_3 Hb \operatorname{Rsp} Ha *$
- 3 10 3. * ${}^{t_1}Ha \operatorname{Inq} Hb \cdot ? x \cdot 1010 x = 101^{t_2}$: $Hb \operatorname{Inq} Ha \cdot 10^{t_2}$: $Hc \operatorname{Inq} Hd \cdot t_2 t_3 Hb \operatorname{Rsp} Ha \cdot \operatorname{Tan} \operatorname{Fal} *$
- 3 10 4. $*^{t_1}Ha \operatorname{Inq} Hb \cdot 10111111 \in \operatorname{Pri} \cdot {}^{t_2}Hb \operatorname{Inq} Ha \operatorname{Fal} {}^{t_3} \cdot Hc \operatorname{Inq} Hd \cdot t_2 t_3 Hb \operatorname{Rsp} Ha *$
- 3 11 1. $^{\# t_1}\mathfrak{p}.\mathfrak{q}.\mathfrak{r}^{t_2}$: $Ha \operatorname{Inq} Hb \cdot ?x. t_1 t_2 \operatorname{Fit} x$: $^{t_3}Hb \operatorname{Inq} Ha : t_1 t_2 \operatorname{Fit} \cdot \mathfrak{p}'.\mathfrak{q}. \mathfrak{r}^{t_4}$: $Ha \operatorname{Inq} Hb \cdot \operatorname{Fal} : t_3 t_4 Hb \operatorname{Mut} \cdot {}^{\vee}x. t_1 t_2 \operatorname{Fit} x \cdot \operatorname{Mod} \cdot \mathfrak{p}' \operatorname{Ilo} \mathfrak{p} \ ^{\#}$

The word written Mut (fL mutat = changes) means "changes". The word written Ilo (fL in loco = instead of) means "instead of".

3 11 2. $^{\sharp t_1} \mathfrak{p} \cdot \mathfrak{q} \cdot \mathfrak{r}^{t_2}$: $Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x :$ $^{t_3} Hb \operatorname{Inq} Ha \cdot \mathfrak{q} \cdot \mathfrak{r}^{t_4} :$ $Ha \operatorname{Inq} Hb \cdot \operatorname{Fal} : t_3 t_4 Hb \operatorname{Mut} \cdot {}^{\vee} x \cdot t_1 t_2 \operatorname{Fit} x \cdot \operatorname{Mod} \cdot \operatorname{Sin} \mathfrak{p}$

The word written Sin (fL sine=without) means "omits".

```
3 11 3.  ^{*t_1}p \cdot q \cdot r^{t_2}:

Ha \operatorname{Inq} Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x:

^{t_2}Hb \operatorname{Inq} Ha \cdot p \cdot \hat{s} \cdot q \cdot r^{t_4}
Ha \operatorname{Inq} Hb \cdot \operatorname{Fal} : t_3 t_4 Hb \operatorname{Mut} \cdot {}^{\checkmark}x \cdot t_1 t_2 \operatorname{Fit} x \cdot \operatorname{Mod} \cdot \operatorname{Add} \hat{s} 

The word written Add (fL addit = adds) means "adds".
```

```
3 11 4. # Ha \operatorname{Inq} Hb.? = 111111110:

Hb \operatorname{Cpu}.? = 11111110:

Hb \operatorname{Inq} Ha 111110000001:

Ha \operatorname{Inq} Hb \cdot ? u. Hb \operatorname{Cpu} \operatorname{Mod} u:

Hb \operatorname{Inq} Ha \cdot \operatorname{Mod}: Ha \operatorname{Mut} \cdot (x+y)^{10} = x^{10} + 10 xy + y^{10} \cdot 

\operatorname{Mod}. 1000000 \operatorname{Ilo} x \curvearrowright (-1) \operatorname{Ilo} y:

Ha \operatorname{Inq} Hb \operatorname{Ben} #
```

3 12 0. After the interrogative "how?" we now treat "why?". Its Lincos equivalent will be written Cur (†L cur = why). Answers to questions 'Cur Etc' may begin with a word written Qia (†L quia = because). The latter word will also be used independently, as an inverse 'Erg' as it were. In this case it introduces an answer to a non-explicit question.

The polysignificance of "why" and its equivalents in other languages is a commonplace. One should at least distinguish "whys" asking for the reason, for the cause, and for the aim. If one tries to examine these notions more closely one runs the risk of getting entangled in a muddle of relations. From experiences of this kind one ultimately learns that in some cases it may be better to avoid such analyses or at least to delay them, if progress is to be made. An interlinguistic homonymy can have deep roots in human behaviour. Perhaps the questioner uses a polysemantic "why" because he often does not know which kind of "because" he is asking for. I think it is a useful rule to discriminate homonyms where this is easy, and to justify them when not. In the latter case one should invent syntactic rules that cover the polysemantic use of the word.

The question 'Cur' as we shall put it, will ask for a justification of the fact that some event or some series of events has happened within a larger series of events. These events may also be of a linguistic kind, i.e. occurrences of certain words or groups of words in a certain context. In any case the things to be justified will be mentioned by means of the word Cur.

```
3 12 1. * Ha \operatorname{Inq} Hb \cdot \sqrt{10} \notin \operatorname{Rat}!
Hb \operatorname{Inq} Ha \operatorname{Cur}!
Ha \operatorname{Inq} Hb \cdot \operatorname{Qia} \cdot \operatorname{Qqm} \cdot \sqrt{10} \notin \operatorname{Rat} \cdot \operatorname{Pon} \cdot \sqrt{10} \in \operatorname{Rat}!
\operatorname{Erg} : \vee \lceil p \cdot q \rceil \cdot \operatorname{Put} \lceil p \cdot q \rceil \cdot p \cup q \cdot C \operatorname{Num} \cdot \land : \neg \cdot 10 \operatorname{Div} p \cdot \land \cdot 10 \operatorname{Div} q : \land \cdot p / q = \sqrt{10}!
\operatorname{Erg} \cdot p^{10} / q^{10} = 10 : \operatorname{Erg} \cdot p^{10} = 10 q^{10}!
\operatorname{Erg} \cdot 10 \operatorname{Div} p : \operatorname{Erg} \cdot 100 \operatorname{Div} p^{10}!
\operatorname{Erg} \cdot 100 \operatorname{Div} (10 q^{10}) : \operatorname{Erg} \cdot 10 \operatorname{Div} q^{10} : \operatorname{Erg} 10 \operatorname{Div} q :
\operatorname{Erg} : 10 \operatorname{Div} p \cdot \land \cdot 10 \operatorname{Div} q \cdot \land : \neg \cdot 10 \operatorname{Div} p \cdot \land \cdot 10 \operatorname{Div} q :
\operatorname{Erg} : \vee s \cdot s \in \operatorname{Prp} \cdot \land \cdot s \wedge \cdot \neg s \cdot s \cdot fal :
\operatorname{Erg} \cdot \sqrt{10} \in \operatorname{Rat} \cdot \in \operatorname{Fal}!
\operatorname{Hb} \operatorname{Inq} \operatorname{Ha} \operatorname{Ben} *
```

The 'Cur' of this text asks for the reason.

The question, could also have been asked in the form 'Cur \notin ' or 'Cur $\cdot \notin$ Rat' or 'Cur $\cdot \bigvee 10 \notin$ Rat' or 'Cur Ver', and so on. The talk should be repeated with these variations.

```
3 12 2. * t_0 Ha \operatorname{Inq} Hb^{t_1} p \in \operatorname{Pri}^{t_1} \land \cdot a \cup b \cdot \subset \operatorname{Num} \cdot \land \cdot p \operatorname{Div}(ab):
                               \rightarrow p \operatorname{Div} a \cdot \vee p \operatorname{Div} b^{t_1}
                         Hb \operatorname{Ing} Ha : \operatorname{Cur} \cdot \mathbf{t}_1 \mathbf{t}_2 Ha \operatorname{Ing} Hb \cdot p \in \operatorname{Pri} \mathbf{l}
                         Ha Inq Hb : Qia
                               110 Div (11 \times 100). A: -110 Div 11. \vee .110 Div 100:
                              tt' Ha Mut: {}^{\checkmark}x \cdot t_0 t_3 Ha Inq Hb x: Mod \cdot Sin \cdot p \in Pri
                              \rightarrow: \forall y \cdot tt' Ha \operatorname{Inq} Hb y \in \operatorname{Fal:Nes} t_1 t_2 Ha \operatorname{Inq} Hb \cdot p \in \operatorname{Pri} I
                         Hb \operatorname{Inq} Ha \cdot \operatorname{Cur} : t_1 t_2 Ha \operatorname{Inq} Hb \cdot p \in \operatorname{Pri} :
                              Sed: \neg 'Etc: \lor \neg q \cdot n \neg \neg q \in \text{Pri} \cdot \land \cdot n \in \text{Num} \cdot \land \cdot p = q^n \mid
                         Ha\operatorname{Inq} Hb ' Qia'
                               100 \, \text{Div} \, (10 \times 110) \, \land : \neg \cdot 100 \, \text{Div} \, 10 \, \lor \cdot 100 \, \text{Div} \, 110:
                              tt' Ha \operatorname{Mut}: {}^{\vee}x \cdot t_0 t_3 Ha \operatorname{Inq} Hb x:
                              \operatorname{Mod} \cdot \vee \lceil q \cdot n \rceil \operatorname{Etc} \cdot \operatorname{Ilo} \cdot p \in \operatorname{Pri}
                               \rightarrow: \forall y \cdot tt' Ha \operatorname{Inq} Hb y \in \operatorname{Fal} \operatorname{Nes} \cdot \operatorname{t}_1 \operatorname{t}_2 Ha \operatorname{Inq} Hb \cdot p \in \operatorname{Pri} \operatorname{I}
                         Hb \operatorname{Ing} Ha \operatorname{Cur} - \operatorname{t}_1 \operatorname{t}_2 Ha \operatorname{Ing} Hb \cdot p \in \operatorname{Pri} \cdot \vee \cdot p = 1 \mathsf{I}
                         t_4Ha \operatorname{Inq} Hb : \neg : \forall s \cdot t_4 t_5 Ha \operatorname{Inq} Hb \cdot \operatorname{Qia} s:
                              \rightarrow: \forall y \cdot tt' Ha \operatorname{Inq} Hb y \in \operatorname{Ver} : \neg \operatorname{Nes} \cdot \operatorname{Ete} \cdot p \in \operatorname{Pri} *
```

The 'Cur' of this talk asks what is aimed at by some act? Ha is giving a justification for having made some supposition.

We have incidentally learned the word written Ncs (fL necesse = necessary) that means "it is necessary that". The vocabulary of modality will be dealt with in detail later on.

```
3 12 3. * t<sub>1</sub> p t<sub>1</sub>:

Ha Inq Hb · ? x . t<sub>1</sub> t<sub>2</sub> Fit x:

t<sub>1</sub>Hb Inq Ha p t<sub>1</sub>:

Hc Inq Hb · Cur . t<sub>3</sub> t<sub>4</sub> Hb Inq Ha p:

Hb Inq Hc : Qia * t<sub>2</sub> t<sub>3</sub> Ha Inq Hb · ? x . t<sub>1</sub> t<sub>2</sub> Fit x: ∧ . t<sub>1</sub> t<sub>2</sub> Fit p:

Hc Inq Hb Ben *
```

This is a 'Cur' asking for something like the cause of an event.

3 13 0. We now turn to the word written Sci (fL scit=knows). In many languages the semantical domain of "to know" is claimed by two words (or clans of words) though the division between them is made rather on syntactical than on semantic grounds. In English this distinction may be circumscribed as that between "to know a fact" and "to know a thing", Latin "scire" and "noscere", French "savoir" and "connaître", German "wissen" and "kennen".

On the same grounds as given in the case of "why" we will stick to one word 'Sci'. This use will be legalized by the rule: A sentence beginning with 'Sci' is to be considered as an answer to a (not necessarily explicit) question. Or more syntactically: the sentence that depends on 'Sci' (the thing known) is to be interrogative (perhaps elliptical).

We now give a few examples of intermediate translations of sentences to be translated into Lincos:

"Ich weiss, wo mich der Schuh drückt" becomes "Ich weiss: "wo drückt mich der Schuh?""

"Ich weiss nicht, was soll es bedeuten" becomes "Ich weiss nicht: "was soll es bedeuten?""

"Kennst du das Land, wo die Zitronen blühn?" becomes "Weisst du: "in welchem Lande blühen die Zitronen?""

"I know Caesar" becomes "there is an x with the properties that x is Caesar and that I know: "is this x Caesar?""

"Eritis sicut Deus scientes bonum et malum" becomes "Like God for all x you will know: "which is the value of x?""

"Scio nil scire" becomes "There is no p such that I know: "is p true?" and I know: "is there any p such that I know: "is p true?"?""

"Γνωτι σε αυτον" becomes "You shall know: "who are you?""

"I know a Roman citizen" may not be translated by "I know: "who is a Roman citizen?"", but "There is an x such that x is a Roman citizen, and such that I know: "is x a Roman citizen?"".

The following is to justify the syntactical rule we have imposed on 'Sci'. Let us suppose we wish to say that Ha knows the solution x of $100 \ x = 1010$. As 'Hb·Sci $100 \ x = 1010$ ' clearly does not fit, we could try Hb Sci ' $^{\gamma}x$. $100 \ x = 1010$. But then it would not be clear what Hb is said to know. It may be the "bad" solution 1010/100, i.e. the fact that $100 \ x = 1010 \ . \rightarrow . x = 1010/100$, whereas we wish to assert that he knows the good one, i.e. $100 \ x = 1010 \ . \rightarrow . x = 11/10$. Hb's ignorance may be still worse; perhaps he does not know any more than that 'x. $100 \ x = 1010$ is a solution of $100 \ x = 1010$, i.e. he can apply a syntactical pattern, but he has no notion of algebra.

Fortunately, we have created the means and established the habit of distinguishing between good and bad solutions. The approved method of asking for good solutions is to use the word written "?". So 'Hb Sci · ? x · 100 x = 1010' expresses fairly well what we wish to say. More fully 'Sci' then reads: "knows the good answer to . . .".

As soon as a syntactical rule for 'Sci' is agreed on in one case, the best we can do is to stick to this rule generally. Thus we were obliged to transform "Ha knows that p" into the interrogative mode. The inverse way is not practicable. Indeed, we are not allowed to translate

Ha knows the solution of $?x \cdot 100x = 1010$

by

$$Ha \, \text{Sci} \cdot 100 \, x = 1010 \, . \rightarrow . \, x = 101/10,$$

because this translation would convey not only the fact that Ha knows the solution, but also the solution itself. We must be able to say that Ha knows something without betraying the precise structure of the thing known.

It might have been useful to introduce a word that announces the x that is considered to be a good answer to a question. If this word is written as a note of exclamation, we should get, e.g.

$$!x.100x = 1010 := 101/10.$$

Under this condition 'Sci' should be subjected to the syntactical rule that the thing known has to have the form of an answer to a (not necessarily explicit) question. We shall not treat the problems that would arise from this alternative convention, but we may mention that they are closely akin to some problems to be dealt with in 3 17 0.

As 'Sci' is always followed by an interrogative sentence one might enquire whether we cannot get rid of the question mark after 'Sci' by means of a fusion with the preceding 'Sci'. But this

is not possible, for we met with interrogative sentences without an introductory question mark, viz. those that start with 'Utr' or 'Cur'. Apart from this fortuitous circumstance, I would not favour dropping the question mark linking 'Sci' to the activity of solving problems.

As a final remark it is perhaps not superfluous to mention that our 'Sci' does not aim at a (more or less mysterious) mental state of the person knowing, but simply at a certain behaviour that will be exhibited by means of a number of examples.

As to the syntactic rules for 'Enu', 'Cpu' and others, the reader will understand why it has been useful to pattern them on that of 'Sci'.

It also goes without saying that 'Ha Sci p' is the written picture of an event, whereas ' t_1 t_2 Ha Sci p' is a report about an event. This report can also be written ' t_1 t_2 Fit. Ha Sci p'.

```
3 13 1. # * ^{t_0}Ha \operatorname{Inq} Hb \cdot ?x \cdot 100 x = 1010 ?
^{t_1}Hb \operatorname{Inq} Ha \cdot 101/10^{t_0} ?
Ha \operatorname{Inq} Hc \cdot t_1 t_2 Hb \operatorname{Sci} \cdot ?x \cdot 100 x = 1010 ?
Hc \operatorname{Inq} Ha \cdot \operatorname{Cur} \operatorname{Ver} ?
Ha \operatorname{Inq} Hc \cdot \operatorname{Qia} * t_0 t_1 Ha \operatorname{Inq} Hb \cdot ?x \cdot 100 x = 1010 ?
\wedge * t_1 t_2 Hb \operatorname{Inq} Ha \cdot 101/10 ?
Hc \operatorname{Inq} Ha \operatorname{Ben} #
# * ^{t_0}Ha \operatorname{Inq} Hb \cdot ?x \cdot 100 x = 1010 ?
^{t_1}Hb \operatorname{Inq} Ha \cdot 1010/100 ?^{t_0} ?
Ha \operatorname{Inq} Hc \cdot \neg : t_1 t_2 Hb \operatorname{Sci} \cdot ?x \cdot 100 x = 1010 ?
Hc \operatorname{Inq} Ha \cdot \operatorname{Cur} \neg :
Ha \operatorname{Inq} Hc \cdot \operatorname{Qia} : \neg : t_1 t_2 Hb \operatorname{Inq} Ha \cdot 101/10 \cdot \operatorname{Sed} 1010/100 #
```

3 13 2. From the last talks the receiver will not yet guess the true meaning of 'Sei'. 'Hb Sei p' could mean "Hb communicates the solution of p". Knowing has been recognized on the strength of symptoms that might appear to be too direct.

```
* {}^{t_1}Ha \operatorname{Inq} Hb \cdot ? x \cdot 100 x = 1010 :

{}^{t_1}Hb \operatorname{Inq} Ha \cdot 1010/100^{t_2} :

Ha \operatorname{Inq} Hc : t_1 t_2 Hb \operatorname{Sci} \cdot ? u \cdot t_0 t_1 Ha \operatorname{Inq} \operatorname{Hb} u :

Hc \operatorname{Inq} Ha \operatorname{Cur} :

{}^{t_2}Ha \operatorname{Inq} Hc : \operatorname{Qia} \cdot t_0 t_1 Ha \operatorname{Inq} Hb \cdot ? x \cdot 100 x = 1010 :

\wedge : {}^{\checkmark}x \cdot 100 x = 1010 \cdot = \cdot {}^{\checkmark}u \cdot t_1 t_2 Hb \operatorname{Inq} Ha u^{t_2} :

Hd \operatorname{Inq} He \cdot ? v \cdot t_3 t_4 Ha \operatorname{Inq} Hc v :

He \operatorname{Inq} Hd : t_3 t_4 Ha \operatorname{Inq} Hc : t_1 t_2 Hb \operatorname{Rsp} Ha \cdot

\operatorname{Erg} : t_1 t_2 Hb \operatorname{Sci} \cdot ? u \cdot t_0 t_1 Ha \operatorname{Inq} Hb u :
```

```
Ha \operatorname{Inq} Hc : t_1 t_2 Hb \operatorname{Sei} ? y : t_0 t_1 y \operatorname{Inq} Hb : ? x . 100 x = 1010 :
Hc \operatorname{Inq} Ha \operatorname{Cur} :
Ha \operatorname{Inq} Hc : \operatorname{Qia} . t_1 t_2 Hb \operatorname{Rsp} Ha :
t_1 t_2 Hb \operatorname{Sei} ? y : t_0 t_1 Ha \operatorname{Inq} y : ? \operatorname{Etc} :
t_1 t_2 Hb \operatorname{Sei} ? z : z Ha \operatorname{Inq} Hb : ? \operatorname{Etc} :
t_1 t_2 Hb \operatorname{Sei} ? z : z Ha \operatorname{Inq} Hb : ? \operatorname{Etc} :
t_1 t_2 Hb \operatorname{Sei} ? x . t_0 t_1 \operatorname{Fit} w #

3 \ 13 \ 3. \ ^* t_0 p ^{t_1} :
Ha \operatorname{Inq} Hb : ? x . t_0 t_1 \operatorname{Fit} x :
t_2 Hb \operatorname{Inq} Ha p ^{t_0} :
Ha \operatorname{Inq} Hc : t_2 t_3 Hb \operatorname{Sei} : ? y . y \operatorname{Fit} p #
```

After these talks it is unlikely that the receiver could understand 'Sci' in the sense of "communicates the solution of . . .". Yet the evidence for 'Sci' may be still more indirect.

3 13 6. Knowledge may begin, last and end:

```
(t_1 < t_2 < t_3 < t_4 < t_5)
# Ha \operatorname{Inq} Hb : ? \operatorname{1um} x \cdot x \in \operatorname{Pri} \cdot \wedge \cdot x > 10^{1001}:

{}^{t_4}Hb \operatorname{Inq} Ha 1000001001^{t_5}:
```

```
Ha \operatorname{Inq} Ha \cdot t_4 t_5 Hb \operatorname{Sci} \operatorname{Etc}:
                         Ha \operatorname{Inq} Hb Utr: t_4 t_5 Hb \operatorname{Pet}? 1^{\operatorname{um}} x. Etc.
                         Hb Inq Ha Fal:
                         Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} \cdot t_4 t_5 Hb \operatorname{Sei} \operatorname{Ete}:
                         Hb \operatorname{Inq} Ha t_1 t_2 Hb \operatorname{Pet} ? l^{\operatorname{um}} x \operatorname{Etc}:
                               t_2 t_3 Hb \operatorname{Rep} : \operatorname{Erg} \cdot t_2 t_3 Hb \operatorname{Sci} : \operatorname{Erg} \cdot t_4 t_5 Hb \operatorname{Sci} :
                               t < t' < t'' < t'''. A. tt' Hb Sci Etc \rightarrow . t'' t''' Hb Sci Etc ^*
3 13 7.
                      (t_1 < t_2 < t_3 < t_4 < t_5 < t_6)
                    * t_1Ha \operatorname{Inq} Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x^{t_2}
                         Hb Inq Ha · Cur . Ha Inq:
                         Ha \operatorname{Ing} Hb \cdot \operatorname{Qia} \longrightarrow \operatorname{t}_5 \operatorname{t}_6 Ha \operatorname{Sei} \cdot \operatorname{?} x \cdot \operatorname{t}_1 \operatorname{t}_2 \operatorname{Fit} x
                         Hb \operatorname{Inq} Ha \operatorname{Tan} \cdot t_3 t_4 Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} a:
                               Erg. t<sub>3</sub> t<sub>4</sub> Ha Sei Ete
                         Ha\operatorname{Inq} Hb \operatorname{Qqm} \cdot t_3 t_4 Ha\operatorname{Sci} : \operatorname{Tan} \cdot - \cdot t_5 t_6 Ha\operatorname{Sci} *
                     + Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x
                         {}^{\mathbf{t_5}}Hb\operatorname{Inq} Ha: \operatorname{Qqm} {}^{\mathbf{t_3}}t_4Hb\operatorname{Sci}: \operatorname{Utr} \cdot \vee p . t_1\,t_2Ha\operatorname{Inq} Hb\,p:
                         \operatorname{Tan}: \neg \cdot t_3 t_4 Hb \operatorname{Sei} \cdot ? = p \operatorname{Etc}^{t_4} :
                         Ha \operatorname{Inq} Hb \cdot t_1 t_2 Ha \operatorname{Inq} Hb a
                         {}^{t_{\bullet}}Hb \operatorname{Inq} Ha {}^{\bullet} \wedge {}^{\vdash} t \cdot t' \cap : t_{5} < t < t' \cdot \rightarrow : tt' Hb \operatorname{Sei} \cdot ? x \cdot t_{1} t_{2} \operatorname{Fit} x #
3 13 8. # Ha \operatorname{Ing} Hb : ! \operatorname{1um} x \cdot x \in \operatorname{Pri} A \cdot x > 10^{1001}
                         t_1Hb \operatorname{Ing} Ha \operatorname{Qqm} : - t_1 t_2 Hb \operatorname{Sci} \operatorname{Etc} \cdot \operatorname{Tan} t_3 t_4 Hb \operatorname{Sci}^{t_5}.
                         Hb Pet: 1^{\text{um}} x \cdot x \in \text{Pri} \cdot \land x > 10^{1001} \cdot 10^{1001}
                         Hb \operatorname{Inq} Ha 1000001001: t_i
                         Ha \operatorname{Inq} Ha \cdot t_3 t_4 Hb \operatorname{Sei}:
                         Ha Inq Hb. Cur Sci.
                         Hb \operatorname{Inq} Ha: t_2 t_3 Hb \operatorname{Pet} \operatorname{Pep} \cdot \operatorname{Erg} t_3 < t' < t'' . \to . t' t'' Hb \operatorname{Sci} *
3 13 9. (t_1 < t_2 < t_3 < t_4 < t_5 < t_6)
                     # t_5 Ha Inq Hb Utr: t_5 t_6 Hb Sci ? x . t_1 t_2 Fit x^{t_6}:
                          Hb \operatorname{Inq} Ha \operatorname{Ver}
                         Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} : t_1 t_2 Hb \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:
                          Hb Inq Ha Fal t<sub>3</sub>t<sub>4</sub> Hc Inq Hb . t<sub>1</sub>t<sub>2</sub> Fit a
                          t < t' < t'' < t''' < t^{(4)} < t^{(5)}. A t' Fit x \cdot \wedge \cdot t'' t''' Hc Inq Hb . <math>tt' Fit x:
                                \rightarrow: t^{(4)} t^{(5)} Hb Sci \cdot ? y \cdot tt' Fit y #
 3 14 1. # {}^{t_3}Ha \operatorname{Ing} Hb \cdot \operatorname{Utr} : {}^{t_3}{}^{t_4}Hb \operatorname{Sei} \cdot {}^{?}x \cdot {}^{t_1}{}^{t_2}\operatorname{Fit} x^{t_4} :
                          Hb \operatorname{Inq} Ha \operatorname{Ver} :
                          Ha \operatorname{Inq} Hb : \operatorname{Utr} \lor \ulcorner t \cdot t' \cdot x \cdot y \urcorner : t' < t_3 \cdot \land tt' y \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} x :
                          Hb Inq Ha Fal:
                          Ha \operatorname{Inq} Hb \operatorname{Utr} : \operatorname{t}_1 \operatorname{t}_2 Hb \operatorname{Sei} \cdot \operatorname{?} x \cdot \operatorname{t}_1 \operatorname{t}_2 \operatorname{Fit} x:
```

```
Hb \operatorname{Inq} Ha \operatorname{Ver} :
Ha \operatorname{Inq} Hb \operatorname{Cur} :
Hb \operatorname{Inq} Ha : t_1 t_2 Hb \operatorname{Ani} : ? x . t_1 t_2 \operatorname{Fit} x #
```

 $(t_1 < t_2 < t_3 < t_4 < t_5 < t_6)$

Hb knows that something has happened though nobody told him about it. He knew it as soon as it had happened. The reason he assigns is that he witnessed that event.

The word written Ani (fL animadvertit = perceives, is aware of) means "witnesses". Syntactically it behaves as "Sci". Of course 'Hb Ani Etc' is the written picture of an event that is reported on by ' t_1 t_2 Hb Ani Etc'.

```
Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:
                          ^{t_{5}}Hb\operatorname{Inq}Ha . t_{1}\,t_{2}\operatorname{Fit}\mathfrak{a}^{t_{6}} :
                          Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} \cdot t_1 t_2 Hb \operatorname{Ani} \operatorname{Etc} :
                          Hb \operatorname{Inq} Ha \operatorname{Fal} :
                          Ha \operatorname{Ing} Hb \operatorname{LUtr} \vee \operatorname{d} t \cdot t' \cdot y \operatorname{lt}_2 < t \cdot \wedge \cdot tt' y \operatorname{Ing} Hb \cdot t_1 t_2 \operatorname{Fit} \mathfrak{a}
                          Hb \operatorname{Ing} Ha \operatorname{Fal} :
                          Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} : t_5 t_6 Hb \operatorname{Sci} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:
                          Hb\operatorname{Inq} Ha: t_3t_4Hc\operatorname{Inq} Hd. t_1t_2\operatorname{Fit} \mathfrak{a}\cdot \wedge \cdot t_3t_4Hb\operatorname{Ani.Utr} t_3t_4\operatorname{Etc} *
3 14 2.
                          (t_1 < t_2 < t_3)
                          ^{t_2}Ha \operatorname{Ing} Hb \operatorname{Utr}: t_3 t_4 Hb \operatorname{Sci} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x^{t_4} 
                          ^{t_{i}}Hb\operatorname{Inq}Ha . Fal | Qia : - : t_{1} t_{2} Ani · ? x . t_{1} t_{2} Fit x.
                                \operatorname{Erg} \cdot - t_3 t_4 Hb \operatorname{Sci}^{t_6} I
                          Ha\operatorname{Inq} Hb' \longrightarrow \bigvee \operatorname{Etc} . \in \operatorname{Fal} : \bigvee x \cdot \mathsf{t_7} \, \mathsf{t_8} \, Ha\operatorname{Inq} Hb . \, \mathsf{t_1} \, \mathsf{t_2} \operatorname{Fit} x \, \mathsf{I}
                          Mod \forall \lceil t \cdot t' \cdot x \cdot y \rceil \cdot t' < t_5 \cdot \land \cdot \text{Etc} : \text{Ilo} \cdot \forall \lceil t \cdot t' \cdot x \cdot y \rceil \text{ Etc} :
                                \Rightarrow: \forall z \cdot t'' t''' Hb \operatorname{Ing} Haz \in \operatorname{Ver} I
                          ^{t_7}Ha \operatorname{Ing} Hb \cdot t_1 t_2 \operatorname{Fit} \mathfrak{a}^{t_8} *
```

```
3 14 3. (t_1 < t_2 < t_3 < t_4 < t_5)

* {}^{t_1}Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} : t_5 t_6 Hb \operatorname{Sci} \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x^{t_4} : Hb \operatorname{Inq} Ha \operatorname{Ver} : Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} : t_1 t_2 Hb \operatorname{Ani} \operatorname{Etc} : Hb \operatorname{Inq} Ha \cdot \operatorname{Fal} : \operatorname{Sed} \cdot t_3 t_4 Hc \operatorname{Inq} Ha \cdot t_1 t_2 \operatorname{Fit} a : Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} : t_3 t_4 Hc \operatorname{Sci} \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x : {}^{t_7}Hb \operatorname{Inq} Ha : - \cdot t_7 t_8 Hb \operatorname{Sci} \operatorname{Cur} \cdot \operatorname{Tan} \cdot t_9 t_{10} Hb \operatorname{Sci} \operatorname{Cur}^{t_8} : Hb \operatorname{Inq} Hc \cdot \operatorname{Utr} : t_1 t_2 Hc \operatorname{Ani} \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x : Hc \operatorname{Inq} Hb \operatorname{Ver} : Hb \operatorname{Inq} Ha \cdot t_1 t_2 Hc \operatorname{Ani} \operatorname{Etc} *
```

```
3 14 4. The talk 3 13 8 will be repeated. Then we continue.
                    * Ha \operatorname{Inq} Hc : ? \operatorname{1um} x \cdot x \in \operatorname{Pri} A \cdot x > 10^{1001}
                        {}^{t_i}Hc \operatorname{Inq} Ha \ 1000001001^{t_i} \ I
                        Ha Inq Hc. Cur Scil
                        Hc \operatorname{Inq} Ha : t_3 t_4 Hc \operatorname{Ani} \cdot ? y \cdot t_3 t_4 Hb \operatorname{Inq} Ha y
                        Ha\operatorname{Inq} Hc : \operatorname{Mal} : t_5 t_6 Hb \operatorname{Sei} \cdot ? u \cdot t_3 t_4 \operatorname{Fit} u
                              \operatorname{Tan} : \neg : t_5 t_6 Hb \operatorname{Sci} : \operatorname{Utr} :
                              \lor \ulcorner t \cdot t' \urcorner : t < t' < t_3 \cdot \land \because tt' Ha \operatorname{Inq} Hb \cdot ? \operatorname{1um} x \operatorname{Etc} \mathsf{I}
                        Ha \operatorname{Ing} Ha \wedge : t_1 t_2 Hc \operatorname{Ani} \cdot ? u \cdot t_1 t_2 \operatorname{Fit} u
                        Ha \operatorname{Inq} Hc : \operatorname{Mal} : - t_5 t_6 Hc \operatorname{Sci} : \operatorname{Utr} u : t_3 t_4 \operatorname{Fit} u : \in \operatorname{Ben} I
                        Hc \operatorname{Inq} Ha \land : t_2 t_3 Hc \operatorname{Ani} ? u \cdot t_2 t_3 \operatorname{Fit} u : \land : \ u \cdot t_2 t_3 \operatorname{Fit} u \in \operatorname{Ben} I
                        Ha Inq Hc Ben #
3 14 5.
                        (t_1 < t_2 < t_3)
                   # Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 x \operatorname{Inq} Hb \mathfrak{p}
                        ^{\mathsf{t_3}}Hb\operatorname{Inq} Haı. — . \mathsf{t_3}\,\mathsf{t_4}\,Hb\operatorname{Sei}ı
                              Qia t_1 t_2 Hb Ani y \cdot \forall x \cdot t_1 t_2 x \text{ Inq } Hb y
                             Tan: \neg Ani: ? x \cdot \lor y \cdot t_1 t_2 x \operatorname{Inq} Hb y^{t_i} *
                    # Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x 
                        Hb\operatorname{Inq} Ha = \operatorname{Sei} \operatorname{Qqm} \operatorname{t}_1 \operatorname{t}_2 Hb\operatorname{Ani} : \operatorname{Utr} \cdot \vee x \cdot \operatorname{t}_1 \operatorname{t}_2 \operatorname{Fit} x:
                              \operatorname{Tan} - \operatorname{Ani} \cdot x \cdot t_1 t_2 \operatorname{Fit} x +
3 14 6. tha Ing Ha lum. Ha Ing Ha 10um.
                             Ha Inq Ha 11um Ha Inq Ha 100um to
                        Hb \operatorname{Ing} Hc t_1 t_2 Ha \operatorname{Ani} : ?x \cdot \operatorname{Fre} x = a \cdot \wedge \cdot t_1 \operatorname{Ant} x \operatorname{Ant} t_2
                        Hc \operatorname{Inq} Hb \cdot \operatorname{Cur} \operatorname{Sci}
                        Hb \operatorname{Inq} Hc : \operatorname{Qia} \cdot t_1 t_2 Ha \operatorname{Enu} \cdot ? \operatorname{Etc} *
3 14 7. * t1p.q.rt2:
                        Ha \operatorname{Inq} Ha : \mathfrak{p.q.r.p.q.r.p.q.r^{t_0}}
                        Hb \operatorname{Inq} Hc : t_1 t_2 Ha \operatorname{Ani} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x
                        Hc \operatorname{Inq} Hb \operatorname{Cur} \operatorname{Sei} :
                        Hb \operatorname{Inq} Hc : \operatorname{Qia} : t_2 t_3 Ha \operatorname{Sei} \cdot
                              ?x.t_1t_2 Fit x: \land `-: \lor \ulcorner t.t'.x.y \urcorner \cdot tt'y Inq Ha.t_1t_2 Fit x #
3 14 8. # Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 Hc \operatorname{Inq} Hdx
                        Hb \operatorname{Ing} Ha \cdot t_1 t_2 Hc \operatorname{Ing} Hd \operatorname{Fal}:
                        Ha \operatorname{Ing} Hb Utr: t_1 t_2 Hb \operatorname{Ani} \cdot \operatorname{Utr} \cdot t_1 t_2 Hc \operatorname{Ing} Hd \operatorname{Fal}:
                        Hb \operatorname{Ing} Ha - t_1 t_2 Hb \operatorname{Ani} \operatorname{Etc}
                        t_0 Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} \vee \mathsf{r} t \cdot t' \cdot x = t_2 < t < t' < t_3.
                              \wedge \cdot t t' x \operatorname{Inq} Hb \cdot t_1 t_2 Hc \operatorname{Inq} Hd \operatorname{Fal}
                        Hb \operatorname{Inq} Ha : \neg \cdot \vee \lceil t \cdot t' \cdot x \rceil \operatorname{Ete} :
```

Ha Inq Hb. Cur Sei!

```
Hb \operatorname{Inq} Ha Qia \cdot to t_1 Hd \operatorname{Inq} Hc \cdot 10 \times 10 = 101:

\operatorname{Erg.} t_1 t_2 Hc \operatorname{Inq} Hd \operatorname{Fal}:

\operatorname{Erg.} t_1 t_2 Hc \operatorname{Inq} Hd \operatorname{Fal} : tt' Hd \operatorname{Inq} Hc y \cdot \wedge \cdot y \in \operatorname{Fal} \cdot \rightarrow \cdot \vee t'' \cdot t' t'' Hc \operatorname{Inq} Hd \operatorname{Fal} *
```

Knowledge about events by inference.

```
3 14 9. (t_0 < t_1 < t_2 < t_3 < t_4)

# Ha \operatorname{Inq} Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x

Hb \operatorname{Inq} Ha : \lor y \cdot t_1 t_2 Hd \operatorname{Inq} Hc y \cdot \land \cdot y \in \operatorname{Fal} :

Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} \cdot t_1 t_2 Hb \operatorname{Ani} \operatorname{Etc} :

Hb \operatorname{Inq} Ha \operatorname{Fal} :

^t \cdot Ha \operatorname{Inq} Hb : \operatorname{Utr} : \lor \vdash t' \cdot t'' \cdot x \cdot u \cdot v \vdash : t_2 < t' < t'' < t_4 .

^t \cdot t' t'' Hb \operatorname{Ani} : \operatorname{Utr} \cdot u \operatorname{Inq} v \cdot t_1 t_2 \operatorname{Fit} x :

Hb \operatorname{Inq} Ha : \operatorname{Fal} : \operatorname{Sed} \cdot t_2 t_3 Hb \operatorname{Ani} \cdot \operatorname{Utr} \cdot t_2 t_3 Hc \operatorname{Inq} Hd \operatorname{Fal} :

^t \cdot t_2 t_3 Hc \operatorname{Inq} Hd \operatorname{Fal} : \operatorname{Erg} : \lor y \cdot t_1 t_2 Hd \operatorname{Inq} Hc y \cdot \land \cdot y \in \operatorname{Fal} :

Hc \operatorname{Inq} Hb : \operatorname{Fal} : \operatorname{Sed} : t_0 t_1 Hc \operatorname{Inq} Hd : \operatorname{Utr} \cdot 10^{1001} + 1 \cdot \in \operatorname{Pri} :

^t \cdot t_1 t_2 Hd \operatorname{Inq} Hc \operatorname{Fal} :
```

3 15 1. We repeat 3 07 6, and then we continue:

```
* Ha Inq Hd.? = 11010 × 1011:

Hd Inq Ha 100011110:

Ha Inq Hd 'Utr: V \( \text{t'} \) \cdot tt' Hd Ani.? = 11010 × 1011:

Hd Inq Ha 'Fal: tt' Hd Ani.? = 11010 × 1011 \( \delta \) Inq Hd 'Utr: V \( \text{t'} \) \cdot tt' Hd Cpu.? = Etc:

Hd Inq Ha Fal:

Ha Inq Hd. Cur Sei:

Hd Inq Ha 't_1 t_2 Hd Ani:

Utr \( \delta \) Hd Inq Ha.11010 × 1011 = 10001110 *
```

There is no perception of mathematical truth.

```
3 15 2. * Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x .

Hb \operatorname{Inq} Ha \cdot t_1 t_2 \operatorname{Fit} \alpha .

Ha \operatorname{Inq} Hb \cdot ? \vdash t \cdot t' \vdash : tt' Hb \operatorname{Sci} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x :

\land \vdash \neg : \lor \vdash u \cdot u' \vdash : u < u' < t \cdot \land : uu' Hb \operatorname{Sci} \operatorname{Etc} :

Hb \operatorname{Inq} Ha : t_3 t_4 Hb \operatorname{Cog} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x :

Ha \operatorname{Inq} Hb \cdot ? = \operatorname{Cog} :

Hb \operatorname{Inq} Ha : \operatorname{Exg} \vdash \neg : tt' Hb \operatorname{Sci} \operatorname{Etc} :

\land \cdot t' t'' Hb \operatorname{Sci} \operatorname{Etc} : \rightarrow : t' t'' Hb \operatorname{Cog} \operatorname{Etc} :

\operatorname{Ini} \cdot Hb \operatorname{Sci} \operatorname{Etc} : = \cdot \operatorname{Ini} Hb \operatorname{Cog} \operatorname{Etc} :
```

The word written Cog (fL cognoscit = recognizes) means "comes to know", "finds out". Syntactically it behaves like "Sci'.

```
3 15 3. # Ha Inq Hb · ? x · t<sub>1</sub> t<sub>2</sub> Fit x :

Hb Inq Ha a :

Ha Inq Hb · Cur Sci :

Hb Inq Ha · t<sub>3</sub> t<sub>4</sub> Hb Cog Etc :

Ha Inq Hb · Utr : t<sub>3</sub> t<sub>4</sub> Hb Cog · Etc · Mod · t<sub>3</sub> t<sub>4</sub> Ani :

Hb Inq Ha Fal :

Ha Inq Hb · Utr : Etc · Mod · \( \nu \) · t<sub>3</sub> t<sub>4</sub> x Inq Hb Etc :

Hb Inq Ha Fal :

Ha Inq Hb · ? u : t<sub>3</sub> t<sub>4</sub> Hb Cog · Etc · Mod u :

Hb Inq Ha : Mod · Hb Ani : Utr · Hc Inq Hd · t<sub>1</sub> t<sub>2</sub> Fit a *
```

3 15 4. We repeat 3 07 6, and then we continue:

```
* Ha Inq Hd.?=11010 × 1011:

Hd Inq Ha 100011110:

Ha Inq Hd. Cur Sei:

Hd Inq Ha.t<sub>1</sub>t<sub>2</sub>Hd Cog Etc:

Ha Inq Hd Utr Mod Cpu:

Hd Inq Ha 'Fal: Mod · Hb Cpu. A. Hd Ani *
```

3 15 5. The greater part of 3 05 1 is repeated, till the moment t₂. Then one continues:

```
* Ha \operatorname{Inq} Hc : \operatorname{Utr} : \land n \cdot n \in \operatorname{Num} . \rightarrow : \lor p \cdot p \in \operatorname{Pri} . \land . p > n \mid
Hc \operatorname{Inq} Ha \operatorname{Ver} \mid
Ha \operatorname{Inq} Hc \operatorname{Cur} \operatorname{Sci} \mid
Hc \operatorname{Inq} Ha . t_1 t_2 Hc \operatorname{Cog} \mid
Ha \operatorname{Inq} Hc : ? \lor z \cdot t_1 t_2 Hc \operatorname{Cog} . \operatorname{Etc} . \operatorname{Mod} z \mid
Hc \operatorname{Inq} Ha . t_1 t_2 Hc \operatorname{Ani} . \operatorname{Utr} t_1 t_2 Hb \operatorname{Dem} \cdot \land . t_1 t_2 Hb \operatorname{Dem} \cdot
\operatorname{Erg} . \operatorname{Ver} \cdot \land \vdash t . t' \vdash : t_2 < t < t' . \rightarrow \cdot tt' \operatorname{Hc} \operatorname{Sci} . \operatorname{Utr} \operatorname{Ver} \mid
Ha \operatorname{Inq} Hc : ? = : \operatorname{Dem}^{io} : \land n \cdot n \in \operatorname{Num} . \rightarrow : \lor p \cdot p \in \operatorname{Pri} . \land . p > n \mid
t_3 \operatorname{Hc} \operatorname{Inq} Ha : \neg : \lor v \cdot t_3 t_4 \operatorname{Sci} : \operatorname{Utr} \cdot v \in . \operatorname{Dem}^{io} \operatorname{Etc} :
\operatorname{Qia} : \lor \vdash u \cdot v \cdot t_3 t_4 \operatorname{Sci} : \operatorname{Utr} \cdot u \in . \operatorname{Anl}^{io} v \cdot \land
t_3 t_4 \operatorname{Hc} \operatorname{Sci} : \operatorname{Utr} \cdot u \in . \operatorname{Anl}^{io} v^{t_4} =
```

Ha Inq Hb:
$$? = {}^{\checkmark} {}^{\lnot}a \cdot b \cdot c^{\lnot} \cdot {}^{\bigstar}x$$
 Etc!
Hb Inq Ha ${}^{\diamond}a \neq 0 \cdot {}^{\rightarrow} \cdot y^2 = b^{10} - 100 ac \cdot {}^{\rightarrow} \cdot x = (y - b)/10 a : (-y - b)/10 a : a = 0 \cdot {}^{\land}b \neq 0 \cdot {}^{\rightarrow} \cdot x = -c/b : a = b = 0 \cdot {}^{\land}c \neq 0 \cdot {}^{\rightarrow} \cdot {}^{\rightarrow} \cdot {}^{\backprime}x$ Etc: $a = b = c = 0 \cdot {}^{\rightarrow} \cdot {}^{\land}x$ Etc:

Note that with respect to oblique speech 'Sci' (and also 'Ani' and so on) behaves as 'Inq'. (See 3 01 5, and 2 04 0-2 04 6.) Ha asked whether Hb can solve the quadratic equation $ax^2 + bx + c = 0$ for every value of a, b, c (quantification outside the working sphere of 'Sci'). Hb, however, stresses that he can solve even the general quadratic equation.

The word written Eti (fL etiam=also) means "even".

3 16 1. Since we interpret

$$Ha \operatorname{Sci} \cdot ?x \cdot t_1 t_2 Hb \operatorname{Inq} Hc$$

as though Ha is answering the question? x Etc, it cannot happen under this formulation that Ha knows what was said, while ignorant of who said it, or to whom it was said, or when it was said. If we wish to express such an occurrence, we can avail ourselves of the pattern of $3\ 14\ 4-3\ 14\ 5$:

```
* Ha \operatorname{Inq} Hb \cdot ? x \cdot t_1 t_2 Hb \operatorname{Inq} Hd x  

^{t_1}Hb \operatorname{Inq} Ha \cdot t_3 t_4 Hb \operatorname{Sei} : ? x \cdot \vee y \cdot t_1 t_2 y \operatorname{Inq} Hd x  

\operatorname{Tan} : \neg \cdot \operatorname{Sei} : ? y \cdot \vee x \cdot t_1 t_2 y \operatorname{Inq} Hd x^{t_1}  

Ha \operatorname{Inq} Ha \cdot t_3 t_4 Hb \operatorname{Sei} : ? x \cdot t_1 t_2 : \forall y \cdot t_3 t_4 Ha \operatorname{Sei} : ? = y : \operatorname{Inq} Hc x  *
```

At some stage of a case under trial the evidence may look like that reported by Ha in his last sentence or something even more intricate:

A murder has been committed, and the detective has traced a person, who knows the identity of the victim (though he does not know that this person has been murdered); the detective will also have found somebody who knows the author of the crime (though he does not know that a crime has been committed); he has discovered the possessor of the weapon (who is as little in the secret); and so on. The pattern of such evidence looks like:

```
* \forall u \cdot t_1 t_2 Ha \operatorname{Sci} \cdot ? = u : \forall x \cdot t_1 t_2 Hb \operatorname{Sci} \cdot ? = x :
\operatorname{Inq} : \forall y \cdot t_1 t_2 Hc \operatorname{Sci} \cdot ? = y : \forall p \cdot t_1 t_2 Hd \operatorname{Sci} \cdot ? = p *
```

At a time u known to Ha a person x known to Hb said to a person y known to Hc a saying p known to Hd. The use of the definite article can be avoided by saying:

- * $\vee \lceil u \cdot x \cdot y \cdot p \rceil$: $ux \operatorname{Inq} y p \cdot \wedge \cdot t_1 t_2 Ha \operatorname{Sci} \cdot ? = u \cdot \wedge \cdot t_1 t_2 Hb \operatorname{Sci} \cdot ? = x \cdot \wedge \cdot t_1 t_2 Hc \operatorname{Sci} \cdot ? = y \cdot \wedge \cdot t_1 t_2 Hd \operatorname{Sci} \cdot ? = p$ *
- 3 17 0. In 3 01 5 we remarked that in the affirmative mode the direct and oblique speech are semantically equivalent. In the interrogative mode, however, the difference is enormous.

For instance, in French:

Il me disait: "où sont les neiges d'antan?"

Il me disait où sont les neiges d'antan.

In German:

Ich sagte ihm: "wer kommt?"

Ich sagte ihm, wer kommt.

In English this paradox is less striking on account of the use of the verb "to tell" in the second case:

I said to him: "who is coming?"

I told him who is coming.

It is clear that the second mode is not truly interrogative.

How shall we treat pseudo-interrogative oblique speech in Lincos? Of course we cannot translate

"Ha tells Hb what has happened from t_1 to t_2 " by

 $Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$.

because the latter sentence means

"Ha says to Hb: "what has happened from t1 to t2?"

The translation

 $Ha \operatorname{Inq} Hb \cdot {}^{\vee}x \cdot t_1 t_2 \operatorname{Fit} x$

would be just as impossible (or at least unsatisfactory) because this means (or could mean)

"Ha says to Hb: $\forall x.t_1 t_2 \text{ Fit } x$ ".

If we were to translate not the event

"Ha tells Hb what has happened from t₁ to t₂",

but the report

"From t₃ to t₄ Ha told Hb what has happened from t₁ to t₂", we could use the syntactical device of outside binding. Then we should translate:

 $t_1 t_2 \operatorname{Fit} x . \rightarrow . t_3 t_4 Ha \operatorname{Inq} Hb x$.

But this way out is blocked in the present situation.

We need a new device for pseudo-interrogative oblique speech. This will be the exclamation mark, borrowed from the alternative method alluded to in 3 13 0.

"Ha tells Hb what has happened from t_1 to t_2 "

could be translated by

 $Ha \operatorname{Inq} Hb \cdot ! ? x \cdot t_1 t_2 \operatorname{Fit} x$.

But this is still unsatisfactory. The receiver might be tempted to believe that

! ? x . t1 t2 Fit x

was something that could be said. Moreover, we could not prevent the actors from saying it.

I have therefore decided to write

 $Ha \operatorname{Inq} ! Hb ? x \cdot t_1 t_2 \operatorname{Fit} x$

with the exclamation mark after the 'Inq'.

3 17 1. $\# Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$:

 $Hb \operatorname{Inq} Ha \cdot \forall x \cdot t_1 t_2 \operatorname{Fit} x$:

 $Ha \operatorname{Ing} Hb : \operatorname{Qgm} \operatorname{Ver} . \operatorname{Tan} \operatorname{Mal} t_1 t_2 \operatorname{Fit} \cdot x : t_1 t_2 \operatorname{Fit} x : \in \operatorname{Mal} :$

 $Hb \operatorname{Inq} Ha a$:

Ha Inq Hb Ben #

3 17 2. $(t_1 < t_2 < t_3)$

 $Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$

taHb Inq Ha ata

 $Hc \operatorname{Inq} Hc: t_3 t_4 Hb \operatorname{Inq} ! Ha: ?x \cdot t_1 t_2 \operatorname{Fit} x$

3 17 3. $(t_1 < t_2 < t_3)$

$Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$

 $t_3Hb \operatorname{Inq} Ha \cdot \forall x \cdot t_1 t_2 \operatorname{Fit} x^{t_4}$

 $Ha\operatorname{Inq} Hb\operatorname{Mal}$ ı

 $Hc \operatorname{Inq} Hc \colon t_3 t_4 Hb \operatorname{Inq} Ha \cdot \Upsilon x \cdot t_1 t_2 \operatorname{Fit} x \colon$

Sed $-: t_3 t_4 Hb \operatorname{Inq} ! Ha \cdot ? x \cdot t_1 t_2 Fit x #$

3 17 4. # t₁ p t₂.

 $Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} \mathfrak{p}^{t_3}$

 $Hc \operatorname{Inq} Hc: t_2 t_3 Ha \operatorname{Inq} ! Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x =$

3 17 5. # t₁ p t₂:

Ha Inq Hb p: ts

 $Hc \operatorname{Inq} Hc - : t_2 t_3 Ha \operatorname{Inq} ! Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x +$

```
3 17 6. * ^{t_1} p ^{t_2}:

Ha \operatorname{Inq} ! Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x^{t_2}:

Hc \operatorname{Inq} Hd \cdot ? y \cdot t_2 t_3 Ha \operatorname{Inq} Hb y:

Hd \operatorname{Inq} Hc \cdot t_2 t_3 Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} p *
```

3 18 1. The (intentional) misprint in the following text is to reflect a slip of the speaker's tongue.

The word written Itg (fL intelligit=understands) means "understands". 'Inq Mod Mal (Mod Ben)' is an abbreviation of ' $\vee \cdot u \dots$ Inq Mod $u \cdot \wedge \cdot u \in$ Mal (Ben)'. This may be stated by a behaviour rule.

```
3 18 2. * {}^{t_1}Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x \overset{t_1}{:} t_2

Hb \operatorname{Inq} Ha \cdot \neg_{1} \operatorname{Itg} \overset{\cdot}{:}

Ha \operatorname{Inq} Hb : t_3 t_4 Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x \overset{\cdot}{:}

Hb \operatorname{Inq} Ha : \mathfrak{a} \cdot t_3 t_4 Hb \operatorname{Ani} \cdot ?y \cdot t_3 t_4 \operatorname{Fit} y :

\wedge : t_3 t_4 Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Ftt} \overset{\cdot}{:}

Hc \operatorname{Inq} Hb : t_3 t_4 Ha \operatorname{Inq} Hb \operatorname{Etc} \cdot \operatorname{Mod} \operatorname{Ben} \cdot

\operatorname{Tan} \cdot t_3 t_4 Hb \operatorname{Ani} \cdot ? \operatorname{Etc} \cdot \operatorname{Mod} \operatorname{Mal} \overset{*}{:}
```

```
3 18 3. * {}^{t_1}Ha \operatorname{Inq} Hb \cdot \operatorname{Utr.100001} \in \operatorname{Pri:} {}^{t_2}
Hb \operatorname{Inq} Hc \operatorname{Fal:}
{}^{t_3}Hc \operatorname{Inq} Hb \cdot \neg : t_3 t_4 Hc \operatorname{Itg} \cdot ? x \cdot t_2 t_3 Hb \operatorname{Inq} Hc x \cdot {}^{t_4}
Hb \operatorname{Inq} Hc \cdot \neg : \forall x \cdot t_2 t_3 Hb \operatorname{Inq} Hc x : \operatorname{Sed} \cdot t_2 t_3 Hb \operatorname{Inq} Ha \operatorname{Fal:}
\operatorname{Qia:} t_1 t_2 Ha \operatorname{Inq} Hb \cdot \operatorname{Utr.100001} \in \operatorname{Pri:}
Ha \operatorname{Inq} Hb \cdot \neg \cdot t_2 t_3 Hb \operatorname{Inq} Ha \operatorname{Fal:}
Hd \operatorname{Inq} Ha \circ Hb \circ Hc \cdot t_2 t_3 Hb \operatorname{Inq} Hc \operatorname{Fal.}
\land : t_2 t_3 Hb \operatorname{Inq} Ha \operatorname{Fal.Mod Mal} \cdot
\land : t_2 t_3 Hb \operatorname{Inq} Ha \operatorname{Fal.Mod } Hb \operatorname{Inq} Hc \operatorname{Fal.}
```

```
3 18 4. \# t_1 Ha \text{ Inq } Hb \cdot \text{Utr } 100001 \in \text{Pri} : t_2
                        Hb Inq Ha Fal:
                        t_2Hc \operatorname{Ing} Hb \longrightarrow t_3 t_4 Hc \operatorname{Itg} ? x \cdot t_2 t_3 Hb \operatorname{Ing} Hc x \cdot t_4
                        Hb \operatorname{Inq} Hc - \vee x \cdot t_2 t_3 Hb \operatorname{Inq} Hc x : \operatorname{Sed} \cdot t_2 t_3 Hb \operatorname{Inq} Ha \operatorname{Fal}:
                        Hd \operatorname{Ing} Hb \sim Hc \cdot t_2 t_3 Hc \operatorname{Ani}. Etc. Mod Mal *
3 18 5.
                        (t_1 < t_2 < t_3)
                    * t_a Ha \operatorname{Ing} Hb a^{t_a}:
                         Hb \operatorname{Inq} Ha \cdot - \operatorname{Itg}:
                        ^{t_a}Ha \operatorname{Inq} Hb \cdot ^{\vee}x \cdot t_1 t_2 \operatorname{Fit} x^{t_a}:
                         Hb \operatorname{Inq} Ha . \neg \operatorname{Itg}:
                         Ha \operatorname{Ing} Hb \cdot t_3 t_4 Ha \operatorname{Ing} ! Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:
                         Hb \operatorname{Inq} Ha \cdot \operatorname{Itg} \cdot t_1 t_2 \operatorname{Fit} a:
                         Hc \operatorname{Inq} Ha - Hb \cdot t_3 t_4 - t_5 t_6 \cdot Ha \operatorname{Inq} \cdot \operatorname{Mod} \operatorname{Mal} +
3 18 6.
                       (t_1 < t_2 < t_3)
                   * t_{\bullet} Ha \operatorname{Inq} Hb \cdot Hc^{t_{\bullet}}
                        Hb \operatorname{Inq} Ha . - \operatorname{Itg}:
                        Ha \operatorname{Ing} Hb t_3 t_4 Ha \operatorname{Ing} ! Hb : ?x \cdot t_1 t_2 x \operatorname{Ing} Ha \cdot 10 \times 10 = 100.
                        Hb \operatorname{Ing} Ha : \operatorname{Itg} \cdot t_1 t_2 Hc \operatorname{Ing} Ha : 10 \times 10 = 100 *
3 18 7.
                        (t_1 < t_2 < t_3)
                    * t_a Ha \operatorname{Ing} Hb - t_a
                        Hb \operatorname{Inq} Ha . \neg \operatorname{Itg}
                        Ha \operatorname{Ing} Hb : t_3 t_4 Ha \operatorname{Ing} ! Hb \cdot \operatorname{Utr} \cdot 100001 \in \operatorname{Pri}^*
                        Hb \operatorname{Ing} Ha : \operatorname{Itg} \cdot - 100001 \in \operatorname{Pri} *
3 18 8.
                        (t_1 < t_2 < t_3)
                   * t<sub>a</sub>Ha Inq Hb · Qia . 11 Div 100001 * t<sub>a</sub>
                        Hb \operatorname{Inq} Ha - \operatorname{Itg}
                        Ha \operatorname{Ing} Hb : t_3 t_4 Ha \operatorname{Ing} ! Hb \cdot \operatorname{Cur} . 100001 \notin \operatorname{Pri}
                        Hb Inq Ha: 11 Div 100001 · Erg. 100001 ∉ Pri *
3 18 9. \# Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} \cdot Hc \operatorname{Inq} Hb \operatorname{Pri}^*
                        Hb \operatorname{Inq} Ha - \operatorname{Itg} : \operatorname{Qia} \cdot Hc \operatorname{Inq} Hb \operatorname{Pri} \cdot \notin \operatorname{Prp}
                        Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot \vee \operatorname{ft'} \cdot \operatorname{tt'} Hc \operatorname{Inq} Hb \operatorname{Pri} \cdot
                        Hb \operatorname{Inq} Ha \cdot \operatorname{Itg.Tan.} \longrightarrow \operatorname{Sci} *
                   * Ha \operatorname{Inq} Hb \cdot t_1 t_2 Hc \operatorname{Inq} \operatorname{Pri} Hb
                        Hb \operatorname{Inq} Ha . \neg \operatorname{Itg}
                        Ha \operatorname{Inq} Hb \cdot t_1 t_2 Hc \operatorname{Inq} Hb \operatorname{Pri} \cdot
                        Hb \operatorname{Inq} Ha \cdot \operatorname{Itg} *
                    # Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot \forall x \cdot x \in \operatorname{Div}
                        Hb \operatorname{Inq} Ha : \neg \operatorname{Itg} : \operatorname{Qia} : \neg : \forall x \cdot x \in \operatorname{Div} : \in \operatorname{Prp}
                              \land \lor x \cdot x \in \text{Div} \cdot \notin \text{Prp} \land \land \neg \cdot \lor x \cdot x \in \text{Div} : \notin \text{Prp}
```

These are a few examples for the use of '- Itg' as a means of rejecting meaningless speech.

3 19 1. The time-signal in the next program texts will have a length of little more than 111,01 seconds and a frequency of about 10¹¹¹¹¹ Hertz.

```
# Ha \operatorname{Inq} Hb \cdot \operatorname{Ecc}^{\mathfrak{t}_1} - t_2
    Ha\operatorname{Inq} Hb !=:\operatorname{Fre} \cdot {}^{\vee}x \cdot \operatorname{t}_1\operatorname{t}_2\operatorname{Fit} x
    Hb \operatorname{Ing} Ha " > \cdot \operatorname{Sec}^{-1} \cdot 10^{10000} ; \Leftrightarrow : < \cdot \operatorname{Sec}^{-1} \cdot 10^{100000} "
         Hc \operatorname{Inq} Ha : \operatorname{Cea} \cdot \operatorname{Sec}^{-1} \cdot 10^{11111}
   Ha \operatorname{Ing} Hb - Hc \cdot \operatorname{Err} 10^{11111} < \operatorname{Err} 10^{11100}
   Ha \operatorname{Inq} Hb ? =: \operatorname{Dur} \cdot {}^{\vee}x \cdot \operatorname{t}_1 \operatorname{t}_2 \operatorname{Fit} x :
    Hb \operatorname{Inq} Ha \cdot \operatorname{Sec} 111:
    Ha \operatorname{Inq} Hb \operatorname{Mal}:
    Hb \operatorname{Inq} Ha \cdot \operatorname{Cea. Sec} 111:
   Ha \operatorname{Inq} Hb: {}^{?} a \operatorname{Dur} x \operatorname{Etc} = \operatorname{Cca} a
         \cdot \wedge : \operatorname{Err} a : < \cdot \operatorname{Err} \cdot \operatorname{Sec} 111 :
   Hb \operatorname{Inq} Ha \cdot \operatorname{Sec} 111.01:
   Ha \operatorname{Inq} Hb Ben: Dur x \operatorname{Etc}. Pau > . Sec 111,01
         \wedge \cdot Dur \Upsilon x \text{ Etc. Mul} < \cdot \text{Sec } 111.10:
         Err . Sec 111.01 · Mul < · Err . Sec 111 *
```

The word written Err has been known from 3 09 1. 'Err' is a function from the set of propositions to the (partially ordered) set of errors. We have used and shall often use this abbreviated construction. The complete sentences would be

```
# Err ' Fre ' x \cdot t_1 t_2 Fit x := \text{Cea} \cdot \text{Sec}^{-1} \cdot 10^{11111}' 
 < \text{Err} \cdot \text{Fre} \cdot \text{Y} x \cdot t_1 t_2 Fit x := \text{Cea} \cdot \text{Sec}^{-1} \cdot 10^{11100} #
```

and so on. These formulations should be linked together by a behaviour rule.

Cea (fL circa=about, nearly) means "about, approximately". Later on it will also be used in less numerical contexts, in the sense of "nearly".

The words written Mul (fL multo = much) and Pau (fL paulo = little) mean "much" and "little" respectively. They behave like adverbs added to '>' and '<'.

```
3 19 2. * Ha \operatorname{Inq} Hb^{??} x : x \in \operatorname{Rat.} \wedge \cdot x := \operatorname{Cea.} \sqrt{10} :

Hb \operatorname{Inq} Ha 1010/111 :

Ha \operatorname{Inq} Hb^{??} x : x \in \operatorname{Rat.} \wedge \cdot x := \operatorname{Cea.} \sqrt{10} \cdot

\wedge \cdot \operatorname{Err} x . \operatorname{Mul} < \cdot \operatorname{Err} 1010/11 :

Hb \operatorname{Inq} Ha 111/101 :
```

```
Ha Inq Hb · Err 111/101 · Pau < . Err 1010/11 :

Hb Inq Ha 1100011/1000110 :

Ha Inq Hb Ben *
```

3 19 3. About 250 time-signals of frequency a and length h are given in a relatively short time, from t₁ to t₂.

```
# Ha \operatorname{Inq} Hb \operatorname{Ece}^{t_1, \dots, t_2} |

Ha \operatorname{Inq} Hb : ? = \operatorname{iCar}^*t \cdot t_1 \le t \le t_2 - \mathfrak{h}.

\wedge : \forall x \cdot t t + \mathfrak{h} \operatorname{Fit} x \cdot \wedge . \operatorname{Fre} x = \mathfrak{a}. \wedge . \operatorname{Dur} x = \mathfrak{h} |

Hb \operatorname{Inq} Ha \cdot \operatorname{Mul} > . 1 |

Ha \operatorname{Inq} Hb : ? \forall y \cdot \operatorname{Err} y \cdot \operatorname{Mul} < : \operatorname{Err} \cdot \operatorname{Mul} > . 1 |

Hb \operatorname{Inq} Ha \cdot \operatorname{Pau} < . 1000000000 = 0
```

3 19 4. \mathfrak{p} will be the metatextual substitute for some noise, \mathfrak{p}' that for a bad copy of \mathfrak{p} , \mathfrak{p}'' that for a better copy.

```
# Ha \operatorname{Inq} Hb \cdot \operatorname{Ecc} \cdot t_{1} \mathfrak{p}^{t_{2}} |

Ha \operatorname{Inq} Hb \cdot ? x \cdot t_{1} t_{2} \operatorname{Fit} x |

Hb \operatorname{Inq} Ha t_{2} \mathfrak{p}' t_{1} |

Ha \operatorname{Inq} Hb : \operatorname{Mal} : ? \vee y : y \cdot = \operatorname{Cea} \cdot \vee x \cdot t_{1} t_{2} \operatorname{Fit} x |

\wedge \cdot \operatorname{Err} y \cdot < : \operatorname{Err} \cdot \vee z \cdot t_{3} t_{4} \operatorname{Fit} z |

Hb \operatorname{Inq} Ha \mathfrak{p}'' \cdot :

Ha \operatorname{Inq} Hb \operatorname{Ben} \cdot :

Hc \operatorname{Inq} Hd \cdot \mathfrak{p} \operatorname{Ant} \mathfrak{p}' \cdot :

Ha \operatorname{Inq} Hc : ? x \cdot \mathfrak{p} \cdot x \operatorname{Ant} \cdot \mathfrak{p}' \cdot :

Hc \operatorname{Inq} Ha \operatorname{Pau} \cdot :

Ha \operatorname{Inq} Hc : ? \vee y : \operatorname{Err} : \mathfrak{p} \cdot \operatorname{Cea} y \cdot \operatorname{Ant} \cdot \mathfrak{p}' \cdot < \cdot \operatorname{Err} : \mathfrak{p} \cdot \operatorname{Pau} \operatorname{Ant} \cdot \mathfrak{p}' \cdot :

Hc \operatorname{Inq} Hd \cdot \mathfrak{p} \cdot (t_{3} - t_{2}) \operatorname{Ant} \cdot \mathfrak{p}' \cdot :

Hd \operatorname{Inq} Hc \operatorname{Ben} \cdot #
```

```
3 19 5. * Ha \operatorname{Inq} Hb \cdot ?x \cdot x^{10} - 1110 x - 111111001111 = 0:

Hb \operatorname{Inq} Ha \cdot x = 1000111 \sim -111001:

Ha \operatorname{Inq} Hb \operatorname{Cur}:

Hb \operatorname{Inq} Ha \cdot \operatorname{Qia} \operatorname{Sei}:

Ha \operatorname{Inq} Hb \cdot ?u \cdot \operatorname{Cog} \operatorname{Mod} u:

Hb \operatorname{Inq} Ha \cdot Hb \operatorname{Cpu} \cdot ?x \cdot x^{10} - 1110 x - 111111001111 = 0:

\operatorname{Mul} \operatorname{Ant} \cdot Ha \operatorname{Inq} Hb \operatorname{Ete}:

Ha \operatorname{Inq} Hb : ? \mathsf{'t} \cdot \operatorname{Err} t \cdot < \cdot \operatorname{Err} \operatorname{Mul}:

Hb \operatorname{Inq} Ha \cdot \operatorname{Sec} 10^{11001}:

Ha \operatorname{Inq} Hb \operatorname{Ben} *
```

3 20 1. The first time Ha pronounces the word written Nnc (fL nunc = now) meaning "now", a time-signal of frequency a is given.

'Nnc Fit \mathfrak{p} ' is a proposition the truth-value of which depends on time, whereas the value of ' t_1 t_2 Fit \mathfrak{p} ' does not depend on time. Extensionally they cannot be equivalent. In an intentional context they can.

Words such as "now", "here", "I", "this", "that" can be avoided. But for this economy we should have to pay clearly in the form of clumsy circumlocutions. We have not introduced a word meaning "I". Up to now, when speaking about himself, everybody used his proper name. Of course there may be situations where the speaker wants to conceal his name though speaking about himself. Then he can use bound variables, e.g.

$t_1 t_2$: $\forall x \cdot \forall p$. Nne $x \operatorname{Inq} Ha p$: $\operatorname{Inq} Ha \mathfrak{q}$ #,

but if a word meaning "I" is available, he can simplify his speech considerably. As long as we are not interested in situations like this, we can dispense with "I". Likewise we have not introduced words meaning "this" and "that" though a more sophisticated analysis may reveal their occurrence in our texts: often the 'Etc' is used as if it were a demonstrative pronoun. It would be worthwhile considering whether this use should not be formalized by the introduction of formal demonstrative pronouns.

In any case it would not be wise to avoid words like "now" and "here". We will now deal with "now", and we shall see that its use is a real advantage.

"Now", like "I" and the other words mentioned, conceals a variable. So we must agree upon the place where this variable is considered to be bound, when Nnc occurs in a report such as

 $* Ha \operatorname{Inq} Hb : t_1 t_2 Hc \operatorname{Inq} Hb \cdot \operatorname{Nne} Hd \operatorname{Sei} ? \operatorname{Ete} *$

In all occurrences of hidden variables we shall observe the rule of binding inside. The 'Nnc' of the example is (approximately) the moment of Hc's pronouncing Nnc (i.e. $t_1 t_2$) not that of Ha's reporting on Hc's pronouncing Nnc. Likewise a word meaning "T", if available, would mean the direct user, not the person reporting about its use. Thus syntactically oblique speech with respect to hidden variables, which is a well-known feature of common languages, is unknown in Lincos.

3 20 2. Parts of our former programm will be repeated, with 'Nnc' instead of an explicit date.

3 20 3. $(t_1 < t_2 < t_3)$ * $t_3 Ha \operatorname{Inq} Hb : \neg : \operatorname{Nne} Ha \operatorname{Sei} \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x^{t_4} : Hb \operatorname{Inq} Ha \cdot t_1 t_2 \operatorname{Fit} \alpha : t_3 Ha \operatorname{Inq} Hb : \neg : \operatorname{Pau} \operatorname{Ant} \cdot \operatorname{Nne} \cdot Ha \operatorname{Sei} \cdot ? \operatorname{Ete} : \operatorname{Tan} \cdot \operatorname{Nne} Ha \operatorname{Sei} \cdot ? \operatorname{Ete} : \operatorname{Nne} Ha \operatorname{Cog} \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x : Hc \operatorname{Inq} Hc : t_3 t_4 Ha \operatorname{Inq} Hb : \neg : t_3 t_4 Ha \operatorname{Sei} \cdot ? \operatorname{Ete} : t_5 t_6 Ha \operatorname{Inq} Hb : \neg : t_3 t_4 \cdot \operatorname{Tan} t_5 t_6 \cdot Ha \operatorname{Sei} \cdot ? \operatorname{Ete} : \operatorname{Compare} 2039 \text{ on this use of } 'h \operatorname{Ant} \cdot y' \text{ for dating.}$

3 20 4. $(t_1 < t_2 < t_3)$ # ${}^{t_4}Ha \operatorname{Inq} Hb {}^{\bullet} - : t_3 t_4 Ha \operatorname{Sci} \cdot ? x . t_1 t_2 \operatorname{Fit} x^{t_4} |$ # ${}^{t_6}Ha \operatorname{Inq} Hb {}^{\bullet} - : t_3 t_4 Ha \operatorname{Sci} \cdot ? x . t_1 t_2 \operatorname{Fit} x^{t_4} |$ # ${}^{t_6}Ha \operatorname{Inq} Hb {}^{\bullet} - : t_3 t_4 . \operatorname{Tan} t_5 t_6 \cdot \operatorname{Ha} \operatorname{Sci} {}^{t_6} |$ # ${}^{t_6}Ha \operatorname{Inq} Hb {}^{\bullet} - : t_3 t_4 Hc \operatorname{Inq} Hb . - : \operatorname{Nnc} Ha \operatorname{Sci} {}^{\bullet} |$ # ${}^{t_6}Ha \operatorname{Inq} Hb {}^{\bullet} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} : \operatorname{Ha} \operatorname{Inq} Hb {}^{\bullet} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Sci} {}^{\bullet} |$ # ${}^{t_6}Ha \operatorname{Inq} Hb {}^{\bullet} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Sci} {}^{\bullet} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Sci} {}^{\bullet} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Sci} {}^{\bullet} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Cog} {}^{\#} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Cog} {}^{\#} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Cog} {}^{\#} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Cog} {}^{\#} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Cog} {}^{\#} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Cog} {}^{\#} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Cog} {}^{\#} - : \operatorname{Pau} \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Pau} \cdot \operatorname{Ant} . \operatorname{Nnc} \cdot \operatorname{Ant} \cdot \operatorname{$

3 20 5. # $Ha \operatorname{Inq} Hb$: ?. $1^{\operatorname{um}} x \cdot x \in \operatorname{Pri} \cdot \wedge \cdot x > 10^{1001}$! $Hb \operatorname{Inq} Ha$: $\neg \cdot \operatorname{Nne} Hb \operatorname{Sei}$. ? Etc: $Ha \operatorname{Inq} Ha \cdot \operatorname{Nne} Hb \operatorname{Pet}$. ? Etc: $Hb \operatorname{Inq} Ha \cdot \operatorname{10^{1001}} + 1001$! $Hb \operatorname{Inq} Hb$: Pau Ant. $\operatorname{Nne} \cdot Hb \operatorname{Pet}$. ? Etc:

3 20 6. # $Ha \operatorname{Inq} Hb$: ?. $1^{\operatorname{um}} x \cdot x \in \operatorname{Pri.} \wedge . x > 10^{1001}$: $Hb \operatorname{Inq} Ha \cdot 10^{1001} + 1001$: $Ha \operatorname{Inq} Hb \cdot \operatorname{Cur.} \operatorname{Nne} Hb \operatorname{Sei}$: $Hb \operatorname{Inq} Ha$: Qia : $\operatorname{Mul} \operatorname{Ant.} \operatorname{Nne} \cdot Hb \operatorname{Cpu}$: $h > 0 \cdot \rightarrow h \operatorname{Pst} \cdot \operatorname{Mul} \operatorname{Ant.} \operatorname{Nne} : Hb \operatorname{Sei}$ #

Clearly this is a syntactically new appearance of 'Mul Ant. Nnc' and such expressions. In 'Mul Ant. Nnc' Fit p' an existentially-bound variable is hidden in 'Mul Ant. Nnc'. In expressions like $h\operatorname{Pst}$ 'Mul Ant. Nnc: Fit q this existentially bound variable is used as an instance out of the set to which it has been bound. In other cases we have accounted for this use in common mathematical and everyday language by the word 'Put' (see 3 05 2). We could justify the new use of 'Mul Ant. Nnc' by a similar artifice, but we had perhaps better avoid subtleties of this kind. A behaviour rule stressing the new use could be added.

```
3 20 7. * p:

Ha Inq Hb ? x: Pau Ant. Nnc · Fit x:

Hb Inq Ha ' Pau Ant. Ha Inq Hb Etc · Fit p:

Pau Pst. p. Ha Inq Hb Etc *
```

3 21 1. We introduce some abbreviations:

* Pau Ant. Nnc·Etc: ↔ . PAN Etc '
Pau Pst. Nnc·Etc: ↔ . PPN Etc '
Mul Ant. Nnc·Etc: ↔ . MAN Etc '
Mul Pst. Nnc·Etc: ↔ MPN Etc *

3 21 2. * Unq Ant $.u \cdot \text{Fit } p : \leftrightarrow .p \text{ Ant } u$ Unq Pst $.u \cdot \text{Fit } p : \leftrightarrow .p \text{ Pst } u$ Unq Ant $.\text{Nnc} \cdot \text{Etc} : \leftrightarrow \text{UAN Etc}$ Unq Pst $.\text{Nnc} \cdot \text{Etc} : \leftrightarrow \text{UPN Etc}$ Unq Fit $p : \leftrightarrow . \forall u . u \text{ Fit } p *$

The word written Unq (fL unquam = ever) means "at some time".

3 21 3. # $Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$: ${}^{t_1}Hb \operatorname{Inq} Ha \cdot t_1 t_2 Hb \operatorname{Ani} \cdot ? \operatorname{Etc} \cdot \operatorname{Tan} : --- \cdot t_3 t_4 Hb \operatorname{Sei} \cdot ? \operatorname{Etc} \cdot {}^{t_4}$ $Ha \operatorname{Inq} Ha \cdot --- \cdot t_3 t_4 Hb \operatorname{Sei} ? \operatorname{Etc} \cdot \operatorname{Qqm} \cdot \operatorname{Unq} \operatorname{Ant} \cdot t_3 t_4 : Hb \operatorname{Sei} ? \operatorname{Etc} \cdot {}^{\bullet}$ $--- \cdot t_3 t_4 \cdot \operatorname{QUA} \cdot Hb \operatorname{Sei} \cdot ? \operatorname{Etc} \cdot {}^{\bullet}$

3 21 4. # t₀Ha Inq Hb · ? x · t₃t₄ Fit x^{t₁}

 $Hb \operatorname{Inq} Ha : - t_3 t_4 Hb \operatorname{Sci} \cdot ? \operatorname{Etc} \cdot \operatorname{Qia} t_0 < t_1 < t_2 < t_3 < t_4 \cdot$

 $\wedge : t_0 t_1 Ha \operatorname{Inq} Hb \cdot ?x \cdot t_3 t_4 \operatorname{Fit} x$

-: $t_1 t_2 Hb Sci \cdot Tan \cdot t_3 t_4 Hb Sci \cdot ? x \cdot t_3 t_4 Fit x :$

 $-_1$: $t_1 t_2 \cdot Qqm \cdot Unq Pst \cdot Hb Sci \cdot ? Etc :$

-, $t_1 t_2$. QUP. Hb Sci.? Etc. #

In 3 21 3 Hb has forgotten, in 3 21 4 Hb does not yet know. So QUA means "longer" as used in "no longer"; QUP means "yet" as used in "not yet".

3 22 1. \mathfrak{p} is a metatextual substitute for a noise that lasts longer than t_2-t_1 ; \mathfrak{q} is a metatextual substitute for a noise that lasts less than t_2-t_1 .

```
# Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} \mathfrak{p}!

Hb \operatorname{Inq} Ha \operatorname{Fal}!

Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x!

Hb \operatorname{Inq} Ha \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x!

Hb \operatorname{Inq} Ha \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x!

Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x!

Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x!

Hb \operatorname{Inq} Ha \cdot \operatorname{Qia} \cdot \operatorname{Dur} \mathfrak{p} \cdot > t_2 - t_1!

Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} \mathfrak{q}!

Hb \operatorname{Inq} Ha \cdot \operatorname{Dur} \mathfrak{q} \cdot < t_2 - t_1 \cdot \operatorname{Erg} \cdot \operatorname{Pot} \cdot t_1 t_2 \operatorname{Fit} \mathfrak{q}!

\operatorname{Tan} : ? \cdot \operatorname{Nes} \cdot t_1 t_2 \operatorname{Fit} \mathfrak{q}! ? \cdot \operatorname{Nnc} Hb \operatorname{Sei} \cdot \operatorname{Utr} \operatorname{Ver} #
```

The word written Pot (fL potest=can) means "it is possible that". 'Nes' is known from 3 12 2. Both 'Pot' and 'Nes' are followed by propositions.

 $3\ 22\ 2.$ $(t_1 < t_2 < t_3)$

 $t_3Hc \operatorname{Inq} Hb \cdot t_1 t_2 Ha \operatorname{Inq} Hc \mathfrak{p}^{t_4}$

 $Hd \operatorname{Inq} Hb \cdot t_1 t_2 Ha \operatorname{Inq} Hd \mathfrak{q}^{t_2}$:

 $Hb \operatorname{Inq} Hb : \forall x \cdot t_3 t_4 Hc \operatorname{Inq} Hb x : \in \operatorname{Fal}$:

 $3\ 22\ 3.$ $(t_1 < t_2 < t_3 < t_4)$

$Ha \operatorname{Inq} Hb : t_1 t_2 Ha \operatorname{Ani} \cdot ? x \cdot t_3 t_4 \operatorname{Fit} x$

 $Hb \operatorname{Inq} Ha : \operatorname{Fal} \cdot \operatorname{Qia} \cdot \operatorname{t_2} < \operatorname{t_3} :$

 \neg Pot $t_2 < t_3$. $\land : t_1 t_2$ Ani $? x . t_3 t_4$ Fit x

Pot $t_4 < t_1$. $\wedge : t_1 t_2 Ha$ Ani $? x \cdot t_3 t_4$ Fit $x : t_3 t_4$

Tan : \neg Nes: $t_4 < t_1 \rightarrow t_1 t_2 Ha$ Ani . ? $t_3 t_4$ Fit x #

3 22 4. * Ha Inq Hb. Utr $t_3 t_4$ Hc Sci $? x \cdot t_1 t_2$ Fit x:

 \rightarrow : $t_1 t_2 Hc$ Ani \cdot ? $x \cdot t_1 t_2$ Fit x

 $Hb \operatorname{Inq} Ha \cdot \operatorname{Pot.Tan.} - \operatorname{Nes}$

 $Ha \operatorname{Inq} Hb$ 'Utr: $\vee \lceil t \cdot t' \rceil \cdot tt' Hc \operatorname{Ani}$? Etc:

 $Hb \operatorname{Inq} Ha : \operatorname{Pot.Tan.} \operatorname{\neg Nes}$

Qia Pot: UAN Hd Inq! Hc ? x . t_1 t_2 Fit x #

3 22 5. # $Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot 10^{10000} + 1. \in \operatorname{Pri}$

 $^{t_1}Hb \operatorname{Inq} Ha: 10^{10000} + 1 \in \operatorname{Pri}^{t_2}$

Ha Inq Hb. Cur Ver

 $Hb \operatorname{Inq} Ha: t_1 t_2 Hb \operatorname{Pet}: ?x \cdot x \operatorname{Div} \cdot 10^{10000} + 1$

Ha Inq Hb: Fal: Qia:

Pot: tt' HbPet: $x \cdot Etc' \leftrightarrow t' - t \cdot Mul > t_2 - t_1$

Hb Inq Ha: t1 t2 Hb Sci. Utr Etc. Qia · UAN Hb Pet.? Etc.

 $Ha \operatorname{Inq} Hb : \operatorname{Pot.Tan.} \rightarrow \operatorname{Nes} :$

Pot: $\forall y \cdot \text{UAN } y \text{ Inq } ! Hb \cdot \text{Utr Etc} :$

Eti Pot Qqm: - t₁ t₂ Hb Sei Utr Ete

 $Tan: t_1 t_2 Hb Inq \cdot 10^{10000} + 1. \in Pri *$

3 22 6. # $Ha \operatorname{Inq} Hb \operatorname{PAN} Ha \operatorname{Inq} Hc \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x$:

Tan - - . PAN Hc Rsp Ha

 $Hb \operatorname{Ing} Ha : \neg \cdot \operatorname{Pot} \cdot \operatorname{PAN} Hb \operatorname{Rsp} Ha :$

Qia \neg PAN Hb Ani: Utr $\lor \neg x \cdot y \neg y$ Inq Hcx

 $Ha\operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2 \operatorname{Ftt} x$

 $Hb\operatorname{Inq} Ha$ \longrightarrow Pot Nne $Hb\operatorname{Rsp} Ha$:

Qia PAN Ha Inq Hc.? Etc. Mod Mal

```
Ho Inq Ha. — Itg.

Hb Inq Ha! — Pot. Nnc Hc Rsp Ha:

Qia! — : Nnc Ha Itg! y. PAN Ha Inq Hc y *
```

- 3 22 7. # $Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} \wedge n : \forall p \cdot p \in \operatorname{Pri} \wedge p > n$ $Hb \operatorname{Inq} Ha \cdot \neg \cdot \operatorname{Pot} \cdot \operatorname{Nnc} Hb \operatorname{Rsp} Ha \cdot \operatorname{Qia} : \neg \cdot \operatorname{Nnc} \operatorname{Sci} \cdot ? = \operatorname{Pri} *$
- 3 22 8. # Ha Inq Hb: ? lum x·x \in Pri. \lambda .x > 10^{1001} \cdot Hb Inq Ha \cdot -: Pot \cdot Nnc Hb Inq! Ha .? Etc:

 Ncs \cdot Nnc Hb Pet .? Etc: -- \cdot Nnc .QUP . Hb Sci. ? Etc:

 -- : Pot \cdot Unq \cdot Ant .\(\hat{h}\) Pst . Nnc: Hb Rsp Ha *

 Compare the remark in 3 20 6.
- 3 23 1. * $^{t_1}Ha \operatorname{Inq} Hb \cdot \operatorname{Utr}: \operatorname{Nne} Hb \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:$ $Hb \operatorname{Inq} Ha : \neg \cdot \operatorname{Pot}: \operatorname{Nne} Hb \operatorname{Inq} ! Ha \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:$ $\operatorname{Qia} : \neg : t_1 t_2 Hb \operatorname{Ani} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:$ $Ha \operatorname{Inq} Hc \cdot \operatorname{Utr}: \operatorname{Nne} Hc \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:$ $Hc \operatorname{Inq} Ha : \operatorname{Pot}: \operatorname{Nne} Hc \operatorname{Inq} ! Ha \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:$ $\operatorname{Qia} : t_1 t_2 Hc \operatorname{Ani} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:$ $Ha \operatorname{Inq} Hc \cdot \operatorname{Utr}: \operatorname{MAN} Hc \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:$ $Hc \operatorname{Inq} Ha : \operatorname{Pot} \operatorname{Etc} \cdot \wedge \cdot \operatorname{Pot}: \neg \cdot \operatorname{MAN} \cdot \operatorname{QUA} \cdot Hc \operatorname{Sei} \operatorname{Etc} *$
- 3 23 2. $Ha \operatorname{Inq} Hb : \operatorname{Utr} : \land n : n \in \operatorname{Num} . \rightarrow : \forall p : p \in \operatorname{Pri} . \land . p > n \mid$ $Hb \operatorname{Inq} Ha \operatorname{Ver} \mid$ $Ha \operatorname{Inq} Hb : \operatorname{Utr} : \operatorname{Pot} . \operatorname{Nnc} Hb \operatorname{Dem} \operatorname{Etc} \mid$ $Hb \operatorname{Inq} Ha : \operatorname{Qqm} : \operatorname{MAN} : Hd \operatorname{Dem} \operatorname{Etc} \mid$ $\Rightarrow Hb \operatorname{Ani} : \operatorname{Utr} . \operatorname{Nnc} Hd \operatorname{Dem} \operatorname{Etc} :$ $\operatorname{Tan} : \neg : \operatorname{Pot} . \operatorname{Nnc} Hd \operatorname{Dem} \operatorname{Etc} :$ $\operatorname{Qia} : \neg : \operatorname{MAN} Hb \operatorname{Itg} : ? u : \operatorname{Nnc} Hd \operatorname{Dem} . \operatorname{Etc} . \operatorname{Mod} u \mid$ $Ha \operatorname{Inq} Hc : \operatorname{Utr} : \operatorname{Pot} . \operatorname{Nnc} Hc \operatorname{Dem} \operatorname{Etc} \mid$ $Hc \operatorname{Inq} Ha : \operatorname{Pot} : \neg : \operatorname{Nnc} . \operatorname{Qqm} . \operatorname{PPN} : Hc \operatorname{Dem} \operatorname{Etc} :$ $\operatorname{Qia} : \operatorname{Nnc} Hc \operatorname{Sei} : ? \vdash u : \land \vdash t : x \vdash : \operatorname{Pot} : tt' x \operatorname{Dem} . \operatorname{Etc} . \operatorname{Mod} u \mid$ $\operatorname{Reformulation} \text{ of } 3 : 15 : 5 \text{ with the new means of expression.}$
- 3 23 3. $Ha \operatorname{Inq} Hb \cdot ?x \cdot x^{10} 100 x 111111001111 = 0$ $Hb \operatorname{Inq} Ha : \operatorname{Qqm} \cdot \neg \cdot t_1 t_2 : \operatorname{Tan} \cdot t_3 t_4 \cdot Hb \operatorname{Sci} \cdot Ha \operatorname{Inq} Hb \operatorname{Cur} \cdot Hb \operatorname{Inq} Ha : \operatorname{Qia} : \wedge \lceil a \cdot b \cdot c \rceil : a \wedge b \wedge c \cdot \in \operatorname{Com} \cdot \rightarrow \operatorname{Pot} : \operatorname{Nnc} Hb \operatorname{Cpu} \cdot ?x \cdot ax^{10} + bx + c = 0 *$

Reformulation of 3156 with the new means of expression.

3 23 4. # $Ha \operatorname{Inq} Ha \cdot tt' x \operatorname{Sei} \cdot ? \operatorname{Ete} \cdot \Leftrightarrow : \operatorname{Pot} \cdot tt' x \operatorname{Inq} ! y \cdot ? \operatorname{Ete} *$

```
3 24 1. # Ha \operatorname{Inq} Hb \cdot \operatorname{Ecc} | ^{t_1} - - - - - ^{t_1} |
Ha \operatorname{Inq} Hb \cdot ? x \cdot \operatorname{PAN} \cdot x^{es} \operatorname{Fit} - |
Hb \operatorname{Inq} Ha \cdot 100 \mid
Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} \operatorname{Sci} \mid
Hb \operatorname{Inq} Ha \cdot \operatorname{PAN} Hb \operatorname{Ani} \cdot ? y \cdot \operatorname{Nne} \operatorname{Fit} y :
\wedge \cdot \operatorname{PAN} Hb \cdot \operatorname{Enu} \cdot ? \operatorname{Etc} \mid
Ha \operatorname{Inq} Hb \cdot t_1 \cdot t_2 Ha \cdot \operatorname{Ani} \cdot \operatorname{Utr} \cdot \vee z \cdot t_1 \cdot t_2 Hb \operatorname{Inq} Hb z \cdot 
\operatorname{Tan} : - - \cdot \vee z \cdot t_1 \cdot t_2 Hb \operatorname{Inq} Hb z \mid
Hb \operatorname{Inq} Ha \cdot t_1 \cdot t_2 Hb \cdot \operatorname{Enu} \cdot ? \operatorname{Etc} \cdot :
\operatorname{Mod} \cdot \wedge \cap t \cdot t' \cdot z \cap z \neq Hb \cdot \rightarrow - \operatorname{Pot} \cdot tt' z \operatorname{Ani} \cdot :
? y \cdot t_1 \cdot t_2 Hb \operatorname{Inq} Hb \cdot y \cdot =
\operatorname{Imperceptible} \operatorname{speaking} \operatorname{(thinking)}.
```

imperception speaking (mining)

3 24 2. * $Ha \operatorname{Inq} Hb \cdot \operatorname{PAN} Ha \operatorname{Inq} Hc \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x :$ $\operatorname{Tan} \cdot - \cdot \cdot \operatorname{PAN} Hc \operatorname{Rsp} Ha :$ $Hb \operatorname{Inq} Ha \cdot - \cdot \cdot \operatorname{Pot} \cdot \operatorname{PAN} Hb \operatorname{Rsp} Ha :$ $\operatorname{Qia} \cdot \operatorname{PAN} Ha \operatorname{Inq} Hc \operatorname{Etc} \cdot \operatorname{Mod} : - \cdot \cdot \operatorname{Pot} :$ $\operatorname{Nnc} Hc \operatorname{Ani} \cdot ? y \cdot Ha \operatorname{Inq} Hc y *$

This may be varied to convey that Hc could not understand whether Ha said something, or to whom Ha spoke.

```
3 24 3. * Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:

Hb \operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:

Hc \operatorname{Inq} Hb \circ :

Hb \operatorname{Inq} Ha \circ :

Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} \operatorname{Sei} :

Hb \operatorname{Inq} Ha \cdot \operatorname{PAN} Hb \operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:

\land : \operatorname{PAN} Hc \operatorname{Inq} ! Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:

Ha \operatorname{Inq} Hb \cdot \operatorname{Cur} : \operatorname{Qqm} Hc \cdot \operatorname{Tan} \cdot \neg \cdot Hb \operatorname{Sei} :

Hb \operatorname{Inq} Ha \cdot \operatorname{Qia} : t_1 t_2 \operatorname{Fit} a

\operatorname{Mod} \cdot \operatorname{Qqm} \cdot \operatorname{Pot} \cdot Hc \operatorname{Ani} \cdot \operatorname{Tan} : \neg \cdot \operatorname{Pot} \cdot Hb \operatorname{Ani} *
```

3 24 4. We repeat 3 22 1 with the following variation.

* $Hb \operatorname{Inq} Ha$ ' $\neg \neg$. Pot . $t_1 t_2 \operatorname{Fit} \mathfrak{p}$: Qia · Dur \mathfrak{p} . > Nim I $Ha \operatorname{Inq} Hb$. ? = Nim I $Hb \operatorname{Inq} Ha$. Exg : PAN $Ha \operatorname{Inq} Hb$. $t_1 t_2 \operatorname{Fit} \mathfrak{p}$: Sed : Dur \mathfrak{p} > 'sup : ^ . Dur : ^ x · Pot . $t_1 t_2 \operatorname{Fit} x$ *

The word written Nim (fL nimis = too (as used in "too large")) means "what is admissible". '> Nim' may be translated "too large"; '<Nim', "too small".

3 24 5. Likewise we may vary 3 22 5.

```
* Hb \operatorname{Inq} Ha : \neg \cdot \operatorname{Pot} t_1 t_2 Hb \operatorname{Pet} \cdot ? \operatorname{Etc} :
\operatorname{Qia} : t_2 - t_1 < \operatorname{inf} \cdot ^\operatorname{Dur} : ^\operatorname{r} x : x = \cdot \operatorname{Pet}^{\operatorname{io}} \cdot ? \operatorname{Etc} :
Ha \operatorname{Inq} Hb : \operatorname{Erg} \cdot \operatorname{Qia} \cdot t_2 - t_1 < \operatorname{Nim} *
```

- 3 24 6. I is the written picture of a very large series of time-signals sent within a relatively short time.
 - * $Ha \operatorname{Inq} Hb \operatorname{Ecc} | {}^{t_1} {}^{t_2} |$ $Ha \operatorname{Inq} Hb : {}^{?} x \cdot {}^{t_1} {}^{t_2} \cdot x^{es} \cdot \operatorname{Fit} |$ $Hb \operatorname{Inq} Ha : {}^{?} x \cdot {}^{t_1} {}^{t_2} \cdot x^{es} \cdot \operatorname{Fit} x :> \operatorname{!sup!} {}^{?} \cdot \operatorname{Car} |$ ${}^{*2} : \operatorname{Pot} \cdot Hb \operatorname{Enu} : {}^{?} = \operatorname{Car} z : {}^{?} x \cdot {}^{t_1} {}^{t_2} x^{es} \operatorname{Fit} x \cdot > \operatorname{Nim} *$
- 3 25 2. We shall now adopt the habit of dropping punctuation between and words like \land , \lor , Nes, Pot, and so on. This may be stated by some behaviour rule
 - $* Ha \operatorname{Inq} Ha : \neg \cdot \operatorname{Pot} \operatorname{Etc} \cdot \leftrightarrow \neg \operatorname{Pot} \cdot \operatorname{Etc} *$

```
3 26 1. \# Ha Inq Hb ·? u \cdot u Fit \mathfrak{p}:
                         Hb Inq Ha MAN
                         Ha Inq Hb Cur
                         Hb\operatorname{Inq} Ha \operatorname{Qia}:\mathfrak{p}\operatorname{Ant}\mathfrak{q} \wedge \mathfrak{q}\operatorname{Mul}\operatorname{Ant}\operatorname{Nnc}:
                        Ha Inq Hb. Cur Mul:
                         Hb\operatorname{Inq} Ha: \neg \operatorname{UAN}: \mathfrak{q} \wedge \cdot Hb\operatorname{Ani}.\operatorname{Utr} \mathfrak{q} Erg. MAN:
                         Hc \operatorname{Inq} Hd : \operatorname{Erg} \cdot \operatorname{Nnc} \operatorname{Aet} Hb < \operatorname{Sec} 10^{11101}
                         Hd \operatorname{Inq} Hc \operatorname{Cur}:
                         Hc\operatorname{Inq} Hd: Qia. -\operatorname{UAN} \cdot Hb\operatorname{Ani}. Utr \mathfrak{g}:
                              ^ Cea. Sec 1011101 · Ant Nnc: Fit q:
                        Hd \operatorname{Inq} Hc : \operatorname{Sec} 10^{11101}. Ant \operatorname{Nnc} \cdot Hd \operatorname{Ani}. Utr \mathfrak{q}:
                        Hc \operatorname{Inq} Hd : \operatorname{Erg} \cdot \operatorname{Nne} \operatorname{Aet} Hd > \operatorname{Sec} 10^{11101}
                        Hd \operatorname{Inq} Hc: {}^{\checkmark}t \cdot t \operatorname{Fit} \mathfrak{q} \cdot \operatorname{Aet} Hd \cdot = \operatorname{Cea} \cdot \operatorname{Sec} 10^{11110}
                        Hc \operatorname{Ing} Hd : \operatorname{Erg} \cdot \operatorname{Nne} \operatorname{Aet} Hd := \operatorname{Cea} \cdot \operatorname{Sec} (10^{11101} + 10^{11110})
                        Hd \operatorname{Inq} Hc : \operatorname{Nnc} \operatorname{Aet} Hh \cdot \operatorname{Mul} > .
                              Nnc Aet Ha 	riangleq \operatorname{Nnc} \operatorname{Aet} Hb 	riangleq \operatorname{Nnc} \operatorname{Aet} Hc
```

Nnc Aet $Hh = \text{Cca} \cdot \text{Sec } 10^{11111} \text{ }^{\#}$

In this program text \mathfrak{p} and \mathfrak{q} are metatextual substitutes for some words the meaning of which is still unknown; it might be "earthquake" and "total eclipse of the sun".

The word written Aet (fL aetas=age) means "age". t Aet Ha means "Ha's age at the moment t". The beginning of individual human life has been identified with beginning of individual faculty of perception. Of course this is not correct, but it will be relatively easy to set things right when more refined considerations are needed.

```
3 26 2. * Ha \operatorname{Inq} Ha : \vee \lceil a \cdot b \rceil : a \curvearrowright b \in \operatorname{Tem} \wedge \wedge \upharpoonright \wedge \lceil A \cdot B \rceil:
A = {}^{\uparrow}t \cdot a \leq t \leq b : \wedge : B = {}^{\uparrow}t \cdot 0 \leq t \leq b - a
\rightarrow {}^{\downarrow}t \cdot t \operatorname{Aet} x \in A \curvearrowright B : \wedge
\wedge t : t \in A \cdot \rightarrow \cdot t \operatorname{Aet} x = \cdot t - a \wedge :
\wedge t \cdot t \in A \cdot \rightarrow \cdot t x \operatorname{Ext} : \wedge
\wedge t : t < a \cdot \vee \cdot t > b \cdot \rightarrow \cdot \neg \cdot t x \operatorname{Ext}
\wedge {}^{\dagger}a = {}^{\downarrow}\operatorname{Ini} : x \operatorname{Ext} : \wedge : b = {}^{\downarrow}\operatorname{Fin} : x \operatorname{Ext} +
```

The word written Ext (fL extat = exists) means "exists". It should not be confused with the word \vee . There is (' \vee ') a Roman general who crossed the Rubicon in 49 B.C., but he does not actually exist ('Ext').

It would be premature to analyze the difference between these two notions.

```
3\ 26\ 3. \ # Ha Inq Ha:
```

```
Pot: \forall y \cdot t_1 t_2 x \, \text{Ani} ? y \cdot \Leftrightarrow: Ini . x \, \text{Ext} \cdot < t_1 < t_2 < \cdot \, \text{Fin} . x \, \text{Ext} :

Pot: \forall y \cdot t_1 t_2 x \, \text{Inq} y \, . \, \text{Nnc} \simeq y \simeq Hc \simeq \neg \simeq \text{Etc}

\Leftrightarrow Ini . x \, \text{Ext} \cdot + \cdot \, \text{Cca} \cdot \, \text{Sec} \, 10^{11001} : < t_1 < t_2 < \cdot \, \text{Fin} \, . x \, \text{Ext}:

Pot: t_1 t_2 x \, \text{Enu} . ? Etc . Mod 1 . 10 . 11 . 100

\Leftrightarrow Ini . x \, \text{Ext} \cdot + \cdot \, \text{Cca} \cdot \, \text{Sec} \, 100 \times 10^{11001} : < t_1

< t_2 < \cdot \, \text{Fin} \, x \, \text{Ext}:

Pot: t_1 t_2 x \, \text{Cpu} . ? = 101 × 101

\Leftrightarrow Ini . x \, \text{Ext} \cdot + \cdot \, \text{Cca} \cdot \, \text{Sec} \, 111 \times 10^{11001} : < t_1

< t_2 < \cdot \, \text{Fin} \, x \, \text{Ext}:

Pot: t_1 t_2 x \, \text{Cpu} \cdot ? \, x \cdot x^{10} + 111 \, x - 1000

\Leftrightarrow Ini . x \, \text{Ext} \cdot + \, \text{Cca} \cdot \, \text{Sec} \, 1110 \times 10^{11001} : < t_1

< t_2 < \cdot \, \text{Fin} \, x \, \text{Ext}
```

Average development of humans has been sketched.

3 26 4. p is again a metatextual substitute for some noise.

```
# Ha \operatorname{Inq} Ha:

\operatorname{Pot} \cdot \vee \lceil t_1 \cdot t_2 \cdot x \rceil \cdot t_1 t_2 Ba \operatorname{Inq} x \mathfrak{p}:

\neg \operatorname{Pot} : \vee \lceil t_1 \cdot t_2 \cdot x \rceil \cdot t_1 t_2 Ba \operatorname{Inq} x.

\operatorname{Nnc} \simeq x \simeq Hc \simeq \neg \simeq \operatorname{Etc}:

\operatorname{Pot} : \vee \lceil t_1 \cdot t_2 \cdot y \rceil \cdot t_1 t_2 Ba \operatorname{Ani} \cdot ? y:

\neg \operatorname{Pot} : \vee \lceil t_1 \cdot t_2 \cdot t_3 \cdot t_4 \cdot y \rceil \cdot t_1 t_2 y \operatorname{Inq} Ba \cdot 10 \times 10 = 100 \cdot 10

\wedge : t_3 t_4 Ba \operatorname{Itg} \cdot ? x \cdot t_1 t_2 y \operatorname{Inq} Ba x
```

3 26 5. $*Ha \in \text{Hom.} Hb \in \text{Hom.} Hc \in \text{Hom.} Etc.$

 $Ba \in Bes$, $Bb \in Bes$, $Bc \in Bes$, Etc.

Ha∉Bes. Etc. Ba∉Hom. Etc.

 $\operatorname{Hom} \cap \operatorname{Bes} = \lceil \rceil *$

The words written Hom (fL homo=man) and Bes (fL bestia=animal) mean "human" and "animal" (not human) respectively. The words written Ba, Bb, Bc and so on are names of animals.

3 26 6. # Ha Inq Ha

 $\neg \forall x : x \in \text{Hom.} \land . \text{Nnc} x \text{ Ext.} \land \cdot \text{Nnc Act} x > . \text{Sec } 10^{100000} \text{ #}$

3 26 7. # Ha Inq Ha:

 $\wedge \cdot \operatorname{Nnc} \operatorname{Aet} x \cdot \leq m \land \cdot$

$$a = \sup F = F \cdot \text{Sec } 10^{100000} \longrightarrow a = \text{Cea.} 11 \times 10^{11110}$$

 \wedge F. Sec 10^{11001} /a:=Cca. 0,00001

$$\land$$
 F. Sec 101×10^{11011} /a:= Cca. 0,1

$$\wedge$$
 F. Sec $101 \times 10^{11100} \cdot /a := \text{Cca.} \cdot 0.11$

 \land * F . Sec 101 × 10¹¹¹⁰¹ */a:= Cca.1 *

Short demographic statistic of actually existing mankind.

 $3\ 26\ 8.\ # Ha Inq Ha$:

 $Car: ^{\uparrow}x \cdot x \in Bes. \land .Nnc x Ext ^{\bullet}Mul > .10^{100100} *$

 $3\ 27\ 1.$ # $Ha\operatorname{Inq} Ha$

$$\land n: Fn . \Leftrightarrow \lor \ulcorner a.b.c \urcorner : a \cup b \cup c. \subset \text{Num} \cdot \land . a^n + b^n = c^n :$$

$$\wedge \cdot - n \operatorname{Div} a \cdot \wedge \cdot - n \operatorname{Div} b \cdot \wedge \cdot - n \operatorname{Div} c \cdot \wedge \cdot a^n + b^n = c^n$$

$$\land : G \Leftrightarrow \neg \lor n : n - 10 \in \mathbf{Num} \cdot \land : Fn$$

$$\rightarrow :: \neg \forall x : x \in \operatorname{Hom} . \land . \operatorname{Nne} x \operatorname{Ext} . \land \cdot \operatorname{Nne} x \operatorname{Sei} . \operatorname{Utr} G$$

Tan . . $\forall x \mid x \in \text{Hom} . \land : \text{Unq} \cdot \text{Sec } 101 \times 10^{11111} . \text{Ant} \cdot \text{Nnc} : x \text{ Ext}$ $\text{$\sim$ Inq } x : \forall y : y \in . \text{Dem}^{\text{io}} G \cdot$

 \wedge Nnc x Sci : Utr $\cdot y \in \operatorname{Dem}^{\mathrm{io}} G \cap A$

 $\forall x : x \in \text{Hom.} \land : \text{Unq} \cdot \text{Sec } 111 \times 10^{11101} . \text{Ant} \cdot \text{Nnc } x \text{ Ext} :$

 $x \operatorname{Dem} An: n-10 \in \operatorname{Num} An = 1100100 \longrightarrow An = Fn$

 $\forall x : x \in \text{Hom.} \land . \text{Nne } x \text{ Ext.} \land : PAN x \text{ Dem}$

$$\wedge n$$
 : $n-10$ \in Num . \wedge . n $<$ 1001101100 o . o Fn . \wedge .

 $\forall x : x \in \text{Hom.} \land . \text{Nnc} x \text{ Ext.} \land . \text{PAN} x \text{ Dem}$

 $\land n:n\,{\in}\,\mathrm{Pri}\,\ldotp\,\land\,\ldotp\, n\,{\neq}\,10\,\ldotp\,\land\,\ldotp\, n\,{<}\,10^{11100}\,\lq\,{\rightarrow}\,\ldotp\,\neg\, F'\,n$

Sketch of the history of Fermat's theorem.

3 28 1. Ha Inq Ha 1

$$t_1t_2Ha \operatorname{Inq} Hb \curvearrowright Hc \curvearrowright Hd \curvearrowright He \curvearrowright Hf \cdot p$$

 $\rightarrow t_1 t_2 Ha \operatorname{Inq} \cdot 101 \operatorname{\bullet} \operatorname{Hom} \cdot p$

```
t_1 t_2 \cdot Ha \wedge Hb \wedge Hc \cdot \text{Ani } p \cdot \rightarrow \cdot t_1 t_2 \cdot 11^\epsilon \text{Hom . Ani } p :
t_1 t_2 Ha \text{Inq . } n^\epsilon \text{Hom . } p \cdot \leftrightarrow : \lor A \cdot A \cdot C \text{ Hom . } \wedge \cdot \text{Car } A \cdot = n \cdot \wedge : \land x \cdot x \in A \cdot \rightarrow \cdot t_1 t_2 Ha \text{Inq } x p :
t_1 t_2 \cdot n^\epsilon \text{Hom . Ani } p \cdot \leftrightarrow : \lor A \cdot A \cdot C \text{ Hom . } \wedge \cdot \text{Car } A \cdot = n \cdot \wedge : \land x \cdot x \in A \cdot \rightarrow \cdot t_1 t_2 x \text{Ani } p :
t_1 t_2 Ha \text{Inq Mul}^\epsilon \text{Hom } p \cdot \leftrightarrow : \lor n \cdot \text{Mul} > \cdot 1 \cdot \wedge \cdot t_1 t_2 Ha \text{Inq . } n^\epsilon \text{Hom . } p :
Etc. n^\epsilon \text{Rea . Etc} \cdot \leftrightarrow : \lor A \cdot A \cdot C \text{Rea . } \wedge \cdot \text{Car } A \cdot = n \cdot \wedge : \land x \cdot x \in A \cdot \rightarrow \cdot \text{Etc } x \text{ Etc} :
Etc. \sharp
```

By the use of the pattern ' $n^{\epsilon}A$ ', meaning "n things from A", Lincos speech becomes still more akin to common languages which admit e.g. "three people" as a subject or an object of a sentence.

Common languages do not distinguish between the "three apples" of

I bought three apples,

and of

Three apples of this kind go to a pound,

or between the "little money" of

I have little money,

and of

With little money you can succeed at Smith's,

though the variable that ranges over the triples of apples or over the small quantities of money has undergone essentially different binding.

We shall carefully use the pattern ' $n^{\epsilon}A$ ', always remembering that it has been defined by existential binding.

 1^{ϵ} can be used as indefinite article.

```
3 28 2. * Ha \operatorname{Inq} Hb \cdot ?x \cdot x = 110/100 :
	Hb \operatorname{Inq} Ha \cdot 11/10 :
	Hc \operatorname{Inq} Hd \cdot ?y \cdot \operatorname{PAN} \operatorname{Fit} y :
	Hd \operatorname{Inq} Hc \cdot \operatorname{PAN} : Ha \operatorname{Inq} Hb \cdot ?x \cdot \operatorname{Etc} \curvearrowright Hb \operatorname{Rsp} Ha :
	Hc \operatorname{Inq} Hd \cdot \operatorname{PAN} Ha \cdot Hb \operatorname{Ise} \operatorname{Inq} :
	Hd \operatorname{Inq} Hc \cdot \operatorname{Ben} *
```

The word written Ise (fL inter se=among themselves) means "to each other". ' $Ha\ Hb$ Ise Inq' is to be translated "Ha and Hb have a talk".

```
3 28 3. # Ha \operatorname{Inq} Hb \mathfrak{p} \cdot Hb \operatorname{Inq} Hc \mathfrak{q} \cdot Hc \operatorname{Inq} Ha \mathfrak{r} \cdot Ha \operatorname{Inq} Hb \smile Hc \smile Hd \mathfrak{F} \cdot Hd \operatorname{Inq} Ha \mathfrak{t} : He \operatorname{Inq} Hf \cdot \mathfrak{f} \cdot x \cdot \operatorname{PAN} \operatorname{Fit} x : Hf \operatorname{Inq} He \cdot \operatorname{PAN} Ha Hb Hc Hd \operatorname{Ise} \operatorname{Inq} #
```

3 28 4. $* Ha \operatorname{Inq} Hb \mathfrak{p} . Hb \operatorname{Inq} Hc \mathfrak{q} . Hc \operatorname{Inq} Ha \mathfrak{r}$.

Ha Inq Hb & . Hb Inq Hct. Etc.

He Inq Hf · PAN Ha Hb Hc Alt Inq #

The word written Alt (fL alternando=alternating) means "alternating".

```
3 29 1. # {}^{t_1}Ha \operatorname{Inq} Hb \cdot ?x \cdot 101 x = 1111 {}^{t_2} - - - - :
Ha \operatorname{Inq} Hc : \neg \operatorname{Nnc} Hb \operatorname{Sci} \cdot ?x \cdot 101 x = 1111 :
Hc \operatorname{Inq} Ha \operatorname{Cur} :
Ha \operatorname{Inq} Hc \cdot \neg \operatorname{PAN} Hb \operatorname{Rsp} :
Hc \operatorname{Inq} Ha \cdot \operatorname{MAN} \operatorname{Usq} \operatorname{Nnc} \cdot Hb \operatorname{Sci} : \operatorname{Tan} \cdot \neg \operatorname{Pot} \cdot \operatorname{Nnc} Hb \operatorname{Rsp} :
Ha \operatorname{Inq} Hc \operatorname{Cur} :
Hc \operatorname{Inq} Ha \cdot \neg t_1 t_2 Hb \operatorname{Ani} : ?u \cdot \operatorname{Nnc} Ha \operatorname{Inq} Hb u \cdot #
Compare the remark in 3 20 6.
```

- 3 29 2. A variation on 3 29 1 with Itg instead of Ani.
- 3 29 3. * $Ha \operatorname{Inq} Hb \cdot ?x \cdot 101x = 1111 - Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} \cdot \operatorname{Nne} Hb \operatorname{Sei} \operatorname{Ete} :$ $Hb \operatorname{Inq} Ha \cdot \operatorname{Nne} Hb \operatorname{Sei} \operatorname{Ete} :$ $Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot - \operatorname{Pot} \cdot \operatorname{Nne} Hb \operatorname{Rsp} :$ $Hb \operatorname{Inq} Ha \cdot \operatorname{Pot} \cdot \operatorname{Nne} Hb \operatorname{Rsp} :$ $Ha \operatorname{Inq} Hb : \operatorname{Cur} \cdot - \cdot \operatorname{Nne} Hb \operatorname{Rsp} :$ $Hb \operatorname{Inq} Ha : \operatorname{Qia} \cdot \operatorname{Nne} Hb \operatorname{Vul} :$ $- \operatorname{Nne} Hb \operatorname{Inq} ! Ha \cdot ?x \cdot 101x = 1111 *$

The word written Vul (fL vult=wishes) means (the verb) "wishes". More texts are needed in order to grasp this meaning, but even now the receiver can conclude that somebody might refuse to answer a question though able to do so.

'Vul' obeys the following syntactical rules: As in the case of 'Pot' and 'Ncs' the object (the thing wanted) is to be a (possibly elliptical) proposition, not an event. Yet the use of 'Vul' differs from that of 'Pot' and 'Ncs' insofar as Ha Vul p is to be treated as an event. Of course ' t_1 t_2 Ha Vul p' is a report, hence a proposition.

By the use of 'Vul' neither a (more or less mysterious) state of mind of the person wishing nor an introspection of the person saying "I wish" is aimed at, 'Vul' describes the behaviour of people who use this word explicitly in the first texts that follow. Later on it will appear that a person may show the behaviour of wishing though he does not explicitly state "I wish". In the case of 'Vul' we shall proceed in a way inverse to that we followed when treating 'Enu', 'Cpu', and so on. There we started with instances of the behaviour of counting and computing, and afterwards introduced words to describe these activities. Now we use 'Vul' from the beginning, whereas in a more advanced stage we shall learn to do without it.

Note that as a symptom the explicit use of 'Vul' is not sufficient to prove the fact of wishing. It may happen that someone claims to wish something though it is clear by circumstancial evidence that he does not wish at all or that he wishes something else. In this respect statements preceded by "I wish" do not differ from other statements about other people or about the speaker himself. Statements may always be false, intentionally or unintentionally.

 $Hb \operatorname{Inq} Ha : \neg : \forall y : t_1 t_2 y \operatorname{Ani} : ?x : t_1 t_2 \operatorname{Fit} x : \land : y \neq Hc :$

 $# Ha Inq Hb \cdot ? x \cdot t_1 t_2 Fit x$

 $Ha \operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$

```
Hc \operatorname{Ing} Ha : \neg 'Nne Hc \operatorname{Vul} : \operatorname{Nne} Ha \operatorname{Ing} ! Hc \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x :
                     Ha \operatorname{Ing} Hb : \neg : \operatorname{Nnc} Hc \operatorname{Vul} : \operatorname{Nnc} Ha \operatorname{Sci} : ? \operatorname{Etc} :
                     Hc \operatorname{Ing} Ha : \operatorname{Eti} \operatorname{Nne} Hc \operatorname{Vul} : - \operatorname{Nne} Ha \operatorname{Sei} ? \operatorname{Ete} :
                      Ha \operatorname{Ing} Hb \operatorname{Erg}^* \longrightarrow \operatorname{Pot}^* \operatorname{Nnc} Ha \operatorname{Sei} ? \operatorname{Etc} *
3 29 4. # Ha Ing Hc · Nne Ha Vul . PPN Hc Ing Hd p :
                     Hc\operatorname{Inq} Hd\mathfrak{p}ı
                      Ha \operatorname{Inq} Hc \operatorname{Ben}
                      Hb \operatorname{Ing} Ha 'Nne Ha \operatorname{Vul} \cdot \operatorname{PPN} Hc \operatorname{Etc} \cdot \rightarrow \cdot \operatorname{PPN} Hc \operatorname{Etc} :
                           Nnc Ha \operatorname{Vul} q \longrightarrow \operatorname{PPN} Hc \operatorname{Vul} q
                      Ha Ing Hb ' Nnc Ha Vul . PPN Hc Ing Hbr
                           Tan. \neg PPNHc Inq Hbr:
                           Qia: \neg Nnc Hb Ani ?x. Nnc Ha Inq Hcx
                      Hb \operatorname{Inq} Ha : \operatorname{Nne} Ha \operatorname{Vul} q . \wedge : \operatorname{Nne} Hb \operatorname{Sei} \cdot
                           Utr. Nnc Ha Vul q \rightarrow PPN Hc Vul q \rightarrow
                      Hb \operatorname{Inq} Hc : \operatorname{Cur} t Ha \operatorname{Vul} q \longrightarrow : \operatorname{Pau} \operatorname{Pst} t \cdot Hc \operatorname{Vul} q
                      Hc\operatorname{Ing} Hb : tHa\operatorname{Vul} \cdot \operatorname{PPNFit} s \longrightarrow \operatorname{PauPst} \cdot t \cdot \operatorname{Fit} s : \longrightarrow s \in \operatorname{Ben} I
                      Ha \operatorname{Ing} Hc : \operatorname{Nnc} Ha \operatorname{Vul} \cdot \operatorname{PPN} Hc \operatorname{Ing} Hb \cdot 10 \times 10 = 101
                      Hc \operatorname{Inq} Hb \cdot 10 \times 10 \neq 101
                      Ha\operatorname{Inq}Hc\operatorname{Mal}
                      Hc Inq Ha: - Pot tHa Vul: Pau Pst. t. Hc Inq Hb. 1 Fal:
                           Erg \rightarrow Nne Ha Vul \cdot PPN Hc Inq Hb . 10 \times 10 = 101:
```

```
Sed Vul Etc≠ı
Ha \operatorname{Inq} Hc \cdot \operatorname{Pot} \cdot t Ha \operatorname{Vul} \cdot
Hc \operatorname{Inq} Hb \cdot \operatorname{PAN} Ha \operatorname{Vul} \cdot \operatorname{PPN} Hc \operatorname{Inq} Hb \cdot 10 \times 10 = 101:
      Erg · Nnc Hc Inq Hb · 10 \times 10 = 101 : Tan Fal
Ha \operatorname{Inq} Hc \cdot \operatorname{Nnc} Ha \operatorname{Vul} \cdot 10 \times 10 = 101
Hc \operatorname{Inq} Ha : \operatorname{Pot} \operatorname{Qqm} \cdot \operatorname{PAN} Ha \operatorname{Inq} Hc \cdot \operatorname{Nnc} Ha \operatorname{Vul} :
      Tan . \neg Nnc Ha Vul:
      Qia '\forall w \cdot PAN Ha Inq Hc \cdot Nnc Ha Vul w : \in Fal
{\it Ha}\,{
m Inq}\,{\it Hc}\,{
m Ben} ı
Hc\operatorname{Inq} Ha 'Utr: \operatorname{Nnc} Ha\operatorname{Vul}\cdot \vee x. \operatorname{Nnc} Hc\operatorname{Inq} Hbx.
Ha\operatorname{Inq} Hc . \operatorname{\neg}\operatorname{Nnc} Ha\operatorname{Vul} .
Hc\operatorname{Inq} Ha 'Utr: Nne Ha\operatorname{Vul} · \neg \lor x. Nne Hc\operatorname{Inq} Hbx.
Ha \operatorname{Inq} Hc . \longrightarrow \operatorname{Nnc} Ha \operatorname{Vul} I
Hb \operatorname{Inq} Hb - \forall x \cdot \operatorname{Nnc} Ha \operatorname{Vul} x \cdot
Ha\operatorname{Inq} Hb . PPN Hc\operatorname{Inq} Hd\mathfrak{p} ,
Hb\operatorname{Inq} Ha ? u: Nne Ha\operatorname{Ani} · Utr. PPN Hc\operatorname{Inq} Hd\mathfrak{p} · \operatorname{Mod} u .
Ha \operatorname{Inq} Hb : \neg \lor u \operatorname{Etc} \operatorname{Mod} u
      Qia: — Pot · Nne Ha Ani . Utr PPN Etc .
Hb\operatorname{Ing} Ha ? u:\operatorname{Nnc} Ha\operatorname{Sei} · Utr. PPN Hc\operatorname{Ing} Hd\mathfrak{p} · \operatorname{Mod} u .
Ha \operatorname{Inq} Hb \operatorname{Mod} \cdot \operatorname{Nnc} Ha \operatorname{Vul} \cdot \operatorname{PPN} Hc \operatorname{Inq} Hd \mathfrak{p}:
      Qia · Nnc Ha Vul q \rightarrow . PPN Hc Vul q \rightarrow .
Ha\operatorname{Inq} Hc \cdot \operatorname{Nnc} Ha\operatorname{Vul} \cdot \operatorname{PPN} Hc\operatorname{Inq} Hd\mathfrak{p}ı
Hc\operatorname{Inq} Hd\mathfrak{p}
He Inq Hb. PPN Hc Inq Hdr.
Ha \operatorname{Inq} Hb . - \operatorname{PPN} Hc \operatorname{Inq} Hd \mathfrak{r}
He Inq Hc · Nnc He Vul. PPN Hc Inq Hdr.
Ha \operatorname{Ing} Hc \cdot \operatorname{Nnc} Ha \operatorname{Vul} = \operatorname{PPN} Hc \operatorname{Ing} Hd \mathfrak{r}
Hc \operatorname{Inq} Hb^{-\gamma}x \cdot \operatorname{Nnc} He \operatorname{Vul} \cdot \operatorname{PPN} \operatorname{Fit} x:
             \neq : \forall x \cdot \text{Nnc} \, Ha \, \text{Vul} \cdot \text{PPN Fit} \, x \, \text{Erg} \cdot He \, \text{Vul Etc} \cdot \in \text{Mal} \cdot
 Hb\operatorname{Ing} Hc\cdot\operatorname{Nnc} Hb\operatorname{Vul}\cdot\operatorname{PPN} Hc\operatorname{Ing} Hd\mathfrak{a}
Hc \operatorname{Inq} Ha : \operatorname{Utr} \cdot \operatorname{Nnc} Ha \operatorname{Vul} \cdot - \operatorname{PPN} Hc \operatorname{Inq} Hd a
 Ha \operatorname{Inq} Hc . - \operatorname{Nnc} Ha \operatorname{Vul}
 Hc \operatorname{Inq} Ha : \operatorname{Utr} \cdot \operatorname{Nnc} Ha \operatorname{Vul} \cdot \operatorname{PPN} Hc \operatorname{Inq} Hd a
 Ha \operatorname{Ing} Hc \cdot - \operatorname{Nnc} Ha \operatorname{Vul} \cdot
 Hb\operatorname{Inq} Hc\operatorname{\cdot}\operatorname{Nne} Ha\operatorname{Ced}\operatorname{\cdot}\operatorname{Nne} Hc\operatorname{Vul}\operatorname{\cdot}^{\vee}x\operatorname{\cdot}\operatorname{Nne} Hb\operatorname{Vul}x\operatorname{\cdot}
 Hc\operatorname{Inq} Hb \wedge \operatorname{Ced} : -\operatorname{Nne} Hc\operatorname{Vul} \cdot \Upsilon x . \operatorname{Nne} Hb\operatorname{Vul} x .
 Hc \operatorname{Inq} Hd \mathfrak{a} *
```

The word written Ccd (fL concedit = allows) means "allows", "suffers". Syntactically it behaves as 'Vul'. The thing allowed is to be a proposition, though 'Ha Ccd Etc' is to be treated as an event. Again ' t_1 t_2 Ha Ccd Etc' is a proposition. The relation between 'Vul' and 'Ccd' is akin to that between 'Ncs' and 'Pot'.

```
3295. # Hb Inq Hb
                             tt' Ha \operatorname{Ccd} \operatorname{Etc} . \Leftrightarrow - tt' Ha \operatorname{Vul} . - \operatorname{Etc} :
                              -tt' Ha \operatorname{Ced} \operatorname{Etc} . \Leftrightarrow tt' Ha \operatorname{Vul} . - \operatorname{Etc} :
                              \neg tt' Ha \operatorname{Ced} \cdot \neg \operatorname{Ete} \cdot \leftrightarrow \cdot tt' Ha \operatorname{Vul} \operatorname{Ete} :
                             tt' Ha \operatorname{Ccd} \longrightarrow \operatorname{Etc} \longrightarrow \operatorname{tt'} Ha \operatorname{Vul} \operatorname{Etc} =
3 29 6.
                        Ha \operatorname{Inq} Hd \cdot \operatorname{Nnc} Ha \operatorname{Vul} \cdot \operatorname{Nnc} Hd \operatorname{Inq} Hc \mathfrak{p}
                        Hd \operatorname{Inq} Ha . - \operatorname{Nnc} Hd \operatorname{Inq} Hc \mathfrak{p}:
                        Ha Inq Hd Cur:
                        Hd \operatorname{Inq} Ha . \neg \operatorname{Nes}
                        Ha \operatorname{Inq} Hd \cdot \operatorname{Tan} \cdot \operatorname{Nnc} Ha \operatorname{Vul}:
                        Hd Inq Ha Cur !
                        Ha \operatorname{Ing} Hd : \operatorname{Nnc} Ha \operatorname{Vul} \cdot \operatorname{Nnc} Ha \operatorname{Sei} \cdot \operatorname{Utr} \mathfrak{p}:
                             Qqm - Pot : Nnc Hc Ani \cdot ?x \cdot Nnc Ha Inq Hc x
                             Tan Pot: Nne Hc Ani ?x. Nne Hd Inq Hcx:
                        Hd \operatorname{Ing} Hc \mathfrak{p}:
                        Hc\operatorname{Inq} Hd: ?u \cdot \operatorname{Unq} \cdot \operatorname{Ant} \operatorname{Nnc} \cdot Hd\operatorname{Cog} \cdot \operatorname{Utr} \mathfrak{p} \cdot \operatorname{Mod} u
                        Hd \operatorname{Ing} Hc \cdot \operatorname{PAN} Ha \operatorname{Vul} \cdot \operatorname{PPN} Hd \operatorname{Ing} Hc \mathfrak{p}:
                        Hc \operatorname{Inq} Hd : \operatorname{Nnc} Hc \operatorname{Vul} \cdot \operatorname{PPN} Hd \operatorname{Inq} Ha \cdot \operatorname{PAN} Hc \operatorname{Cog} \operatorname{Utr} \mathfrak{p}:
                        Hd \operatorname{Inq} Ha. Nnc Hc \operatorname{Sci} \operatorname{Utr} \mathfrak{p} *
3 29 7. # Ha Inq He · Nnc Ha Vul . Nnc He Inq Hc p
                        He \operatorname{Inq} Ha . \neg \operatorname{Nnc} He \operatorname{Vul}:
                       Ha \operatorname{Ing} He : \neg \operatorname{Pot} : \operatorname{Nnc} Hc \operatorname{Ani} \cdot ?x \cdot \operatorname{Nnc} Ha \operatorname{Ing} Hc x
                             Erg · Nnc Ha Vul . Nnc He Inq Hcp:
                       He \operatorname{Ing} Ha \operatorname{Tan} : - \operatorname{Nnc} He \operatorname{Vul} \cdot x \cdot \operatorname{PPN} Ha \operatorname{Vul} x :
                       Ha \operatorname{Ing} Hd : \operatorname{Utr} \operatorname{^{\bullet}}\operatorname{Nnc} Hd \operatorname{Vul} : \operatorname{Nnc} Hd \operatorname{Ing} He \operatorname{^{\bullet}}
                             Nnc Hd Vul. Nnc He Ing Hc p:
                       Hd \operatorname{Inq} Ha \operatorname{Ver}
                       Hd \operatorname{Ing} He \cdot \operatorname{Utr} \operatorname{Nnc} He \operatorname{Vul} \cdot \operatorname{Nnc} He \operatorname{Ing} Hc \mathfrak{p}
                       He \operatorname{Inq} Hc \mathfrak{p} *
3298. # Ha Inq Hb.? = Itg.
                       Hb \operatorname{Inq} Ha : \operatorname{Exg} : \operatorname{Nnc} Hc \operatorname{Inq} Hd p.
                             \rightarrow 'Nnc Hd Itg \cdot? x . Nnc Hc Inq Hd x:
                                   \Leftrightarrow: Nnc Hd Sci ?x. Nnc Hc Inq Hdx
                       Ha \operatorname{Ing} Hb : \operatorname{PAN} Hc \operatorname{Ing} Hd \cdot ?x \cdot t_1 t_2 \operatorname{Ftt} x :
                             \wedge: Pau Pst. PAN · Hd \operatorname{Inq} Hc. t_1 t_2 \operatorname{Fit} \mathfrak{a}
                             Utr: PAN Hd Itg \cdot ?y \cdot PAN Hc Inq Hd y \cdot
                       Hb \operatorname{Inq} Ha : \operatorname{Ver} : \operatorname{Nne} Hd \operatorname{Itg} : ?x . \operatorname{Nne} Hc \operatorname{Inq} Hd x :
                             \leftrightarrow Nne Hd Sci : ? x · Nne Hc Vul · Nne Hc Inq Hd x #
```

3 29 9. Comments on 3 23 1-3 23 4 in order to state that somebody "can if he will". E.g. on 3 23 4:

```
* Ha Inq Ha: tt'x Sci. ? Etc.
                               \Leftrightarrow "tt' x Vul · Nnc x Inq ! y . ? Etc : \Rightarrow · tt' x Inq ! y . ? Etc #
3 30 1. \#^{t_1}Ha \operatorname{Inq} Hb \operatorname{Utr} : \operatorname{Nnc} Hb \operatorname{Sci} ? x \cdot t_1 t_2 \operatorname{Fit} x
                          Hb \operatorname{Inq} Ha \cdot \operatorname{Ver} : \operatorname{Tan} : \operatorname{Nnc} Hb \operatorname{Vul} :
                                 \wedge \lceil t \cdot t' \rceil \rceil = tt' Hc \operatorname{Sci} ? x \cdot t_1 t_2 \operatorname{Fit} x \rceil
                          Ha \operatorname{Ing} Hb : \operatorname{Nnc} Ha \operatorname{Vul} : \operatorname{PPN} Hb \operatorname{Ing} ! Ha : ? x \cdot t_3 t_2 \operatorname{Fit} x :
                                       \rightarrow: \land \lceil t \cdot t' \rceil \rceil - : tt' Ha \operatorname{Inq} ! Hc \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x \mid
                          Hb \operatorname{Ing} Ha^{: \vee}v \cdot \operatorname{Nnc} Hb \operatorname{Vul} v \cdot \operatorname{Mul} > \cdot^{\vee}u \cdot \operatorname{Nnc} Ha \operatorname{Vul}
                                 Ha \operatorname{Vul} - tt' Ha \operatorname{Inq} Hc \operatorname{Etc}
                                 Eti: Hb \operatorname{Vul} : \neg \vee \lceil t \cdot t' \cdot y \rceil \cdot Ha \operatorname{Inq} ! y \operatorname{Etc} :
                                 \neg \operatorname{Sat} \cdot \neg \vee \lceil t \cdot t' \rceil Ha \operatorname{Inq} ! \operatorname{He} \operatorname{Ete} :
                                 Qia 'Pot: \forall x \cdot x \neq Ha \cdot \land \cdot UPN x \text{ Inq } ! Hc \text{ Etc } !
                          Ha\operatorname{Ing} Hb \cdot \operatorname{Nnc} Ha\operatorname{Pol} Hb \cdot \longrightarrow \vee \lceil t \cdot t' \cdot y \rceil \operatorname{Etc} \mathsf{I}
                          Hb\operatorname{Inq} Ha \cdot t_1 t_2\operatorname{Fit} \mathfrak{a}
                          Hd \operatorname{Inq} Ha^* \operatorname{Utr} : \operatorname{Nne} Ha \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x
                          {}^{\mathsf{t}_{\mathsf{a}}} Ha \operatorname{Inq} Hd : \operatorname{Ver} : \operatorname{Tan}^{\mathsf{c}} \operatorname{PAN} Ha \operatorname{Pol} Hb :
                                 \neg UPN Ha Inq! Hd ·? x · t<sub>1</sub> t<sub>2</sub> Fit x |
                          Hd \operatorname{Inq} Ha \cdot \operatorname{Utr} \cdot \operatorname{Nne} Ha \operatorname{Inq} ! He \operatorname{Ete} \mathsf{I}
                          Ha Inq Hd: Eti: PAN Ha Pol Hb.
                                       \neg \lor \ulcorner t \cdot t' \cdot y \urcorner \cdot tt' Ha \operatorname{Inq} ! y \operatorname{Etc} "
                                 Erg: Nne Ha Vul \cdot \neg \lor \lceil t \cdot t' \cdot y \rceil \cdot tt' Ha Inq! y Etc *
```

The word written Pol (fL pollicetur = promises) means the verb "promises". Syntactically it behaves as 'Inq'.

The word written Sat (fL satis=enough) means "enough".

3 30 2. # $Ha \operatorname{Inq} Ha \operatorname{Itt}' Ha \operatorname{Inq} Hb$: Nne Ha Vul 'Ete: h Pst. $t' \cdot Ha$ Inq Ete: \leftrightarrow tt' Ha Pol Hb * Etc: h Pst. t' · Ha Inq Etc # Explicit definition of 'Pol'.

```
3 30 3.
                         (t_1 < t_2 < t_3 < t_4)
                    # Ha \operatorname{Inq} Hb : \operatorname{Utr} \operatorname{Nnc} Hb \operatorname{Sei} \cdot ? y \cdot t_3 t_4 \operatorname{Fit} y :
                         Hb \operatorname{Inq} Ha \operatorname{Ver} \Xi
                         Ha \operatorname{Inq} Hb : \operatorname{Nnc} Ha \operatorname{Vul}
                               Nne Hb Inq! Ha \cdot ? y \cdot t_3 t_4 Fit y:
                                     \rightarrow: PPN Ha Inq! Hb ? x . t_1 t_2 Fit x:
                         Hb \operatorname{Inq} Ha : \operatorname{Nne} Hb \operatorname{Vul} \cdot \Upsilon z \cdot \operatorname{PAN} Ha \operatorname{Vul} z :
                         Ha \operatorname{Inq} Hb \operatorname{Ben}:
                         Hb \operatorname{Inq} Ha : t_3 t_4 \operatorname{Fit} \mathfrak{b} \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x :
                         Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} \mathfrak{a}:
                         Hc \operatorname{Inq} Ha \operatorname{Cur} : Ha \operatorname{Inq} ! Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x :
                         Ha \operatorname{Inq} Hc : \operatorname{Qia} \cdot \operatorname{PAN} Ha \operatorname{Pol} Hb. Etc *
```

3 30 4. We repeat the greater part of 3 30 3, but now Ha refuses to answer.

Ha Inq Hb: — Nnc Ha Vul. Nnc Ha Inq! Hb Etc! Hb Inq Ha: PAN Ha Pol Hb. PPN Ha Inq! Hb Etc!

 $Ha \operatorname{Inq} Hb : \operatorname{Fal} : \operatorname{PAN} Ha \operatorname{Pol} : \operatorname{Nne} Hb \operatorname{Inq} : Ha : ?y : t_3 t_4 \operatorname{Fit} y :$

 \rightarrow : PPN Ha Inq! $Hb \cdot ? x \cdot t_1 t_2$ Fit x :

Hb Inq Ha · PAN Hb Inq Ha . t3 t4 Fit b:

t. Ha Inq Hb - PAN Ha Ani · Utr . Hb Inq Ha Ete: to

t1 t2 Fit a:

Hb Inq Ha Ben :

 $Hc \operatorname{Inq} Hc \cdot t_5 t_6 Ha \operatorname{Inq} Hb u \rightarrow u \in \operatorname{Fal} *$

3 31 1. We introduce the abbreviations:

- * $t \cdot p \rightarrow q$ * instead of * $t \operatorname{Fit} p \cdot \rightarrow \cdot t \operatorname{Fit} q$ *
- * $t \cdot p \longrightarrow h \operatorname{Pst} \cdot q$ * instead of * $t \operatorname{Fit} p \longrightarrow : h \operatorname{Pst} \cdot t \cdot \operatorname{Fit} q$ *
- * $t \cdot p \rightarrow h \text{ Ant. } q$ * instead of * $t \text{ Fit } p \rightarrow h \text{ Ant. } t \cdot \text{ Fit } q$ *

This can be done by means of repetitions of former program texts and finally by means of a behaviour rule.

As a matter of course the arrow in ' $t \cdot p \rightarrow q$ ' does not state an implication from p to q, because p and q do not mean propositions, but events. If any doubt about the character of p and q is possible, it is better not to use this abbreviation. Likewise the "retarded and retroactive implications" ' $\rightarrow h$ Pst' and ' $\rightarrow h$ Ant' are not proper implications.

Note that

$$`t \cdot p . \to h \text{ Pst } . q"$$

cannot be expressed by

$$q \cdot h \operatorname{Pst} \cdot p$$
 or $h \operatorname{Pst} \cdot p \cdot \operatorname{Fit} q$

because in the latter sentences all variables are existentially bound, whereas the t in the first sentence is bound by generalizing.

3 32 1. A comment on 3 30 3 and 3 30 4:

* $Hb \operatorname{Inq} Ha : tt' Ha \operatorname{Pol} Hb : h \operatorname{Pst} \cdot t' \cdot Ha \operatorname{Inq} \operatorname{Etc}$ $\rightarrow \operatorname{Deb} : h \operatorname{Pst} \cdot t' \cdot Ha \operatorname{Inq} \operatorname{Etc} :$

Ha Inq Hb: tt' Ha Pol Hb

t'' Fit $p \rightarrow h$ Pst $t'' \cdot Ha$ Inq Etc:

 \rightarrow Deb t'' Fit $p \rightarrow h$ Pst $t'' \cdot Ha$ Inq Etc *

The word written Deb (#L debet=is obliged, ought) means 'obliged'. Syntactically it behaves as 'Ncs' and 'Pot'. In 'Deb Etc', 'Etc' is to mean a proposition. 'Deb Etc' itself means a proposition. Semantically 'Ncs' and 'Deb' differ. This difference will be analyzed later on.

An analogous comment may be added to 3 30 1.

3 32 2. The text 3 30 1 from t₃ onwards is changed into

```
* Ha \operatorname{Inq} Hd ' \operatorname{Ver} ' \operatorname{Tan} ' \operatorname{\sqsubseteq} Lie : UPN Ha \operatorname{Inq} ! Hd ' ? x \cdot t_1 t_2 \operatorname{Fit} x : Hd \operatorname{Inq} Ha \cdot \operatorname{Cur} \cdot \operatorname{\sqsubseteq} Lie : Ha \operatorname{Inq} Hd : \operatorname{Qia} \cdot \operatorname{PAN} Ha \operatorname{Pol} Hb \cdot \operatorname{\sqsubseteq} UPN Etc * and so on.
```

The word written Lic (fL licet=it is allowed that) means "it is allowed (not forbidden) that". Syntactically it behaves like 'Ncs', 'Pot', 'Deb'. Semantically it differs from 'Pot'. 'Ccd' (allows) is both semantically and syntactically different from 'Lic' (it is allowed).

```
3 32 3. * Ha \operatorname{Inq} Ha.

Lie Etc. \leftrightarrow · \neg Deb. \neg Etc:

\neg Lie Etc. \leftrightarrow · Deb. \neg Etc:

\neg Lie. \neg Etc. \leftrightarrow · Deb Etc:

Lie. \neg Etc. \leftrightarrow · \neg Deb Etc. *
```

- 3 32 4. # $Ha \operatorname{Inq} Ha$:

 Lie Etc. \rightarrow $x \operatorname{Vul} \operatorname{Etc.} \in \operatorname{Mal}$ #
- 3 32 5. The receiver will not grasp the meaning of 'Deb' as long as promises (Pol) seem to be the only sources of obligations (Deb). Other instances are indispensable.

The next talk is related to that of 3294.

```
* ^{t} \cdot ^{t}
```

```
Erg: \neg Pot PAN Ha Vul: \lor \lceil t \cdot t' \cdot y \rceil \cdot Hc Inq y \cdot t_1 t_2 Fit \mathfrak{b} |
                 Ha \operatorname{Inq} Hc \cdot \operatorname{Pot} \cdot \operatorname{Qqm} \cdot - \operatorname{Lie} \cdot
                      \neg Lie u \rightarrow \neg Lie Nne Ha Vul u.
                      \wedge: — Lie · Nne Ha Inq Hc . Nne Ha Vul u:
                      Tan: Lie tt' Hc Inq Hy. t_1 t_2 Fit b
                 HcInq Ha -Nnc HcItg: t1 t2 Fit b eFal: Erg Etc: Erg. Lie Etc
                 Ha \operatorname{Inq} Hc \cdot t_1 t_2 \operatorname{Fit} \mathfrak{b} \in \operatorname{Ver} \mathsf{I}
                 Hc Inq Ha: Qqm · PAN Ha Inq Hc. t1 t2 Fit a
                 Ha Inq Hc Deb PAN Ha Inq Hc. t<sub>1</sub> t<sub>2</sub> Fit b:
                      PAN Ha Vul . Nnc Ha Inq Hcb:
                      Tan. PAN Ha Inq Hca:
                      PAN Ha Inq Hc.b. Mod Mal I
                  Hd Inq Hd: Pot · PAN Ha Inq Hc. b. Mod Mal:
                      Tan . - Nes Etc
                      Pot · PAN Ha Vul . Noc Ha Inq Hc a : Qqm . \neg Lie *
3 32 6. # Ha \operatorname{Inq} Hb: \lim x \cdot x \in \operatorname{Pri} A \cdot x > 10^{10010}.
                 Hb Inq Ha · Cur . Ha Inq Hb:
                  Ha\operatorname{Inq} Hb 'Qia: - Nnc Ha\operatorname{Sei} . Tan · Nnc Ha\operatorname{Vul} . Nnc Ha\operatorname{Sei} .
                 Hb \operatorname{Inq} Hc: \operatorname{?} \operatorname{lum} x \cdot x \in \operatorname{Pri} \cdot \wedge \cdot x > 10^{1001}:
                 Hc \operatorname{Inq} Hb 1000001001:
                 Hb \operatorname{Inq} Ha 1000001001:
                 Hd Inq Hd · Lie . PAN Hb Inq Hc Etc *
3 32 7. # Ha \operatorname{Inq} Hb: ? 1^{\operatorname{um}} x \cdot x \in \operatorname{Pri} A \cdot x > 10^{1001}:
                  Hb\operatorname{Inq} Hc: \operatorname{?lum} x \cdot x \in \operatorname{Pri} \cdot \land \cdot x > 10^{1001}
                  Ha\operatorname{Inq} Hc 'Nne Ha\operatorname{Vul}. \operatorname{\neg}\operatorname{Nne} Hc\operatorname{Rsp} Hb.
                       \neg Lie . PAN Hb Inq Hc Etc:
                  Hc\operatorname{Inq} Ha: Utr · PAN Ha\operatorname{Sei} . ? 1^{\operatorname{um}} Ete :
                  {\it Ha}\,{
m Ing}\,{\it Hc}\,{
m Ver} :
                  Hc \operatorname{Inq} Ha \operatorname{Cur} \cdot \operatorname{PAN} Ha \operatorname{Inq} Hb \operatorname{Ete}
                  Ha \operatorname{Ing} Hc : \operatorname{PAN} Ha \operatorname{Vul} \cdot \operatorname{Nne} Ha \operatorname{Sei} \cdot \operatorname{?lum} \operatorname{Etc} :
                      Sed 'PAN Ha Vul: Nnc Ha Sei · Utr Nnc Hb Sei · ? 1um Etc :
                      \operatorname{Erg} \cdot - \operatorname{Lie} \cdot \operatorname{PAN} Hb \operatorname{Inq} Hc \operatorname{Ete} *
3 32 8. # Ha \operatorname{Inq} Hb : ! \operatorname{1um} x \cdot x \in \operatorname{Pri.} A \cdot x > 10^{1001}
                  Hc \operatorname{Inq} Hb \cdot 1000001001 \cdot \operatorname{Mod} \cdot - Ha \operatorname{Ani}:
                  Hb \operatorname{Inq} Ha 1000001001:
                  Ha Inq Hb Ben
                  Hd Inq Hd 'Nne Hb Vul: Nne Ha Inq Ha · Nne Hb Sei .? 1um Etc:
                           Qqm — Sci :
                  Hd \operatorname{Inq} Hb : \operatorname{Mal} - \operatorname{Lie} : \operatorname{Nne} Hb \operatorname{Vul} \cdot \operatorname{Nne} Ha \operatorname{Inq} Ha \cdot 1^{\epsilon} \operatorname{Fal} *
```

```
3 32 9. (t_1 < t_2 < t_3 < t_4 < t_5)

**{}^{t_1}Ha \operatorname{Inq} Hb \cdot ?x \cdot t_3 t_4 \operatorname{Fit} x \cdot Hb \operatorname{Inq} Ha \cdot ?x \cdot t_3 t_4 \operatorname{Fit} x \cdot Hb \operatorname{Inq} Ha \cdot ?x \cdot t_3 t_4 \operatorname{Fit} x \cdot Ha \operatorname{Inq} Hb \cdot : \operatorname{Nnc} Ha \operatorname{Vul} \cdot : \operatorname{Nnc} Ha \operatorname{Sci} \operatorname{Etc} \cdot PAN Hb \operatorname{Inq} ! Ha \cdot ?y \cdot t_1 t_2 \operatorname{Fit} y \cdot : Qqm \cdot Pau \operatorname{Ant} \cdot PAN \cdot Hb \operatorname{Pol} Hc \cdot : Qqm \cdot Pau \operatorname{Ant} \cdot PAN \cdot Hb \operatorname{Pol} Hc \cdot : Qqm \cdot Pau \operatorname{Ant} \cdot : PAN \cdot Hb \operatorname{Inq} ! Ha \cdot ?y \cdot t_1 t_2 \operatorname{Fit} y \cdot : Qqm \cdot Pau \operatorname{Pol} \cdot : \operatorname{Nnc} Hb \operatorname{Inq} ! Ha \cdot ?x \cdot t_3 t_4 \operatorname{Fit} x \cdot : Qqm \cdot : PPN \cdot Ha \operatorname{Inq} ! Hc \cdot : \operatorname{UAN} Hb \operatorname{Inq} ! Ha \cdot ?y \cdot t_1 t_2 \operatorname{Fit} y \cdot : Qqm \cdot : PPN \cdot Ha \operatorname{Inq} ! Hc \cdot : \operatorname{UAN} Hb \operatorname{Inq} ! Ha \cdot ?y \cdot t_1 t_2 \operatorname{Fit} y \cdot : Qqm \cdot : Pqn \cdot : Qqm \cdot : Qqqm \cdot : Qqqm \cdot : Qqqm \cdot : Qqq \cdot : Qqqq \cdot : Qqq \cdot : Qqq \cdot : Qqqq \cdot : Qqqq \cdot : Qqq \cdot : Qqq \cdot : Qqqq \cdot : Qqq \cdot : Qqqq \cdot : Qqqq \cdot
```

Many examples like this may be constructed in order to explain the use of 'Deb' and 'Lic'.

```
3 33 1. * ^{t_1}Ha \operatorname{Inq} Hb \cdot ? = 10 \times 10^{t_2} |

Hb \operatorname{Inq} Ha 100 \cdot 
Hc \operatorname{Inq} Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x \cdot 
Hb \operatorname{Inq} Hc : t_1 t_2 Ha \operatorname{Vul} \cdot \operatorname{Nne} Ha \operatorname{Sci} \cdot ? = 10 \times 10 \cdot 
Hc \operatorname{Inq} Hb \cdot \operatorname{Fal} \cdot \operatorname{Qia} \cdot t_1 t_2 Ha \operatorname{Sci} \cdot ? = 10 \times 10 \cdot 
- t_1 t_2 Ha \operatorname{Vul} \cdot \operatorname{Nne} Ha \operatorname{Sci} \cdot \operatorname{Etc} \cdot 
\operatorname{Sed} : t_1 t_2 Ha \operatorname{Vul} \cdot \operatorname{Nne} Ha \operatorname{Sci} \cdot \operatorname{Utr} \operatorname{Nne} Hb \operatorname{Sci} \cdot \operatorname{Etc} \cdot 
Hb \operatorname{Inq} Hc \cdot \operatorname{Etc} \cdot \operatorname{Pot} \cdot t_1 t_2 Ha \operatorname{Vul} \operatorname{Sci} \cdot \cdot 
\operatorname{Utr} : t_1 t_2 Hb \operatorname{Ani} \cdot ? y \cdot t_1 t_2 Ha \operatorname{Inq} Hb y \cdot 
Hc \operatorname{Inq} Hb \operatorname{Ben} *
```

From this text and the next ones the receiver may learn how to prove 'Ha Vul Etc' by circumstantial evidence though Ha did not use the word Vul.

```
3 33 2.  (t<sub>1</sub> < t<sub>2</sub> < t<sub>3</sub>)

**t<sub>1</sub> Ha Inq Hb · ? x . t<sub>1</sub> t<sub>2</sub> Fit x t<sub>4</sub> |

Hb Inq Ha a |

Hc Inq Hb · ? y . t<sub>3</sub> t<sub>4</sub> Fit y |

Hb Inq Hc * t<sub>3</sub> t<sub>4</sub> Ha Vul : Nne Ha Sei · ? x . t<sub>1</sub> t<sub>2</sub> Fit x |

Hc Inq Hb : ? u · PAN Hb Cog : Utr * t<sub>3</sub> t<sub>4</sub> Ha Vul :

Nne Ha Sei · ? x . t<sub>1</sub> t<sub>2</sub> Fit x : Mod u |

Hb Inq Hc : t<sub>3</sub> t<sub>4</sub> Ha Inq Hb · ? x . t<sub>1</sub> t<sub>2</sub> Fit x |

Hc Inq Hb · ¬ Nes . t<sub>3</sub> t<sub>4</sub> Ha Vul Ete :

Sed · Pot : t<sub>3</sub> t<sub>4</sub> Ha Sei Ete . ^ t<sub>3</sub> t<sub>4</sub> Ha Vul :

Nne Ha Sei · Utr . Nne Hb Sei Ete *
```

3 33 3. $\# t_1Ha \operatorname{Inq} Hb Hb^{t_2}$: $Hb \operatorname{Inq} Ha : \operatorname{Nnc} Hb \operatorname{Ani} \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x$:

```
Hc Inq Hb · ? x · t<sub>1</sub> t<sub>2</sub> Fit x :

Hb Inq Hc · t<sub>1</sub> t<sub>2</sub> Ha Vul : Nnc Ha Sci ·

Utr : Nnc Hb Ani · ? x · t<sub>1</sub> t<sub>2</sub> Fit x :

Hc Inq Hb · ? u : PAN Hb Cog · ? p · t<sub>1</sub> t<sub>2</sub> Ha Vul p : Mod u :

Hb Inq Hc · t<sub>1</sub> t<sub>2</sub> Ha Inq Hb Hb :

Hc Inq Hb Ben *
```

3 33 4. ** t₁Ha Inq Hb: Utr·Nnc Hb Vul. Nnc Hb Inq Hd p t₂ l

Hb Inq Ha Ver l

Hb Inq Hd p l

Hc Inq Hb·? x. t₁t₂ Fit x l

Hb Inq Hc·t₁t₂ Ha Vul. PPN Hb Inq Hd p l

Hc Inq Hb. Cur Sci l

Hb Inq Hc: t₁t₂ Ha Inq Hb: Utr·Nnc Hb Vul. Nnc Hb Inq Hd p

Erg·t₁t₂ Ha Vul. PPN Hb Inq Hd p l

Hc Inq Hb: Cur: — t₁t₂ Ha Inq Hb. Nnc Ha Vul Etc

Sed: t₁t₂ Ha Inq Hb·Utr. Hb Vul Etc l

Hc Inq Hc:. Pot | Pau Ant. t₁ Ha Inq Ha:

Pot: t: Ha Inq Hb. Nnc Ha Vul r

Pau Pst·Hb Inq Ha. — Nnc Ha Vul r

* Ha Inq Hb·Utr. Nnc Hb Vul r:

 \rightarrow Pau Pst · Hb Ing Ha. Nnc Ha Vul r *

3 34 1. Continuation of the last talk:

```
* Hb Inq Hc i. Qia ! Qqm' Lic:

tt' Ha Inq Hb · Nnc Ha Vul . Nnc Hb Etc:

Tan : Dec 'tt' Ha Inq Hb : Utr · Hb Vul . Nnc Hb Etc i

tt' Ha Vul . Nnc Hb Etc ·

⇒ : Dec 'Ha Inq Hb : Utr · Hb Vul . Nnc Hb Etc i

¬ Deb . Sed Dec I

Hc Inq Hb i Utr : Dec : Eti : Ha Inq Hb ·

Utr : Nnc Hb Ccd · Nnc Ha Vul . Nnc Hb Etc I

Hb Inq Hc : Sat : Ha Inq Hb Utr · Nnc Hb Vul . Nnc Hb Etc :

Tan : Plt : Ha Inq Hb ·
```

Utr: Nnc Hb Ccd · Nnc Ha Vul · Nnc Hb Etc *

The words written Dec (fL decet) and Plt (fL placet) mean modalities of decency. I have not succeeded in finding unambiguous translations of these words. Perhaps "it is courteous" (French "courtois", German "höflich") for 'Plt', and "it is convenient" (French "il convient", German "es schickt sich") for 'Dec' will do. The relationship between the two levels of decency 'Plt' and 'Dec' is akin to that between the two levels of duty, 'Deb', and 'Lic',

and to that between the two levels of necessity, 'Nes' and 'Pot'.

The receiver will learn from the last talk that there may be some reason of uttering one's wishes in a more circumstantial way, not for instance saying "I wish that you..." but rather "do you wish to..." or even "do you allow me to ask you whether you...". As early as in 3 29 7 one of the partners has displayed this more polite behaviour. A comment analogous to that on 3 33 4 could be added to 3 29 7.

Though we have still little use for refined behaviour modalities like 'Plt' and 'Dec', a few more texts might be welcome.

```
3 34 2. ** ^{t_a}Ha \operatorname{Inq} Hb `Utr: Nne Hb \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x^{t_a} |
Hb \operatorname{Inq} Ha \cdot t_1 t_2 \operatorname{Fit} \alpha^{t_a} |
Hc \operatorname{Inq} Hb `Cur: - t_4 t_5 Hb \operatorname{Inq} Ha \cdot \alpha \cdot \operatorname{Sed} \cdot t_1 t_2 \operatorname{Fit} \alpha |
Hb \operatorname{Inq} Hc. Qia Plt |
Hc \operatorname{Inq} Hb `Cur': - t_3 t_4 Ha \operatorname{Inq} Hb :
?x \cdot t_1 t_2 \operatorname{Fit} x `Sed `Utr: Nne Hb \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x |
Hb \operatorname{Inq} Hc \cdot : Qia Dec \cdot \cdot \cdot \operatorname{Eti} \cdot \cdot \cdot \operatorname{Plt} \cdot \cdot
t \cdot \cdot Ha \operatorname{Vul} : Nne Ha \operatorname{Sei} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x \cdot \cdot
\rightarrow \cdot \cdot \cdot Ha \operatorname{Inq} Hb : Utr `Nne Hb \operatorname{Ced} :
Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x \cdot +
```

3 34 3. Likewise one can treat talks such as:

```
# Ha \operatorname{Inq} Hb 'Utr: Nne Hb \operatorname{Sei} ? x \cdot t_1 t_2 \operatorname{Fit} x:
Hb \operatorname{Inq} Ha 'Utr: Nne Ha \operatorname{Ced} '
Nne Hb \operatorname{Inq} Ha : \longrightarrow \operatorname{Nne} Hb \operatorname{Sei} ? x \cdot t_1 t_2 \operatorname{Fit} x #
```

and

```
# Ha \operatorname{Inq} Hb 'Utr: Nne Hb \operatorname{Sei} ? x \cdot t_1 t_2 \operatorname{Fit} x!

Hb \operatorname{Inq} Ha: Ver. Tan: Utr' Nne Ha \operatorname{Ced}:

— Nne Hb \operatorname{Inq} ! Ha \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x #
```

3 34 4. * $Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit}$ $Hb \operatorname{Inq} Ha a$ $Ha \operatorname{Inq} Hb x$ $Hc \operatorname{Inq} Hb : \neg \operatorname{Dec} \cdot Hb \operatorname{Inq} Ha a \cdot \operatorname{Ant} \cdot Ha \operatorname{Inq} Hb x *$

```
3 34 5. # Ha \operatorname{Inq} Hb \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x:

Hc \operatorname{Inq} Ha a:

Hd \operatorname{Inq} Hc \cdot \operatorname{Dec} \cdot t : Ha \operatorname{Inq} Hb \cdot ? \operatorname{Etc} \cdot

\rightarrow \operatorname{Pau} \operatorname{Pst} \cdot Hc \operatorname{Ced} \cdot \operatorname{Nnc} Hb \operatorname{Rsp} Ha #
```

```
3 34 6. # Ha Inq Hb p |

Hc Inq Hd 'Utr: Nnc Hd Sci · ? x. PAN Ha Inq Hb x |

Hd Inq Hc · ¬ Sci |

Hc Inq Hd · Utr · ¬ PAN Hd Ani Etc |

Hd Inq Hc Ver |

Hc Inq Hd: Utr · ¬ Pot . PAN Hd Inq Etc |

Hd Inq Hc! Pot . Tan . ¬ PAN Vul!

¬ PAN Hd Vul . Nnc Hd Ani Etc |

Qia : ∀x' PAN Ha Inq Hb: x: Mod · Ha Vul . ¬ Nnc Hd Ani Etc |

Erg 'Plt · PAN Hd Vul . ¬ Nnc Hd Ani Etc |

A: Dec · ¬ PAN Hd Vul . Nnc Hd Ani Etc #
```

3 34 7. # $Ha \operatorname{Inq} Hb \mathfrak{p} I$

```
Hc \operatorname{Inq} Hd 'Utr: Nne Hd \operatorname{Sei} ? x. PAN Ha \operatorname{Inq} Hb x |
Hd \operatorname{Inq} Hc : \operatorname{Qqm} \cdot \operatorname{Nne} Hd \operatorname{Sei} :

\operatorname{Tan} : \operatorname{Nne} Hd \operatorname{Vul} \cdot \neg : \operatorname{Unq} \operatorname{Pst} \operatorname{PAN} : Hd \operatorname{Sei} :

\operatorname{Qia} : \forall x \cdot \operatorname{PAN} Ha \operatorname{Inq} Hb : x : \operatorname{Mod} \cdot Ha \operatorname{Vul} : \neg \operatorname{Nne} Hd \operatorname{Ani} :

\neg \operatorname{Dec} : \operatorname{Nne} Hd \operatorname{Inq} ! Hc \cdot ? x : Ha \operatorname{Inq} Hb x *
```

3348. # Ha Inq Ha

```
Dec Etc. \leftrightarrow · ¬ Plt. ¬ Etc: ¬ Dec Etc. \leftrightarrow · Plt. ¬ Etc: ¬ Dec. ¬ Etc. \leftrightarrow · Plt Etc: Dec. ¬ Etc. \leftrightarrow · ¬ Plt Etc.
```

3 34 9. $# Ha \operatorname{Inq} Ha$

```
Nes Etc. \rightarrow. Deb Etc. \rightarrow. Plt Etc. \rightarrow. Dec Etc. \rightarrow. Lie Etc. \rightarrow. Pot Etc. ^*
```

Perhaps we had better give incidental comments on former texts instead of this general behaviour rule which may become discredited when 'Etc' is a proposition such that e.g. Ncs Etc is true whereas Deb Etc is void of sense, and so on.

- 3 35 1. I have tried to design a vocabulary of modalities. It is still far from satisfactory. I found little help in the literature on this topic. I believe that modalities cannot be treated successfully, if not linked to human behaviour.
- 3 35 2. Modalities allude to some constraint. This constraint may be logical or physical (Ncs), legal or moral (Deb), it may be for decency's sake (Plt), and finally the constraint may be exercised by some person saying "I wish" (Vul). Even in the first three cases the author may be personified: Logic and nature (Ncs), Law (Deb), Decency (Plt).

3 35 3. In the simplest way such a modality is used when somebody justifies an action of his or other people by saying: I did it because I had to. As early as 3 12 2 Ha acted in this way. Compare also 3 29 4 and following texts, 3 34 1 and following texts. The motivation may be necessity, duty, decency or the will of another person.

Note that in our terminology Ncs need not be prohibitive.

Likewise somebody may justify the fact that he or other people abstain from some action by saying that there was a constraint to abstain. See 3 22 6, 3 22 7, 3 23 1, 3 23 3, 3 32 5, 3 32 7, 3 32 8, 3 29 4, 3 30 2 and so on.

Finally some action or the lack of some action may be justified by saying that there was no constraint to commit it or to abstain from it. See 3 22 3, 3 22 4, 3 22 9 and many others.

3 35 4. In a more complicated context modalities are used if somebody is accused of having acted against a constraint or of not having acted in spite of a constraint.

As this cannot be done with impunity, punishments should be shown. If in $3\ 12\ 2$ Ha omits ' $p\in Pri$ ' the punishment will be: falsehood of his assertion. Physical constraints cannot be dealt with, because we do not know anything about the bodies of the acting persons; they are sheer spirits that cannot be punished by physical means.

A man who kills another man in self-defence, acts under the constraint of self-preservation. This constraint is not prohibitive. He could have done otherwise, but then he would have been punished with the loss of his own life. In our system there is no logical implication 'Nes $p cdot \to p$ '.

Punishments against moral transgressors will be treated in 3 36 1.

Justifications and accusations may be combined in the same text.

- 3 35 5. In all these cases the actual event p (or its non-actualization) has been the starting point. It has either been justified or condemned. Yet the inverse process is also possible. The event p (or its non-actualization) is proved or forecast by the argument that there is some constraint to provoke or to prevent it. See 3 22 1, 3 29 4 and many others.
- 3 35 6. 'Pot' is also used to designate some disposition. From potentiality as absence of certain constraints the stress is shifted to potentiality as the presence of certain faculties. If there are no constraints to prevent somebody from committing a certain action, nothing but

his own will is needed for this action to take place. His own will is counter-balanced by his own potentiality. As will is an activity, ability becomes an activity too.

In Lincos syntax we have not accounted for this semantic change from an impersonal "it is possible" to a personal "he can" by any corresponding syntactical device. See 3 23 1, 3, 3 29 8.

3 35 7. Naive physics uses human behaviour-patterns for the understanding of nature. (The sun will not transgress its measures; otherwise the Erinys, Justice's bailiffs, will hunt it down -Herakleitos.) For centuries Aristotle was the philosopher of anthropomorphic physics. Potentiality, will, activity, and passivity are key notions in his natural philosophy. In spite of the law of action and reaction, Newtonian forces (when really applied) still have many features in common with constraints and manifestations of will in human behaviour. Anthropomorphic physics as a first phase of men's intellectual development cannot be discarded. Even in the laboratory, and wherever human actions influence extrahuman nature or are influenced by it, we cannot dispense with an anthropomorphic terminology. When we say that the wind moves the branches of a tree, or can move them, we divide nature into a passive background of potentiality and a causa efficiens that actualizes some potentiality as individual will actualizes the potentialities of the individual. By overstressing this feature we can construct and understand machines. Here the points where the causae efficientes apply are sharply marked by taps, switches, handles and so on.

Physics has not yet been a subject of our texts, but when we turn to it, we shall apply human behaviour terminology in essentially the same way as we do in daily life.

```
3 36 1. * Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} a \operatorname{I}
Hc \operatorname{Inq} Ha : \operatorname{Cur} \cdot \operatorname{PAN} Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} a \operatorname{I}
Ha \operatorname{Inq} Hc : \operatorname{Qia} \cdot \operatorname{Pau} \operatorname{Ant} \cdot \operatorname{PAN} \cdot Ha \operatorname{Pol} Hb :
\operatorname{PPN} Ha \operatorname{Inq} ! Hb \cdot ? x \cdot t_1 t_2 \operatorname{Fit} x : t_1 t_2 \operatorname{Fit} a :
\operatorname{Erg} : \operatorname{Deb} \cdot \operatorname{PAN} Ha \operatorname{Inq} Hb \cdot t_1 t_2 \operatorname{Fit} a :
\operatorname{Deb} p \cdot \rightarrow \cdot Ha \operatorname{Vul} p \cdot \in \operatorname{Ben} :
\operatorname{Deb} \cdot \neg p \cdot \rightarrow \cdot Ha \operatorname{Vul} p \cdot \in \operatorname{Mal} \operatorname{I}
\operatorname{Hc} \operatorname{Inq} Ha \cdot ? q : t : Ha \operatorname{Vul} : 1^{\mathfrak{e}} \cdot \uparrow p \cdot \operatorname{Deb} p \cdot \rightarrow \operatorname{Pau} \operatorname{Pst} \cdot \operatorname{Fit} q :
? r : t : Ha \operatorname{Vul} : 1^{\mathfrak{e}} \cdot \uparrow p \cdot \operatorname{Deb} \neg p \cdot \rightarrow \operatorname{Pau} \operatorname{Pst} \cdot \operatorname{Fit} r \operatorname{I}
\operatorname{Ha} \operatorname{Inq} Hc : q = \cdot \wedge \lceil x \cdot y \rceil : x \wedge y \cdot \in \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Ben} :
r = \cdot \wedge \lceil x \cdot y \rceil : x \wedge y \cdot \in \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Ben} \operatorname{Inq} \operatorname{Ho} \cdot \wedge p : \operatorname{Nnc} \cdot \operatorname{Ha} \operatorname{Vul} p \cdot \rightarrow \operatorname{Pau} \operatorname{Pst} \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Ben} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Ise} \operatorname{Inq} \operatorname{Hom} \operatorname{Inq} \operatorname{Hom} \cdot \rightarrow \cdot xy \operatorname{Inq} \cdot y \operatorname{Inq} \cdot
```

```
Hc \operatorname{Inq} Ha \cdot \operatorname{Utr} : \operatorname{Etc} : t : Ha \operatorname{Vul} : 1^{\mathfrak{c}} \cdot \operatorname{Deb} p
\rightarrow \operatorname{Pau} \operatorname{Pst} \cdot Hi Hj \cdot \operatorname{Ise} \operatorname{Inq} \operatorname{Ben} !
Ha \operatorname{Inq} Hc \cdot \cdot \operatorname{Pot} \cdot Hi Hj \cdot \operatorname{Ise} \operatorname{Inq} \operatorname{Mal} \cdot \cdot 
\vee x : x \in \operatorname{Hom} \wedge \wedge \wedge \wedge t \cdot t \cdot p : tt' x \operatorname{Vul} p \cdot \rightarrow \cdot \operatorname{Deb} \cdot \neg p : \wedge :
\vee x : x \in \operatorname{Hom} \wedge \vee \wedge \wedge t \cdot t \cdot p : tt' x \operatorname{Vul} p \cdot \wedge \cdot \operatorname{Deb} \cdot \neg p : \wedge :
\vee x : x \in \operatorname{Hom} \wedge \wedge \wedge \wedge t \cdot t \cdot p : tt' x \operatorname{Vul} p \cdot \wedge \cdot \operatorname{Deb} p : \wedge :
\vee x : x \in \operatorname{Hom} \wedge \wedge \wedge \wedge t \cdot t \cdot p : tt' x \operatorname{Vul} p \cdot \rightarrow \cdot \operatorname{Deb} p :
A = : ^{\star} x : x \in \operatorname{Hom} \wedge \wedge \wedge \wedge t \cdot t \cdot p : tt' x \operatorname{Vul} p \cdot \rightarrow \cdot \operatorname{Deb} p :
\rightarrow : t'' : u \operatorname{Vul} : 1^{\mathfrak{c}} \cdot \wedge q \cdot \operatorname{Deb} q : \rightarrow \operatorname{Pau} \operatorname{Pst} \cdot \operatorname{Mul} A \cdot \operatorname{Ise} \operatorname{Inq} \operatorname{Ben} : \wedge : t'' : u \operatorname{Vul} : 1^{\mathfrak{c}} \cdot \wedge q \cdot \operatorname{Deb} \cdot \neg q : \rightarrow \operatorname{Pau} \operatorname{Pst} \cdot \operatorname{Mul} A \cdot \operatorname{Ise} \operatorname{Inq} \operatorname{Mal} *
\operatorname{Ise} \operatorname{Inq} \operatorname{Mal} *
```

Clearly this is childish reasoning. But it is no more childish than any conversation with children about fundamentals of ethics or than the categorical imperative when translated into a language lacking the equivalents of "categoric", "imperative", "maxim", "principle". It is also no more childish than the introduction of natural numbers in the first chapter and of time in the second. We may subsequently reformulate the definition of natural number by means of Peano's axioms, though it is not yet certain that Peano's numbers are superior to simple counters. People have sound ideas about practical ethics because they once knew its principles in childish terms. On a higher linguistic level more sophisticated conversations about ethics will be possible, and if there be any rational system for ethics like Peano's axioms for arithmetics, we may try to translate it. At the present stage of our project this would be premature.

After all, the last talk is not as childish as it looks. Ha defines the set of good people. It is those who wish only good things to happen. Then Ha asserts that good people call good actions good and bad actions bad. If we apply to this their own actions, we may conclude that good peoples are committed to those actions they call good themselves. There is some sense in this assertion, though Ha may be wrong in saying that the set of people like this is not void. It is a criterion of goodness though it is not Kant's categorical imperative. It looks more like Shakespeare's

This above all: to thine own self be true. And it must follow, as the night the day, Thou canst not then be false to any man,

but this is a mere accident. I have made no attempt to inquire into the fundamentals of ethics.

Of course the ethical principle stated by Ha is liable to many

variations, e.g. by using Lic instead of Deb or by inverting implication arrows. We do not enter into details.

A conversation about decency could be carried on in a similar way.

3 36 2. When reading the printer's proof, I remark that the analysis in 3 35 e has been too limited. The levels of modality cannot be localized by punishments only. We have had to use rewards too. 'Deb' and 'Plt' have been distinguished from 'Lic' and 'Dec' by means of their relation to the judgements "praiseworthy" (Ben) and "blameless" (¬Mal).

```
3 37 1. # Ha \operatorname{Inq} Hb \cdot ? = 10 \times 10 I
                Hb \operatorname{Inq} Ha 100 I
                Ha \text{ Inq } Hb .? = 1010111 \times 100101 I
                Hb\operatorname{Inq} Ha. \neg Scil
                Ha Inq Hb : Cur ' Nne Hb Sei:
                    Qqm \cdot ? = 10 \times 10 \cdot Tan - ? = 1010111 \times 100101 I
                Hb \text{ Inq } Ha : \text{Qia 'Dif 'Nnc } Hb \text{ Pet . } ? = 1010111 \times 100101 :
                    >: Dif · Nnc Hb Pet . ? = 10 \times 10
                Hc \operatorname{Inq} Ha : \operatorname{Pot} \cdot \operatorname{Nnc} x \operatorname{Rep} \cdot ? = 1010111 \times 100101 :
                    Qqm \rightarrow .Tan . \rightarrow :Pot \cdot Nnc x Rep .? = 10 \times 10 I
                Hd \operatorname{Inq} Ha: = \operatorname{Dif} I
                Ha \operatorname{Inq} Hd \cdot \operatorname{Exg} \cdot \operatorname{Dif} \cdot ?x \cdot x^{10} - 101x + 110 = 0:
                    > : Dif.? = 1010111 \times 100101 I
                \rightarrow: Dif. ? x Etc. < Dif. ? = 1010111 × 100101 |
                Ha \operatorname{Ing} Hc \operatorname{Tan} : \neg \operatorname{Nne} y \operatorname{Sei} \operatorname{Etc} \rightarrow \operatorname{Dif} \operatorname{Etc} > \operatorname{Dif} \operatorname{Etc} 
                Hc Inq Ha Ver #
```

The word written Dif (fL difficultas=difficulty) means "difficulty (of...)". It might be considered as a modality quite different from the modalities treated above. The syntax of 'Dif' is clear from its occurrence in Hb's third answer. Later on we use an abbreviated notation that can be explained by some behaviour rule.

```
3 37 2. * Ha \operatorname{Inq} Hb : \operatorname{Utr} : \wedge n \cdot n \in \operatorname{Num}.

\rightarrow : \vee \lceil p \cdot m \rceil \cdot p \in \operatorname{Pri} . \wedge . m \in \operatorname{Num} . \wedge . p = 100 \ m - 1 . \wedge . p > n \mid
Hb \operatorname{Inq} Ha : \operatorname{PPN} . \operatorname{Qqm} . \neg \operatorname{Nnc} : Hb \operatorname{Sci} . \operatorname{Utr} \operatorname{Etc} :
\operatorname{Qia} : \operatorname{Pot} : \operatorname{Nnc} Hb \operatorname{Dem} :
\wedge n \cdot n \in \operatorname{Num} . \rightarrow : \vee p \cdot p \in \operatorname{Pri} . \wedge . p > n :
\operatorname{Tan} : \operatorname{Dif} : \operatorname{Utr} : \wedge n \cdot n \in \operatorname{Num} . \rightarrow : \vee r p \cdot m \cap \operatorname{Etc} :
\Rightarrow : \operatorname{Dif} : \operatorname{Utr} : \wedge n \cdot n \in \operatorname{Num} . \rightarrow : \vee p \operatorname{Etc} :
```

More examples may be added, e.g. comparisions of the difficulty of

Utr: $\wedge n$ $n \in \text{Num}$.

$$\rightarrow$$
: $\forall \lceil p \cdot m \rceil \cdot p \in \text{Pri.} \land . m \in \text{Num.} \land . p = 100 m + 1. \land . p > n #$

* Utr: $\wedge \lceil a \cdot b \cdot n \rceil$:

$$a \wedge b \wedge n \in \text{Num} \wedge \wedge q : q \text{ Div } a \wedge q \text{ Div } b \rightarrow q = 1$$
:
 $\Rightarrow : \forall \lceil p \cdot m \rceil \cdot p \in \text{Pri} \wedge m \in \text{Num} \wedge p = am + b \wedge n = n = n$

3 37 3. A comment on 3 27 1:

$$* \wedge \lceil tt'x \rceil$$
: Dif $tt'x \operatorname{Pet} \operatorname{Utr} G \cdot > \operatorname{Vld} *$

The word written Vld (fL valde=very) is akin to 'Nim' (see 3 24 4). '>Vld' may be translated "very large", '<Vld' "very small".

Dif $\cdot tt' x \operatorname{Pet} \cdot ? \operatorname{Etc} :> \operatorname{Vld} \to \operatorname{Cea} := \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Pot} \cdot tt' x \operatorname{Rep} \cdot ? \operatorname{Etc} = \operatorname{Pot} \cdot tt' x \operatorname{Pot} \cdot tt' x \operatorname{Pot} = \operatorname{Pot} \cdot tt' x \operatorname{Pot} = \operatorname{Pot} \cdot tt' x \operatorname{Pot} \cdot tt' x \operatorname{Pot} = \operatorname{Pot} \cdot tt' x \operatorname{Pot} =$

3 37 4. A comment on 3 37 1:

$$\wedge \lceil t \cdot t' \cdot x \rceil$$
 Dif $t' \cdot x$ Pet $? = 10 \times 10 : < \text{Vld}$ #

* Dif •
$$tt'x$$
 Pet .? Etc : $<$ Vld • \Leftrightarrow • Cca : Ncs • $tt'x$ Rep .? Etc *

3 37 5. There is a great variety of contexts in which 'Dif' can appear.

Let p, q, r be sequences of 4, 32 and 256 short signals respectively within a relatively short time.

```
# t, p t, a t, r t,
```

```
Ha \operatorname{Inq} Hb \cdot \operatorname{Dif} \cdot ?x \cdot t_1 t_2 x^{es} \operatorname{Fit} \cdot :< \operatorname{Vld}:
```

$$\operatorname{Dif} \cdot ?x \cdot t_1 t_2 x^{es} \operatorname{Fit} \bullet : < : \operatorname{Dif} \cdot ?y \cdot t_2 t_3 y^{es} \operatorname{Fit} \bullet :$$

$$<$$
: Dif \cdot ? z . $t_3 t_4 z^{es}$ Fit \bullet :

$$\mathrm{Dif} \cdot ?z \cdot t_3 t_4 z^{es} \, \mathrm{Fit} \bullet :> \mathrm{Vld} \, \mathrm{'Qia} : \, ^{\vee}z \cdot t_3 t_4 z^{es} \, \mathrm{Fit} \, \bullet \cdot \mathrm{Mul} > .1 \, *$$

$Ha \operatorname{Inq} Hb \longrightarrow \operatorname{Nnc} Ha \operatorname{Sci} \cdot ?x \cdot t_1 t_2 \operatorname{Fit}$:

$$\wedge: \neg$$
 Nnc Ha Sci $? y \cdot t_3 t_4$ Fit $y \cdot$

$$Hb \operatorname{Inq} Ha$$
 'Nnc. Pau 'Hom. Sci $?x \cdot t_1 t_2 \operatorname{Fit} x$:

Tan: Nnc.
$$Mul^{\epsilon}$$
 Hom. $Sei \cdot ? y \cdot t_1 t_2$ Fit y .

$$Ha \operatorname{Inq} Hb : \operatorname{Erg} : \operatorname{Dif} : \operatorname{Nnc} Ha \operatorname{Pet} \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x \operatorname{Etc}$$

$$\gt$$
 Dif: Nnc Ha Pet $\mathrel{?} y \cdot t_1 t_2$ Fit y Etc $*$

$$+ Ha \operatorname{Inq} Hb$$
 ? $x: t_1 t_2 x \operatorname{Ani} \cdot ? y \cdot t_1 t_2 \operatorname{Fit} y$:

$$Hb\operatorname{Inq} Ha \cdot \neg \lor x \cdot \mathsf{t}_1 \, \mathsf{t}_2 \, x \operatorname{Ani} \operatorname{Etc}$$
:

$$Ha \operatorname{Inq} Hb : \operatorname{Erg} \operatorname{Dif} \cdot ?y \cdot t_1 t_2 \operatorname{Fit} y :> \operatorname{Vld} :$$

$$Ha \operatorname{Inq} Hb$$
 ? $x : t_3 t_4 x \operatorname{Ani} \cdot ? y : t_3 t_4 \operatorname{Fit} y$:

$$Hb \operatorname{Inq} Ha : t_3 t_4 x \operatorname{Ani} : y \cdot t_3 t_4 \operatorname{Fit} y : \leftrightarrow \cdot x = Hc$$
:

$$\operatorname{Tan}: \wedge \lceil t \cdot t' \cdot z \rceil$$
 Nne $\operatorname{He}\operatorname{Vul}: \neg tt' \operatorname{He}\operatorname{Inq}!z \cdot ?x \cdot t_3 t_4 \operatorname{Fit} x$:

$$Ha \operatorname{Inq} Hb : \operatorname{Dif} : \operatorname{Nnc} Ha \operatorname{Cog} \cdot ?x \cdot \mathsf{t}_1 \mathsf{t}_2 \operatorname{Fit} x$$

> Dif: Nne $Ha \operatorname{Cog} \cdot ? x \cdot t_3 t_4 \operatorname{Fit} x$:

Tan: Dif: Nnc $Ha \operatorname{Cog} \cdot ?x \cdot t_3 t_4 \operatorname{Fit} x \cdot > \operatorname{Vld} *$

```
3\ 37\ 6.\ * {}^{t_1}Ha\ \text{Inq}\ Hb\ p^{t_2}
                    Ha Ani · Utr · Pot Hc Ani l
                    Ha \operatorname{Inq} Hc \operatorname{Utr} : \operatorname{PAN} Hc \operatorname{Ani} \cdot ?x . \operatorname{Nnc} Ha \operatorname{Inq} Hb x I
                    Hc \operatorname{Inq} Ha \operatorname{Ver} \mathbf{I}
                    Ha Inq Hc I Utr : Nne Hc Ced : Nne Ha Vul : Nne Hc Vul :
                         \neg: Unq Pst PPN . Hc Sci \cdot? x \cdot t_1 t_2 Fit x \mid
                    Hc \operatorname{Inq} Ha : \operatorname{Qqm} . \operatorname{Nnc} Hb \operatorname{Vul} : \operatorname{Tan} : - \operatorname{Nnc} Hb \operatorname{Sei} :
                             Utr: Pot \neg: Unq Pst PPN. Hc Sci ?x.t_1t_2 Fit x:
                         Qia \land \land y : y \in \text{Hom} \rightarrow \land \land
                             Qqm \cdot \operatorname{Car}^{: \uparrow} x : \forall t \ t \ y \operatorname{Cog} x . \land :
                                       \neg Unq. Pau Pst. t \cdot y Sci x \ge V \text{ld}...
                                  Tan \forall \lceil x \cdot t \rceil \exists t y \operatorname{Cog} x \cdot \wedge \exists
                                       \wedge t' t < t' < \cdot Fin. y \operatorname{Ext} : t' y \operatorname{Sci} x:
                        Exg: t Hb \cos x \cdot \wedge \cdot t' Hb \cos x'
                             \wedge : v = \cdot \neg \text{Unq Pst PPN} \cdot Ha \text{Sci } x :
                              \wedge : v' = \cdot - Unq Pst PPN \cdot Ha Sci x' :
                              \wedge: Pau Pst. t \cdot Hb Vul v:
                              \wedge : \neg \operatorname{Pau} \operatorname{Pst} \cdot t' \cdot Hb \operatorname{Vul} v'
                              \rightarrow Dif v Mul > Dif v' *
```

3 38 0. Our treatment of the word written Vul is still unsatisfactory. The wishes uttered by the acting persons are rather whims, and even if a motivation is given, it is no less "whimsical". Punishments have been displayed, but I am not sure whether they are persuasive. So I do not believe that the receiver will really grasp the meaning of 'Vul'.

Stories usually start with a wish of one of the acting persons. As long as we do not know his past, we can hardly appreciate his wish. One could remove this difficulty by telling longer stories or rather by giving a text that is one connected story. However, in this tentative stage of Lincos construction I do not feel prepared for such a task. I have therefore renounced any attempt to show diversity of human individuals, though not for reasons of principle.

We can compensate this want of personality by a more solid general background of human wishes. This will be done in the remaining paragraphs of this chapter.

```
3 38 1. Ha \operatorname{Inq} Hb : \operatorname{Nnc} Ha \operatorname{Vul} : \operatorname{PPN} Ha \operatorname{Inq} Hb . 1^{\mathfrak{e}} \operatorname{Qus} \cdot \\ \wedge : \operatorname{Pau} \operatorname{Pst} . \operatorname{PPN} : Hb \operatorname{Rsp} \operatorname{Mod} \operatorname{Ben} . \\ \rightarrow \operatorname{Pau} \operatorname{Pst} : Ha \operatorname{Dat} Hb . \operatorname{Den} 10000 \\ : \wedge : \operatorname{Pau} \operatorname{Pst} : \operatorname{PPN} : Hb \operatorname{Rsp} \operatorname{Mod} \operatorname{Mal} . \\ \rightarrow \operatorname{Pau} \operatorname{Pst} : Hb \operatorname{Dat} Ha . \operatorname{Den} 10000 \mathsf{I}
```

```
Hb \operatorname{Ing} Ha \cdot \operatorname{Nnc} Hb \operatorname{Vul} \ \ \ \ \ \operatorname{PAN} Ha \operatorname{Vul} w 
Ha Inq Hb \cdot ?x \cdot x^{11} - 11x - 1000,001 = 0
Hb \operatorname{Ing} Hb \cdot \operatorname{Dif} > \operatorname{Vld} \cdot \operatorname{Qia} - \operatorname{Nnc} Hb \operatorname{Sci}
     ? \forall y : y \in Anl^{io} : 1 \in Cpu^{io} : ? Etc
Hb \operatorname{Ing} Hc : \operatorname{Nnc} Hb \operatorname{Vul} : \operatorname{PPN} Hb \operatorname{Sci}
         ?x.x^{11}-11x-1000.001=0
     Nnc: Hc \operatorname{Inq} ! Hb. ? Etc \rightarrow Pau Pst \cdot Hb \operatorname{Dat} Hc. Den 1000 |
Hc Hd He Ise Inq I
Hc\operatorname{Inq} Hb: \neg \lor y \cdot y \in \operatorname{Hom} . \land .\operatorname{Nnc} y\operatorname{Sci} \operatorname{Etc} \mathsf{I}
Hb\operatorname{Ing} Ha . 
eg\operatorname{Nnc} Hb\operatorname{Sei} \mathsf{I}
Ha\operatorname{Inq} Hb 10,1 lacksquare
Hb Inq Ha Mal I
Ha \operatorname{Inq} Hb \cdot 10,1^{11} - 11 \times 10,1 - 1000,001 = 0
Hb \operatorname{Ing} Ha \operatorname{Ver} \cdot \operatorname{Tan} \cdot - \operatorname{Sat}
     Deb: \forall u \cdot u \in Cpu^{lo} Etc. \land . Nnc Ha Inq Hb u
Ha \operatorname{Ing} Hb \cdot \operatorname{Nnc} Ha \operatorname{Sci} : ? = u \cdot u \in \operatorname{Cpu}^{lo} \operatorname{Etc}
     Tan '- Nne Ha Vul: Nne Ha Inq! Hb ? u \in Cpu^{io} Etc!
Hb \operatorname{Ing} Ha = \operatorname{Deb} \cdot \operatorname{Nnc} Hb \operatorname{Dat} Ha \cdot \operatorname{Den} 10000
     t: Ha \operatorname{Ing} ! Hb' ? ``u: \operatorname{Pot} \cdot \operatorname{Nnc} Hb \operatorname{Cpu} \cdot \operatorname{Etc} \cdot \operatorname{Mod} u: \Leftrightarrow \operatorname{Pau} \operatorname{Pst} :
         Hb \operatorname{Dat} Ha \cdot \operatorname{Den} 10000 I
Ha\operatorname{Ing} Hb\operatorname{Deb} I
Hb \operatorname{Inq} Ha : \wedge \lceil t \cdot t' \cdot p \rceil \cdot tt' Hg \operatorname{Sci} \cdot \operatorname{Utr} \operatorname{Deb} p :
     Erg: Nnc Hb Vul: Nnc Ha \sim Hb. Inq Hg 'Utr: Deb.
         Nnc Hb Dat Ha. Den 10000 |
Ha \operatorname{Ing} Hb \operatorname{Ben} I
Ha \sim Hb \cdot \text{Inq} ! Hg \cdot ! w \cdot PAN Ha Hb \text{ Ise Inq } w \text{ }
Hg \operatorname{Ing} Ha \sim Hb - Deb · Nnc Hb \operatorname{Dat} Ha . Den 10000:
     Deb · Nnc Ha Cpu . ? Etc I
Ha \operatorname{Ing} Hb \operatorname{Pon} x = u + v:
     Erg. u^{11} + v^{11} + (11 uv - 11)(u + v) = 1000,001:
     Pon • uv = 1  \sim u^{11} + v^{11} = 1000,001  \sim
u^{11} = p: Erg. p + 1/p = 1000,001
Hb \operatorname{Inq} Ha : \operatorname{Ben} : \wedge \lceil a \cdot b \rceil \quad a \stackrel{\wedge}{\sim} b \in \operatorname{Rea}
          \rightarrow: Nnc Hb Sci · ? x \cdot x^{11} + ax + b = 0
     Pot · Nnc Hb Cpu. ? Etc I
t1 ¥ t2
Hc \operatorname{Ing} Hc \cdot t_1 t_2 Hb \operatorname{Dat} Ha. Den 10000^{t_3}
 Hd \operatorname{Inq} Hd : t_2 t_3 Hc \operatorname{Anl} \cdot \forall z \cdot t_1 t_2 \operatorname{Fit} z *
```

The word written Dat (fL dat=gives) means "gives". The word written Den (fL denarius=penny) means some unit of value.

‡ is a metatextual substitute for a sequence of 16 striking sounds.

It means the payment of the 16 units. As this is the first action that cannot be derived from the fundamental action Inq, we did not stage it in the form 'Hb Dat Ha. Den 10000', but rather by means of a sequence of characteristic sounds. One could repeat the text with 'Hb Dat Ha. Den 10000' instead of 'f'.

Of course it is not certain that the receiver will guess more details of the event of paying.

```
3 38 2. \# Ha Inq Ha; tt'x Dat y. Den z.
                          \rightarrow t' \operatorname{Pec} y - t \operatorname{Pec} y = t \operatorname{Pec} x - t' \operatorname{Pec} x = \operatorname{Den} z
                     t \in \text{Tem.} \land x \in \text{Hom} \rightarrow \forall u : u \in \text{Rea.} \land t \text{Pec } x = \text{.Den } u *
                      The word written Pec (L pecunia = fortune) means "fortune".
                  't \operatorname{Pec} x' means the fortune of x at the moment t.
3383. The next talk continues those of 3294 and 3297.
                     Hd \operatorname{Inq} He \operatorname{Cur} : t \cdot Ha \operatorname{Vul} q \longrightarrow \operatorname{Pau} \operatorname{Pst} \cdot Hc \operatorname{Vul} q \operatorname{I}
                     He Inq Hd :: Qia .. MAN Ha Pol Hc .
                          \land t 't Pst Nnc. \rightarrow : t \cdot Ha \text{ Vul } q \rightarrow \text{Pau Pst. } Hc \text{ Vul } q:
                               \rightarrow: \land t: \lor n n \in \text{Num}. \land: t - \text{Nnc} = \cdot \text{Sec}. n \times 10^{10100}.
                                    \rightarrow t Ha \text{ Dat } Hc. \text{ Den } 10^{1100}
                     Hd Inq He: Cur - PAN:
                          Ha \operatorname{Ing} He \cdot \operatorname{Nnc} Ha \operatorname{Vul} q \longrightarrow \operatorname{Pau} \operatorname{Pst} \cdot \operatorname{He} \operatorname{Vul} q 
                     He \operatorname{Inq} Hd : \operatorname{Qia} : \forall x \cdot \operatorname{Deb} \cdot Ha \operatorname{Dat} He \cdot \operatorname{Den} x :
                          \operatorname{Tan}: -\operatorname{Nnc} \operatorname{Ha} \operatorname{Vul} \cdot \operatorname{Nnc} \operatorname{Ha} \operatorname{Dat} \operatorname{He} \cdot \operatorname{Den} x
                     Hd Inq He. Cur Deb I
                     He \operatorname{Ing} Hd : \operatorname{Qia} : \forall x : \operatorname{MAN} Ha \operatorname{Pol} He : \operatorname{PPN} Ha \operatorname{Dat} He : \operatorname{Den} x :
```

3 39 1. ${}^{t_1}Ha\ Hb\ \text{Alt Inq } \cdot 1.1.11.1001.110.100100.10.100.}$ $110.100100.1101.10101001^{t_2}$; $Hc\ \text{Inq } Hd\ {}^{t_1} \leq t \leq t_2. \rightarrow : t\cdot Ha\ \text{Inq } Hb\ x.$ $\rightarrow \text{Pau Pst. } Hb\ \text{Inq } Ha\ x^{10}$; $t\ Ha\ \text{Inq } Hb\ x. \rightarrow \text{Deb}: \text{Pau Pst. } t\cdot Hb\ \text{Inq } Ha\ x^{10}$;

Lie' $t Ha \operatorname{Inq} Hb : \forall y \cdot \operatorname{Nne} Ha \operatorname{Vul} \cdot \operatorname{Nne} Ha \operatorname{Inq} Hb y^{t_a}$:

 \land 'MAN Ha Vul: \neg Unq Pst MAN. Ha Dat He. Den $x \neq$

 $Hd \operatorname{Inq} Hc^{\bullet \vee} v \cdot t_2 t_3 Hc \operatorname{Inq} Hd v \cdot \in : \operatorname{Lex} \cdot {}^{\vee} z \cdot t_1 t_2 \operatorname{Fit} z *$

The word written Lex (fL lex=law) means "law". The law of a discourse has been stated.

```
3 39 3. ^{\text{# }t_1}Ha\ Hb\ \text{Alt Inq · 1 . 1 . 10 . 11 . 101 . 1000 . 1101} Etc<sup>t<sub>2</sub></sup>:
                   Hc \operatorname{Inq} Hd : x \cup y := : Ha \cup Hb : \land : t_1 \leq t \leq t_2 : \rightarrow `
                         t: x \operatorname{Inq} y u \cdot y \operatorname{Inq} x v \longrightarrow \operatorname{Pau} \operatorname{Pst} x \operatorname{Inq} y \cdot u + v^{t_0}
                   Hd \operatorname{Inq} Hc \overset{\checkmark}{} v \cdot t_2 t_3 Hc \operatorname{Inq} Hd v \cdot \in : \operatorname{Lex} \overset{\checkmark}{} z \cdot t_1 t_2 \operatorname{Fit} z *
                   <sup>t</sup>, Ha Hb Alt Inq · 110 . 1101 . 10010 . 10100 . 10111 .
3 40 1.
                         11000 - 100000 t<sub>2</sub> L
                    Hb Dat Ha. Den 1000^{t_3}!
                    Hc Inq Hd. Cur Dat I
                    Hd Inq Hc · Qia . PAN Ha Vin I
                    Hc \operatorname{Inq} Hd \cdot ? = \operatorname{Vin} \mathsf{I}
                    Hd \operatorname{Inq} Hc : t_1 t_2 Ha Hb \operatorname{Lud}_1
                        Deb • t<sub>2</sub> t<sub>3</sub> Hb Dat Ha . Den 1000 • Qia : Pau Ant . t<sub>3</sub> • Ha Vin •
                        Pau Ant. t<sub>3</sub> · Ha Vin : Qia : Pau Ant. t<sub>3</sub> · Ha Inq Hb 100000 l
                    Hc \operatorname{Inq} Hd : \operatorname{Cur} x \cdot \operatorname{lum} \operatorname{Par} \cdot v \cdot t_1 t_2 Ha \operatorname{Inq} Hb v:
                         \rightarrow x = 110: Sed \rightarrow x = 100000:
                        \land Cur'x. 1um Par'v. t_1 t_2 Hb Inq Ha v:
                         \rightarrow x = 1101 : \text{Sed} : \neg x = 100000 : \text{Etc}
                    Sed .. Lic x \cdot 1^{um} \operatorname{Par} \cdot \forall v \cdot t_1 t_2 Ha \operatorname{Inq} Hb v:
                             \rightarrow x \in \text{Num.} \land x \leq 1000
                        \land Lie: y. 1^{um} Par. \forall w. t_1 t_2 Hb Inq Ha w:
                             \rightarrow 'y \in \text{Num} . \land : y \leq ``x \cdot x \cdot 1^{\text{um}} \text{ Par } .`v \text{ Etc} : +1000 .
                        \wedge Lie z. 10^{um} Par \dot{v} t_1 t_2 Ha Inq Hb v:
                             \rightarrow: z \in \text{Num} \land z \leq \forall y \cdot y \cdot 1^{\text{um}} \text{Par} \land w \text{Etc} : +1000 \cdot \text{Etc} \cdot 1
                        Hc \operatorname{Inq} Hd : \operatorname{Pot} : tt' x \operatorname{Inq} y : 1^{\bullet} : \uparrow r : \neg \operatorname{Lic} : tt' x \operatorname{Inq} y r I
                   Hd Inq Hc: Qqm Pot. Tan - Lie:
                        Qia:tt'x \operatorname{Inq} y: 1:r \cdot \neg \operatorname{Lie} tt'x \operatorname{Inq} yr
                             \rightarrow Pau Pst t' x Prd \sim y Vin I
                    Hc \operatorname{Inq} Hd \operatorname{!} \operatorname{Utr} \operatorname{!} \operatorname{Lic} \operatorname{!} \neg \lor w
                        Pau Pst. 1^{um} v \cdot t_1 t_2 Ha \operatorname{Inq} Hb v : Hb \operatorname{Inq} Ha w
                   Hd \operatorname{Inq} Hc :: \operatorname{Fal} : \operatorname{Qia} : \operatorname{Deb} : \operatorname{Pau} \operatorname{Pst} : x \operatorname{Inq} y p \cdot y \operatorname{Inq} x
                             1^e: \uparrow q: q \in \text{Num} . \land . p < q \leq p + 1000
                        ^{\mathsf{t_4}}\mathrm{Cca}\,t . Ha\,Hb . \mathrm{Lud}_1 : \rightarrow . Deb : \forall k \colon \lor \ulcorner p_1 \cdot p_2 \cdot \mathrm{Etc} \cdot p_k \urcorner \colon
                             \land i: i \in \text{Num} . \land . 1 \leq i \leq k-1 \rightarrow p_i < p_{i+1} \leq p_i + 1000 . \land .
                                  p_i \in \text{Num} : \land p_1 \leq 1000 \land p_k = 100000
                                  \wedge: Cca t. Ha \sim Hb Alt Inq \cdot p_1 \cdot p_2. Etc. p_k:
                                  \wedge : \lceil x \rceil \cup \lceil y \rceil := Ha \cup Hb \rightarrow \operatorname{Cea} t : \forall x \cdot \operatorname{Cea} t \cdot x \operatorname{Inq} y
                                       100000: Vin ts I
                    Hc\operatorname{Inq} Hd: \Upsilon r \cdot \mathsf{t}_4 \, \mathsf{t}_5 \, Hd\operatorname{Inq} Hc \, r \cdot \in \operatorname{Lex} \cdot Ha \, Hb \operatorname{Lud}_1 :
                        w \in \operatorname{Cca} t_1 \cdot xy \operatorname{Lud}^{\operatorname{io}} : \to \operatorname{LDeb} : \forall k : \forall \lceil p_1 \cdot p_2 \cdot \operatorname{Etc} \cdot p_k \rceil :
```

```
p_1 
hline p_2 
hline Etc 
hline p_k 
hline Num 
hline Pum 
hlin
                                                                                          \land p_1 \leq 1000 \land p_1 < p_2 \leq p_1 + 1000 \land \text{Etc.} \land p_k = 100000
                                                                                          \wedge \cdot \operatorname{lum} \operatorname{Ict} w = x \operatorname{Inq} y p_1 \cdot
                                                                                          \wedge \cdot 10^{\mathrm{um}} \operatorname{Ict} w \cdot = y \operatorname{Inq} x p_2 \cdot
                                                                                          ^ Etc .
                                                                                         \land : Ult^{um} Ict w := : x Inq y p_k : x Vin : y Prd : y Dat x : Den 1000
                                                                                                          \sim = : y \operatorname{Ing} x p_k . y \operatorname{Vin} . x \operatorname{Prd} : x \operatorname{Dat} y . \operatorname{Den} 1000 I
                                                             Hd \operatorname{Inq} Hc \cdot ? = \operatorname{Ict} \mathsf{I}
                                                             Hc\operatorname{Inq} Hd w = xy\operatorname{Lud} : \rightarrow : z = i^{\operatorname{um}}\operatorname{Ict} w \leftrightarrow z \cdot i^{\operatorname{um}}\operatorname{Par} \cdot w + i^{\operatorname{um}} \cdot i^{\operatorname{
                                                               The word written Lud (fL ludit = plays) means "to play a
                                                  game". We shall become acquainted with a variety of games,
                                                  which will be distinguished by indices.
                                                                 The words written Vin (fL vincit = wins) and Prd (fL perdit =
                                                  loses) mean "wins" and "loses" respectively.
                                                                 The word written Ict (fL ictus = blow) means (the substantive)
                                                  "move".
3 40 2. # Ha \operatorname{Inq} Hb \cdot \operatorname{Nnc} Ha \operatorname{Vul} \cdot \vee \lceil a \cdot b \rceil = a \stackrel{\wedge}{\cdot} b \in \operatorname{Pos}.
                                                                                           \land PPN Fit \mathfrak{p} \longrightarrow Pau Pst \cdot Ha Dat Hb . Den a:
                                                                                          \land PPN Fit \mathfrak{g} \rightarrow: Pau Pst \cdot Hb Dat Ha \cdot Den b \mid
                                                             Hb \operatorname{Inq} Ha \operatorname{Ben} I
                                                              Hc \operatorname{Ing} Hc \operatorname{`Nne} Ha \curvearrowright Hb \operatorname{Vul} \operatorname{`Nne} Ha Hb \operatorname{Lud} :
                                                                            PPN • Fit \mathfrak{q} . \rightarrow . Ha \text{ Vin}:
                                                                            PPN · Fit \mathfrak{p} . \rightarrow . Hb \text{ Vin}:
                                                                            t Ha \operatorname{Vin} \cdot \leftrightarrow \cdot t Hb \operatorname{Prd} :
                                                                            t Hb \operatorname{Vin} . \Leftrightarrow . t Ha \operatorname{Prd} *
3 40 3. A continuation of 3 40 1.
                                                              Hc\operatorname{Inq} Hd\operatorname{Lne} Hc\operatorname{Vul}\operatorname{Nne} Hc\operatorname{H} d\operatorname{Lud}_1 "
                                                                            Tan: \dot{x}. PPN x Prd · Dat · \dot{y}. PPN y Vin · Den 10000:
                                                              Ha \operatorname{Inq} Hc \operatorname{Ben}:
                                                              Hc Hd Alt Inq · 1000.1110.1111.10111:
                                                              Hc \operatorname{Inq} Hd. Nnc Hc \operatorname{Prd}:
                                                              Hc \operatorname{Dat} Hd \cdot \operatorname{Den} 10000:
                                                              Hc Inq Hd: Nnc Hc Vul. Nnc Hc Hd Lud<sub>1</sub>
                                                                              PPN: ^{\checkmark}x.x Prd \cdot Dat \cdot ^{\checkmark}y.y Vin \cdot Den 10000:
                                                               Hd \operatorname{Inq} Hc \operatorname{Ben} \Xi
                                                               Hc Hd Alt Inq · 101 . 1101 . 1110 . 1111 . 10111 :
                                                               Hd \operatorname{Inq} Hc \cdot \operatorname{Nnc} Hd \operatorname{Prd}:
                                                               Hd \operatorname{Dat} Hc \cdot \operatorname{Den} 10000:
                                                               Hd Inq Hc · PAN Hc Lud<sub>1</sub> Mod Mal . Nnc Hc Lud<sub>1</sub> Mod Ben :
```

 $Hc \operatorname{Inq} Hd : \operatorname{Nnc} Hc \operatorname{Sci} : ? p \cdot Hc \operatorname{Inq} Hd p \cdot \in \operatorname{Ben}^*$

Nes: 1^{um} Ict. = $101 \cdot \land \cdot 11^{um}$ Ict. = $1110 \cdot \land \cdot 101^{um}$ Ict. = 10111

```
Nne Hc Vul. Nne Hc Vin \rightarrow: Nes · 1 um I et . = 101 : Etc ·
                          \operatorname{Cea} t \cdot xy \operatorname{Lud}_1 \cdot \wedge \cdot x \operatorname{Inq} y p \in \operatorname{Ben} : \to \cdot \operatorname{Cea} t \cdot x \operatorname{Vin}
                          \operatorname{Erg}: -_{1}\operatorname{Pot}\cdot\operatorname{Cea} t.y\operatorname{Vin}^{\bullet}\operatorname{Nes}\cdot\operatorname{Cea} t.y\operatorname{Prd}^{\bullet}
                          t: x \text{ Vul } \cdot \text{Nnc } x \text{ Vin } \rightarrow \text{Pau Pst } \cdot x \text{ Vin } \cdot
                          Lex xy Lud<sub>1</sub> \notin Ius \#
                       The word written Ius (fL iustum=righteous) means "fair".
3 40 4. \# Ha \operatorname{Inq} Hb \cdot 10 \cdot \operatorname{Mod} - \operatorname{Pot} Hb \operatorname{Ani} I
                     Hb \operatorname{Ing} Ha. 11. \operatorname{Mod} - \operatorname{Pot} Ha \operatorname{Ani} I
                     Hc \operatorname{Inq} Ha - Hb \cdot \operatorname{PAN} Ha \operatorname{Inq} Hc 10 - Hb \operatorname{Inq} Hc 11.
                           \operatorname{Erg} Ha\operatorname{Vin} I
                     Ha \operatorname{Inq} Hc.11 \cdot \operatorname{Mod} \operatorname{Etc} I
                     Hb \operatorname{Inq} Hc.1.\operatorname{Mod} \operatorname{Etc} I
                     Hc \operatorname{Ing} Ha \sim Hb: 11.1 \cdot \operatorname{Erg} \cdot Ha \operatorname{Vin} I
                     Ha Inq Hc. 100. Mod Etc!
                     Hb \operatorname{Ing} Hc. 100 \cdot \operatorname{Mod} \operatorname{Etc} I
                     Hc \operatorname{Inq} Ha - Hb : 100 \cdot 100 \cdot \operatorname{Erg} \cdot Hb \operatorname{Vin} I
                     Hd \operatorname{Inq} He \ldots \operatorname{PAN} Ha Hb \operatorname{Lud}_2 \mathbf{I}
                           \operatorname{Cca} t \cdot xy \operatorname{Lud}_2 \longrightarrow \operatorname{Deb} : \vee \lceil z \cdot p \cdot q \rceil
                                \operatorname{Cea} t \cdot x \operatorname{Inq} z \cdot p \cdot \operatorname{Mod} - y \operatorname{Ani} \cdot \wedge \cdot \operatorname{Cea} t \cdot y \operatorname{Inq} z \cdot q
                                      Mod - xAni
                                \wedge \cdot \lceil p \rceil \cup \lceil q \rceil \cdot \subset \cdot 1 \cup 10 \cup 11 \cup 100
                                . \land p \neq q . \rightarrow : \text{Cea } t \cdot z \text{ Inq } x \land y . \text{Nne } x \text{ Vin}
                                \land p = q . \rightarrow : \operatorname{Cea} t \cdot z \operatorname{Inq} x \land y . \operatorname{Nnc} y \operatorname{Vin} I
                           \forall v . PAN Hd Inq He v \cdot \in : Lex \cdot xy Lud_2 I
                      Hc\operatorname{Inq} Hd \cdot \operatorname{Lex} \operatorname{Etc} \cdot \notin \operatorname{Ius} \mathsf{I}
                      Hd\operatorname{Inq} Hc , \operatorname{Lex}\operatorname{Etc} \in\operatorname{Ius} :
                           Qia Deb Nnc: x \text{ Vin} \rightarrow \text{Pau Pst} \cdot y \text{ Dat } x \text{ Den } 1
                                 \land 'Nnc: y Vin. \rightarrow Pau Pst \cdot x Dat y. Den 100:
                           \operatorname{Cca} t \cdot xy \operatorname{Lud}_2 \cdot
                                \rightarrow: Prb · Cca t · x Vin := 0,11 · \land · Prb · Cca t · y Vin := 0,01 |
                      Hd\operatorname{Ing} Hc\operatorname{:}\operatorname{Cea} t\operatorname{.} Hb\operatorname{Vin}
                                \leftrightarrow: Cea t. Hb Inq! Hc \cdot ?u. PAN Ha Inq Hcu
                           \operatorname{Tan} - Hb\operatorname{Sci} \operatorname{Etc} - \operatorname{Qia} - \operatorname{Ani} - \operatorname{Lie} \cdot Hb\operatorname{Ani}
                           - Lie: Pau Ant. Hb \operatorname{Inq} Hcq \cdot Hc \operatorname{Inq} !Hb \cdot ?p. PAN Ha \operatorname{Inq} Hcp
                           Pau Ant. Hb Inq Hcq · Hb Sei Etc:
                                \rightarrow: Cca t. Ha Hb Lud<sub>2</sub>: \notin Ius
                           Erg · Deb . — Sci #
                       The word written Prb (fL probabilitas = probability) means
```

"probability". Though mechanical stochastic devices are not

available, this simple game (akin to "matching pennies") gives us the opportunity to introduce numerical probability. Of course it should be repeated with other numerical constants.

From the next text the syntax of Prb will become clearer.

```
3 40 5. # Ha Hb Lud<sub>4</sub> |
                   Hc \operatorname{Inq} Hd \cdot \operatorname{Nne} Hc \operatorname{Vul} \cdot \operatorname{Nne} Hc Hd \operatorname{Lud}_i
                   Hd \operatorname{Inq} Hc : ? = \operatorname{Prb} : \operatorname{Fin} \cdot xy \operatorname{Lud}_i \cdot x \operatorname{Vin} I
                   Hc \operatorname{Ing} Hd \cdot \operatorname{Prb} : \operatorname{Fin} \cdot xy \operatorname{Lud}_i \cdot x \operatorname{Vin} = 101/111
                         Erg 'Prb: Fin xy Lud<sub>i</sub> y 'Vin = 10/111!
                         Qia 'Prb : Fin xy Lud x Vin '+ 'Prb : Vin xy Lud y Vin '= 1
                   Hd \operatorname{Inq} Hc : ? \cap p \cdot q \cap \operatorname{Deb} t : x \operatorname{Vin} \rightarrow y \operatorname{Dat} x \cdot \operatorname{Den} p
                         \wedge t: y Vin \rightarrow x Dat y . Den q 1
                   Hc \operatorname{Inq} Hd \cdot p = 11 \cdot \wedge \cdot q = 100 
                   Hd \operatorname{Inq} Hd : x \operatorname{Uti} : xy \operatorname{Lud}_{i} = \operatorname{Den} p \times \operatorname{Prb} : \operatorname{Fin} : xy \operatorname{Lud}_{i} : x \operatorname{Vin} :
                         -: Den q \times 'Prb: Fin. xy Lud<sub>i</sub> · y Vin := Den 1 |
                   Hd Inq Hc : Nne Hd Vul
                         t: Hc Hd \operatorname{Lud}_{\mathfrak{l}} . \Leftrightarrow \operatorname{Pau} \operatorname{Ant} \cdot Hc \operatorname{Dat} Hd \cdot \operatorname{Den} 1:
                         Qia \cdot Hc Uti = Den 1 \cdot \wedge \cdot Hd Uti = Den - 1 
                     The word written Uti (/L utilitas = utility) means "utility",
                "expectation".
3 40 6. # Hd Inq Hd
                         Prb: Fin xy Lud x Vin = a:
                         \land Prb: Fin xy Lud y Vin = b = 1 - a.
                         \wedge: w \in \cdot \text{Lex.} xy \text{Lud}:
                         \wedge v = t \cdot x \operatorname{Vin} \rightarrow y \operatorname{Dat} x \cdot \operatorname{Den} p
                              \wedge t: y \text{ Vin} \rightarrow x \text{ Dat } y \cdot \text{Den } q \cdot \wedge v \text{ Par } w
                         \rightarrow "x Uti, xy Lud = Den ap-bq:
                             \wedge : y \text{ Uti } . xy \text{ Lud } = \text{ Den } . bq - ap :
                              \wedge : w \in \text{Ius.} \Leftrightarrow x \text{Uti.} = y \text{Uti.} = \text{Den } 0 \text{ } 
3 40 7. \# Ha Inq Hc: 11: Mod \cdot – Pot. Hb Ani I
                    Hb \operatorname{Inq} Hc: 1: \operatorname{Mod} - \operatorname{Pot} Ha \operatorname{Ani} I
                    Hc\operatorname{Inq} Ha \curvearrowright Hb: 11.1 · Erg · Ha\operatorname{Vin} I
                    Ha Inq Hc. 1. Mod Etc |
                    Hb \operatorname{Inq} Hc. 10 \cdot \operatorname{Mod} \operatorname{Etc} \mathsf{I}
                    Hc \operatorname{Inq} Ha - Hb : 1 \cdot 10 \cdot \operatorname{Erg} \cdot Ha \operatorname{Vin} \mathbf{I}
                    Ha \operatorname{Inq} Hc. 10. \operatorname{Mod} \operatorname{Etc} I
                    Hc \operatorname{Inq} Ha. 10. \operatorname{Mod} \operatorname{Etc} I
                    Hc\operatorname{Inq} Ha \sim Hb:10.10 \cdot \operatorname{Erg} \cdot Hb\operatorname{Vin} I
                    Ha \operatorname{Dat} Hb \cdot \operatorname{Den} 10
```

Hd Inq He. PAN Ha Hb Luds I

```
He \operatorname{Ing} Hd ? \forall w : w \in \operatorname{Lex} xy \operatorname{Lud}_3
                  Hd \operatorname{Ing} He : \operatorname{Deb} : \vee \lceil a \cdot b \cdot z \cdot t \rceil a \rightsquigarrow b \cdot \in 1 \cup 10 \cup 11 \cup 100 \land A
                       "\operatorname{Cca} t \cdot Ha \operatorname{Ing} z : a : \operatorname{Mod} \cdot - \operatorname{Pot} \cdot Hb \operatorname{Ani} \wedge
                      "\operatorname{Cea} t \cdot Hb \operatorname{Ing} z : b : \operatorname{Mod} \cdot - \operatorname{Pot} \cdot Ha \operatorname{Ani} \wedge
                      : \operatorname{Cca} t \cdot Hb \operatorname{Vin} \cdot \leftrightarrow a = b \cdot \land
                       " \operatorname{Cca} t : Hb\operatorname{Vin} . 	o " Ha\operatorname{Dat} Hb . \operatorname{Den} a \mathsf{I}
                  He Ing Hd · Nnc He Vul . Nnc Hd He Lud<sub>3</sub> I
                  Hd Inq He: Nnc Hd Vul 'Nnc: Hd He Lud3.
                       \leftrightarrow Pau Ant · He Dat Hd . Den 0,101 |
                  He \operatorname{Inq} Hd \cdot \operatorname{Cur} 0,101 
                  Hd \operatorname{Ing} He : \operatorname{Qia} : \operatorname{Prb} : Hd \operatorname{Ing} Hc a := \cdot \operatorname{Prb} : He \operatorname{Ing} Hc b := 0.01
                      Erg: Prb \cdot Hd Inq Hc i \land He Inq Hc i = 0.0001
                      \operatorname{Erg} \cdot \operatorname{He} \operatorname{Uti} = \operatorname{Den} 0,101 \operatorname{I}
                  He \operatorname{Ing} Hd \cdot \operatorname{Utr} \cdot \operatorname{Deb} : \operatorname{Prb} \cdot Hd \operatorname{Ing} Hc i := 0.01 \text{ }
                  Hd \operatorname{Ing} He : \neg \operatorname{Deb} : \operatorname{Qia} : \neg \operatorname{Prb} = 0.01 \cdot \operatorname{Par} : \forall w
                      w \in \text{Lex Lud}_3: Tan: Prb = 0.01. \rightarrow
                      He Uti = Den 0,101 I
                  He \operatorname{Ing} Hd: ? = \cdot \operatorname{Prb} \cdot Hd \operatorname{Ing} Hei!
                  Hd \operatorname{Ing} He : - \operatorname{Deb} \cdot \operatorname{Nnc} Hd \operatorname{Ing} ! He \cdot ? = \operatorname{Etc} !
                  He \operatorname{Ing} He : \operatorname{Pon} \operatorname{Prb} \cdot Hd \operatorname{Ing} Hc i = c_i:
                           \wedge: Prb. He Inq Hcj := d_i:
                      Erg • He Uti = c_1 d_1 + 10 c_2 d_2 + 11 c_3 d_3 + 100 c_4 d_4:
                      Exg: Nnc "Hd Vul. 11 c_3 \ge 1 c_1 \land \ge 10 c_2 \land \ge 100 c_4 \rightarrow Pau Pst:
                           He Vul \cdot d_3 = 1 \cdot He Uti = Den 11 c_3 \cdot Etc
                      Erg: Nnc Hd Vul. He Uti \leq Den e: \rightarrow
                           "Nes. Neg Hd Vul "\land i: i=1 \sim 10 \sim 11 \sim 100. \rightarrow ic_i \leq e.
                           \operatorname{Erg} \cdot c_i \leq e/i \cdot \operatorname{Erg} \cdot 1 = c_1 + c_2 + c_3 + c_4 \leq 11001 \, e/1100 \, .
                           Erg. e \ge 1100/11001:
                       \neg Pot: Nnc Hd Vul·He Uti. < 1100/11001:
                       Ncs: Nnc Hd Ccd · He Uti \geq 1100/11001
                       Tan 'Pot: Nnc Hd Vul·He Uti = 1100/11001 |
                   He \operatorname{Inq} Hd Nnc He \operatorname{Ccd}:
                       Ant. HdHeLud_3 \cdot HeDatHd. Den 1100/11001 I
                   Hd\operatorname{Inq} He \cdot \operatorname{Ben} \cdot \operatorname{Qia}\operatorname{Ius} *
               von Neumann's minimax principle.
3 40 8. # HaHb Alt Inq Hc: 1.10.1.1.10.10.10:
                       \operatorname{Mod}: \neg \operatorname{Pot} Hb \operatorname{Ani} ? \Upsilon u \cdot Ha \operatorname{Inq} Hc u :
                            \sim: Ha Ani ·? ^{\checkmark}v. Hb Inq Hcv
                   Hc \operatorname{Inq} Ha \sim Hb \cdot 1 + 10 + 1 + 1 + 10 + 10 + 10 \ge 1010. Ha \operatorname{Vin} I
                   Hd Ing He. PAN Ha Hb Lud41
```

 $He \operatorname{Inq} Hd$ '? $\forall w : w \in \operatorname{Lex} Ha Hb \operatorname{Lud}_4 I$

```
Hd\operatorname{Inq} He: Deb \exists \lor k \exists \lor \ulcorner p_1 . p_2 . \operatorname{Etc} . p_k \urcorner
                     p_1 
ightharpoonup p_2 
ightharpoonup \operatorname{Etc} 
ightharpoonup p_k . \in 1 \cup 10 : \land : Ha Hb Alt Inq
                     p_1 \cdot p_2 \cdot \text{Etc.} p_k : \land : k^{\text{um}} \text{Ict.} = \cdot \text{Ultum Ict.} \leftrightarrow \cdot
                     p_1 + p_2 + \text{Etc} + p_k \ge 1010:
                     \land x \in Ha \cup Hb \rightarrow Ult^{um} Ict x Vin \leftrightarrow
                         Ultum Ict. = x \operatorname{Ing} Hc p_k
                     \land 't: Ha \text{ Vin} \rightarrow \text{Pau Pst} \cdot Hb \text{ Dat } Ha \cdot \text{Den 1}
                     \wedge t: Hb \text{ Vin} . \rightarrow \text{Pau Pst} \cdot Ha \text{ Dat } Hb . \text{ Den 1 I}
                 He\operatorname{Inq} Hd 	riangleq Hc \cdot \operatorname{Nnc} He\operatorname{Vul} . \operatorname{Nnc} Hd\operatorname{He}\operatorname{Lud}_4 \operatorname{I}
                 \mathit{Hd} 	riangleq \mathit{Hc} \operatorname{Inq} \mathit{He} \operatorname{Ben} \mathsf{I}
                 HdHe. Alt InqHc \cdot 1. 10.1.10.10.1.10 · Mod Etc I
                 Hc \operatorname{Inq} Hd - He \cdot 1 + 10 + 1 + 10 + 10 + 1 + 10 \ge 1010 \cdot Hd \operatorname{Vin} I
                 He\operatorname{Inq} Hd 	riangleq Hc \cdot \operatorname{Nnc} He\operatorname{Vul} . Etc I
                 Hd 	riangleq Hc \operatorname{Inq} He \operatorname{Ben} I
                 Hd\,He . Alt Inq Hc \cdot 1 . 1 . 1 . 1 . 10 . 10 . 10 \cdot Mod Etc \mid
                 Hc \operatorname{Inq} Hd - He \cdot 1 + 1 + 1 + 1 + 1 + 10 + 10 + 10 \ge 1010. Hd \operatorname{Vin} I
                 He \operatorname{Inq} He \operatorname{LUtr} - \operatorname{Pot} : \forall t \cdot \operatorname{Cea} t \cdot He \operatorname{Vin} \mathbf{I}
                 HdHe Lud<sub>4</sub>: Mod: 1.1.1.1.10.1.10.1 He Vin
                     Mod: 10.1.10.1.10.1.10 · Hd Vin
                     Mod: 10.1.10.1.10.10.10 He Vin
                     Mod:1.1.1.1.10.10.10.10 Hd Vin I
                 He Inq He: PAN He. Lud4. Mod Mal.
                     A. PPN He. Lud4. Mod Ben:
                     Nnc He Sci : ? \forall w Cca t: He Vul . Nnc He Vin •
                         \rightarrow Ncs. Nnc He Inq Hc \forall w I
                 Hd He Lud4 Mod: 10.1.10.1.10.10 He Vin
                     Mod: 10.10.10.10.10.10.Hd Vin'
                     Mod:1.10.1.10.10.10.10·HeVin
                     Mod: 10.1.10.1.10.10.10 · He Vin
                     Mod:1.1.1.1.10.10.10 · Hd Vin I
                 Ht \operatorname{Ing} Ht : \operatorname{Nnc} Hd He \operatorname{Lud}_{4} : \operatorname{Mod}_{5} :
                     Prb: x Inq Hc \cdot 1.1.10.10.10.Etc
                      = \Pr : x \operatorname{Inq} Hc \cdot 10.10.10. Etc:
                     x \operatorname{Lud}_4 \cdot \operatorname{Mod} \cdot \operatorname{Nnc} Hd He \operatorname{Lud}_4 \longrightarrow x \operatorname{Lud}_4 \in \operatorname{Ben}^{\sharp}
                     Lex Lud₄. ∈ Ius !
                 He \operatorname{Dat} Hd \cdot \operatorname{Den} 110 - \operatorname{Den} 101 = . \operatorname{Den} 1 
3 40 9. # Ha Hb Hc. Alt Inq · 1.11.101: Hc Vin:
                 Ha\,Hb\,Hc. Alt Inq·1.10.11.101: Ha\,\mathrm{Vin}:
                 Ha\,Hb\,Hc . Alt Inq · 10 . 11 . 101 : Hc\,\mathrm{Vin} :
```

 $Ha\,Hb\,Hc$. Alt Inq \cdot 1 . 10 . 11 . 101 : $Ha\,\mathrm{Vin}$

 $\mathit{Hd}\operatorname{Inq}\mathit{He}$. Nnc $\mathit{Ha}\mathit{Hb}\mathit{Hc}\operatorname{Lud}_{5}$.

```
Deb: \forall \lceil t \cdot n \rceil: \forall \lceil p_1 \cdot p_2 \cdot \text{Etc.} p_n \rceil:
                                   p_1 
hline p_2 
hline Etc 
hline p_n 
hline 10 \cdot 11 \cdot 100 \cdot 101 \cdot 10
                                   \land p_1 \leq 10 \land p_2 - p_1 = 1 \checkmark 10 \land p_3 - p_2 = 1 \checkmark 10 \land Etc.
                                   \wedge: Cea t. Ha Hb Hc. Alt Inq \cdot p_1 \cdot p_2. Etc. p_n:
                                    \wedge : n^{\text{um}} \text{ Ict.} = . \text{ Ultum Ict.} \Leftrightarrow p_n \ge 101:
                                    \wedge \ \ \lceil x \rceil \cup \ \lceil y \rceil \cup \ \lceil z \rceil = . Ha \cup Hb \cup Hc \rightarrow :
                                                 Pau Pst. Ultum Ict. x Vin
                                                                \leftrightarrow Ultum Ict = x \operatorname{Inq} yz p_n \wedge
                                   Cca t \cdot Ha \operatorname{Vin} \longrightarrow \operatorname{Pau} \operatorname{Pst} : Hb \operatorname{Dat} Ha \cdot \operatorname{Den} 1000 \cdot
                                                  Hc \operatorname{Dat} Ha \cdot \operatorname{Den} 1000 ^ ^
                                   Cca t \cdot Hb \text{ Vin} \rightarrow \text{Pau Pst} \cdot \text{Etc} \land
                                   Cca t \cdot Hc \operatorname{Vin} \cdot \to \operatorname{Pau} \operatorname{Pst} \cdot \operatorname{Etc} I
        He \operatorname{Ing} Hd : \notin \operatorname{Ius} : \operatorname{Qia} : \neg \operatorname{Pot} : \forall t \cdot \operatorname{Cea} t \cdot Hb \operatorname{Vin} : \neg \operatorname{Cea} t \cdot Hb \operatorname{V
                     t \cdot Ha \operatorname{Ing} Hb \wedge Hc 10 \rightarrow \operatorname{Pau} \operatorname{Pst} \cdot Hc \operatorname{Vin}
                     t \cdot Ha \operatorname{Ing} Hb \wedge Hc 1 \rightarrow \operatorname{Pau} \operatorname{Pst} \cdot Ha \sim Hc \operatorname{Vin} I
        Hd Inq He:
                     t Ha \operatorname{Ing} 1 \cdot \wedge \cdot z = Ha \sim Hc \cdot \wedge : \operatorname{Pau} \operatorname{Pst} \cdot t \cdot Hb \operatorname{Vul} \cdot \operatorname{Nnc} z \operatorname{Vin}
                                   \rightarrow: Pau Pst \cdot t \cdot z Vin:
                     \operatorname{Cea} t \cdot y \operatorname{Vul} : \operatorname{Pau} \operatorname{Pst} \cdot t \cdot z \operatorname{Vin} \hookrightarrow \operatorname{Cea} t \cdot y \operatorname{Aux} z
                     Pot. Hb Aux Ha: Pot. Hb Aux Hc I
        Ha Hb Hc · Vul · Nne Ha Hb Hc . Lud, I
        Ha \operatorname{Ing} Hb \operatorname{Inc} Ha \operatorname{Pol} Hb
                     Nne Hb Vul. PPN Ha Vin \rightarrow: Nne Ha Ccd.
                                     \operatorname{PPN} Hb\operatorname{Dat} Ha . Den 0
        Hc \operatorname{Inq} Hb : \operatorname{Nnc} Hc \operatorname{Pol} Hb:
                                    Nne Hb Aux Hc \rightarrow Nne Hc Ced:
                                     \neg \neg PPN Hb Dat Hc. Den 1000: Sed · Hc Dat Hb. Den 10:
                      Hc Uti Hb Aux Hc = Den 1100:
         \forall a : \text{Nnc } Hb \text{ Aux } Hc \rightarrow PPN \text{ Pec } Hc = a
                                       - b: Nnc Hb Aux Ha. \rightarrow PPN Pec Hc. = b: = Den 1100:
                       \operatorname{Erg}: Hc\operatorname{Pol} Hb \cdot 0.1 \times \operatorname{Den} 1100 \cdot \in \operatorname{Ius} I
         Ha \operatorname{Inq} Hb : \operatorname{Nnc} Ha \operatorname{Pol} Hb :
                       PPN Ha Dat Hb Den: 10 + x \cdot PAN Hc Pol Hb Den x \mid
        Hb Inq Ha Ben I
        Ha \operatorname{Inq} Hb \wedge Hc 11
         Hb \operatorname{Inq} Hc - Ha 10 I
         Hc \operatorname{Inq} Ha - Hb \cdot \operatorname{Nnc} Hc \operatorname{Prd} I
         Hc Dat Ha. Den 1000 |
         Ha \operatorname{Dat} Hb \cdot \operatorname{Den} 100 I
         Hd Ing He · PAN Hb Aux Ha . Ha Hb Soc *
           The word written Aux (fL auxiliatur = aids) means (the verb)
"aids". Soc (fL socii = allies) means "allies".
```

3 41 1. * $^{t_1}Ha \operatorname{Inq} Ha \cdot \neg \lor x \cdot t_1 t_2 Ha \operatorname{Inq} Ha x^{t_2}$! $Hb \operatorname{Inq} Hc \cdot ? p \cdot t_1 t_2 Ha \operatorname{Inq} Ha p$! $Hc \operatorname{Inq} Hb \cdot \neg \lor x \cdot t_1 t_2 Ha \operatorname{Inq} Ha x$! $Hb \operatorname{Inq} Hc \operatorname{Mal}$! $Hc \operatorname{Inq} Hb \cdot t_1 t_2 Ha \operatorname{Inq} Ha : \neg \lor x \cdot t_1 t_2 Ha \operatorname{Inq} Ha x$! $Hb \operatorname{Inq} Hc \operatorname{Ben}$! $Hc \operatorname{Inq} Hb \cdot \operatorname{Tan} : \lor p \cdot t_1 t_2 Ha \operatorname{Inq} Ha \cdot \in \operatorname{Fal}$! $Hb \operatorname{Inq} Hc \operatorname{Ben} = \mathbb{R}$

A text that prepares the way for the statement of the liar:

3 41 2. ** $^{t_1}Ha \operatorname{Inq} Ha: ^{\checkmark}x \cdot t_1 t_2 Ha \operatorname{Inq} Hax \cdot \in \operatorname{Fal}^{t_1} |$ $Hb \operatorname{Inq} Ha: t_1 t_2 Ha \operatorname{Vul}: \operatorname{Nnc} Ha \operatorname{Inq} Ha \cdot 1^{\epsilon} \operatorname{Ver}^{\epsilon}$ $\to ^{\bullet}t_1 t_2 Ha \operatorname{Vul}: \operatorname{Nnc} Ha \operatorname{Inq} Ha \cdot 1^{\epsilon} \operatorname{Fal}^{\epsilon}$ $\wedge ^{\bullet}t_1 t_2 Ha \operatorname{Vul}: \operatorname{Nnc} Ha \operatorname{Inq} Ha \cdot 1^{\epsilon} \operatorname{Fal}^{\epsilon}$ $\to ^{\bullet}t_1 t_2 Ha \operatorname{Vul}: \operatorname{Nnc} Ha \operatorname{Inq} Ha \cdot 1^{\epsilon} \operatorname{Ver}^{\epsilon}$ $\operatorname{Utr}^{\bullet}t_1 t_2 Ha \operatorname{Vul}: ^{\checkmark}p \cdot \operatorname{Nnc} Ha \operatorname{Inq} Ha \cdot \in \operatorname{Ver}: ^{\checkmark}: \in \operatorname{Fal}^{\bullet}$ $Ha \operatorname{Inq} Hb: t_1 t_2 Ha \operatorname{Vul}: t_1 t_2 Ha \operatorname{Inq} Hb p \cdot \cdot \cdot p \in \operatorname{Fal} \cdot \wedge \cdot p \in \operatorname{Fal} \cdot \rightarrow \cdot p \in \operatorname{Ver}^{\bullet}$ $Hb \operatorname{Inq} Ha \cdot t_1 t_2 \operatorname{Fit} q \cdot \rightarrow \cdot q \in \operatorname{Mal}^{\#}$ The talk of the liar.

CHAPTER IV

SPACE, MOTION, MASS

4000. So far the members of the class 'Hom' might be ghosts. The only extension we needed, was time. We shall now introduce space, motion, mass, and other notions of mechanics. We could do so by axioms, but such a procedure would be unsatisfactory. We prefer the behaviouristic approach. Afterwards the crude ideas we have acquired will be refined and settled with more precision by means of an axiomatic system.

```
4 01 1.
                    Ha \operatorname{Ing} Hb : \forall h : h \in \operatorname{Pos}.
                         \land: Sec h. Pst·Hc Inq Hd \mathfrak{p}: Hd Ani: Utr. PAN Hc Inq Hd \mathfrak{p}
                         \land Sec h. Pst · Hd Inq Hc q: Hc Ani: Utr. PAN Hd Inq Hc q:
                    Hb \operatorname{Ing} Ha \cdot \operatorname{Cur} \cdot h > 0
                    Ha \operatorname{Inq} Hb \cdot \operatorname{Qia} \cdot \operatorname{Loe} Hc \neq \cdot \operatorname{Loe} Hd
                         Loc Hc = \text{Cea} \cdot \text{Loc } Hd \rightarrow \text{Pot} : t_1 t_2 Hc \text{ Ani} \cdot i p \cdot t_1 t_2 Hd \text{ Inq } Hc.
                         Loc Hc \neq \text{Loc } Hd \rightarrow \text{Pot Etc.}
                         Loc Hc \subset \operatorname{Spa} : \operatorname{Loc} Hd \subset \operatorname{Spa} :
                         a \in \operatorname{Spa} : \rightarrow a \in a:
                         \land X: X \in \operatorname{Hom} \cup \operatorname{Bes} \to \operatorname{Loc} X \subset \operatorname{Spa}
                         \forall Z \, \dot{} Z \subset \operatorname{Spa} \, \land \dot{} : \neg \, \forall X \, \dot{} Z = \operatorname{Loc} X \, \dot{} 
                         Sec h. Pst \cdot X Inq Y p : Y Ani Etc \wedge
                         Sec h'. Pst X' Inq Y' p': Y' Ani Etc \land Loc X'. \neq . Loc Y':
                         \rightarrow: Dst XY./. Dst X'Y':= Cea.h/h':
                         \forall p: p \in \text{Pos.} \land \cdot \text{Dst } XY = \cdot \text{Cmt } p
                         \operatorname{Dst} XY = \operatorname{Dst} \operatorname{Loc} X \operatorname{Loc} Y
                         a \in \text{Loc } X \cdot \land b \in \text{Loc } Y \cdot \rightarrow \cdot \text{Dst } XY \cdot = \text{Cea. Dst } ab
```

The word written Loc (fL locus = place) means "place" (of). Delay of signals is symptomatic for difference of place.

The word written Spa (fL spatium = space) means "space" (or, more precisely: set of spatial points). All places are in space.

The word written Dst (fL distantia = distance) means "distance". 'Dst' is said to be proportional to the time of delay. Of course this definition is provisional. It can be applied as long as there is not more than one method of signalizing. In due course we shall learn that there are several kinds.

The word written 'Cmt' means "centimeter". The length unit is still unknown. 'Cmt' behaves syntactically like 'Sec'.

401 2. *
$$a
ightharpoonup b \in \text{Rea}$$
.
 $ightharpoonup \text{Cmt.} a + b \cdot = \cdot \text{Cmt} a + \text{Cmt} b$:
 $hi \cdot \text{Cmt.} a - b \cdot = \cdot \text{Cmt} a - \text{Cmt} b$:
 $hi \cdot \text{Cmt} a \cdot \leq \cdot \text{Cmt} b \cdot \leftrightarrow \cdot a \leq b$:
 $hi \cdot \text{Cmt} a \cdot = \text{Cca.} \cdot \text{Cmt} b \cdot \leftrightarrow \cdot a \cdot = \text{Cca.} b$:
 $hi \cdot a \times \cdot \text{Cmt} b \cdot = \cdot \text{Cmt.} a \times b$: hi .
Etc #

4 01 3. *
$$a imes b imes c . \in \operatorname{Spa}$$
.
 $\to : \forall p : p \in \operatorname{Rea} . \land . p \ge 0 . \land \cdot \operatorname{Dst} ab . = . \operatorname{Cmt} p$.
 $\land \cdot \operatorname{Dst} ab . = . \operatorname{Cmt} 0 : \longleftrightarrow . a = b$:
 $\land : \operatorname{Dst} ab . + . \operatorname{Dst} bc : \ge . \operatorname{Dst} ac$ *

Axioms of metric space.

Special properties of the metrics of Euclidian space.

4 01 5. *
$$a \curvearrowright b \in \operatorname{Spa} \rightarrow$$
:
 $\operatorname{Cvx} . a \cup b := \ \ \ c : \operatorname{Dst} ac . + . \operatorname{Dst} bc := . \operatorname{Dst} ab \ |$

$$A \subset \operatorname{Spa} . \rightarrow : \operatorname{Cvx} A .$$

$$= : \bigcap \ \ B : A \subset B \subset \operatorname{Spa} . \land : \land \lceil x . y \rceil \ \ x \curvearrowright y \in B .$$

$$\rightarrow : \operatorname{Cvx} . x \cup y \cdot \subset B \ ^{*}$$

The word written Cvx means "convex envelope".

401 6.
$$*r = . \operatorname{Rct} ab \cdot \leftrightarrow : a \land b \in \operatorname{Spa} . \land . a \neq b . \land : r = : \uparrow c$$

$$c \in \cdot \operatorname{Cvx} . a \cup b : \lor : a \in \cdot \operatorname{Cvx} . b \cup c : \lor : b \in \cdot \operatorname{Cvx} . a \cup c :$$

$$r \in \operatorname{Rct} \cdot \leftrightarrow : \lor \vdash a . b \vdash \cdot r = . \operatorname{Rct} ab$$

The word written Rct (fL recta = straight line) means "straight lines". 'Rct ab' means the straight line through a and b.

4 01 7. *
$$p \in \text{Pla} . \leftrightarrow : p \subset \text{Spa} . \land : \lor \lceil r.s \rceil \cdot r \neq s$$
.
 $\land . r \curvearrowright s \in \text{Rct} . \land \cdot r \cap s . \neq \lceil \rceil \cdot \land : p = \cdot \text{Cvx} . r \cup s \text{ *}$
The word written Pla (fL planum = plane) means "planes".

- 401 8. * a
 ightharpoonup b
 ightharpoonup Car
 ightharpoonup car relations a parabola a production a produ
- 4 02 1. * $a \in \operatorname{Spa} . \land . \varrho \in \operatorname{Pos} \cdot \rightarrow : \operatorname{Sph} a \varrho . = ` \uparrow x : x \in \operatorname{Spa} . \land \cdot \operatorname{Dst} ax . = \varrho :$ $\land : \operatorname{Bul} a \varrho . = ` \uparrow x : x \in \operatorname{Spa} . \land \cdot \operatorname{Dst} ax . \leq \varrho :$ $s \in \operatorname{Sph} . \leftrightarrow : \lor \ulcorner a . \varrho \urcorner \cdot s = . \operatorname{Sph} a \varrho :$ $s \in \operatorname{Bul} . \leftrightarrow : \lor \ulcorner a . \varrho \urcorner \cdot s = . \operatorname{Bul} a \varrho *$ $The words \operatorname{Sph} (fL \operatorname{sphaera} = \operatorname{sphere}) \text{ and } \operatorname{Bul} (fL \operatorname{bulla} = \operatorname{bubble})$

mean "spheres" and "solid spheres" respectively.

- 4 02 3. * $r \curvearrowright s \in \text{Ret.} \land . a \curvearrowright b \curvearrowright c \in \text{Spa.} \land \cdot c = .r \cap s \cdot \land . a \in r. \land . b \in s. \land .r \perp s.$ $\land \cdot \text{Dst } bc. = \text{Cmt } \alpha \cdot \land \cdot \text{Dst } ca. = \text{Cmt } \beta \cdot \land \cdot \text{Dst } ab. = \text{Cmt } \gamma :$ $\Rightarrow .\alpha^{10} + \beta^{10} = \gamma^{10} *$
- 4 02 4. * $\varphi \in \operatorname{Spa} \curvearrowright \operatorname{Spa} \cdot \to : \varphi \in \operatorname{Rfc}$. $\longleftrightarrow : \wedge \lceil x \cdot y \rceil \cdot x \curvearrowright y \in \operatorname{Spa} \cdot \to : \operatorname{Dst} xy \cdot = \cdot \operatorname{Dst} \cdot \varphi x \cdot \varphi y :$ $\wedge : \uparrow x \cdot x \in \operatorname{Spa} \cdot \wedge : x = \varphi x : \in \operatorname{Pla} \mathsf{I}$ $\operatorname{Rfc} \neq \lceil \rceil *$

The word written Rfc means "reflections". Translations and rotations could be defined as products of reflections.

4 03 1. *
$$a \land b \in \operatorname{Spa} \land \land \alpha \land \beta \in \operatorname{Rea} \land \land \alpha + \beta = 1$$

 $\Rightarrow `a = b \Rightarrow \land \alpha a + \beta b = a$:

 $\land `a \neq b \Rightarrow `\alpha a + \beta b \in \operatorname{Ret} ab$:

 $\land `a \neq b \land \land \alpha \geq 0 \land \land \beta \geq 0 \Rightarrow `\alpha a + \beta b \in \operatorname{Cvx} a \cup b$:

 $\land `a \neq b \land \land \alpha \geq 0 \land \land \beta \geq 0 \Rightarrow `\alpha a + \beta b \in \operatorname{Cvx} a \cup b$:

 $\land `a \neq b \land \land \alpha \leq 0 \Rightarrow `\alpha a + \beta b \neq \operatorname{Cvx} a \cup b$:

 $\land `a \neq b \land \land \alpha \leq 0 \Rightarrow `\alpha a + \beta b \neq \operatorname{Cvx} a \cup b$:

 $\land `a \neq b \land \land \alpha \leq 0 \Rightarrow `\alpha a + \beta b \Rightarrow (\beta \mid \alpha \mid a \mid b \in \operatorname{Spa} \land \land \alpha \leq \beta \in \operatorname{Rea} \land \land \alpha \leq \beta \leq \operatorname{Rea} \land \land \alpha \leq \beta \leq \operatorname{Spa} \Rightarrow 1$)

4 03 2.
$$*a \Leftrightarrow b \in \operatorname{Spa} . \to : b - a . \in \operatorname{Cmt} \operatorname{Veo} \cdot \wedge \cdot b - a . \notin \operatorname{Spa} :$$
 $u \in \operatorname{Cmt} \operatorname{Veo} . \to . u \in u *$

' $u \in V$ co' means u is a vector. 'Cmt Vco' is the set of $u \in V$ co labelled with the unit Cmt.

4 03 3. *
$$a imes b imes c imes d \in \operatorname{Spa} . \to \cdot a - b = c - d$$
.
 $\longleftrightarrow 0.1 \times a + 0.1 \times d = 0.1 \times b + 0.1 \times c$:
 $a imes b imes c imes d \in \operatorname{Spa} . \to \cdot a - b = c - d$.
 $\to \cdot \operatorname{Dst} ab . = . \operatorname{Dst} cd \cdot \wedge : a = b . \wedge . c = d \cdot \vee \cdot \operatorname{Ret} ab . // . \operatorname{Ret} cd *$

4 03 4.
$$\# u \in \text{Cmt Vco.} \land .a \in \text{Spa}$$

 $\rightarrow .\text{Car:} ^b \cdot b \in \text{Spa.} \land .b - a = u = 1 #$

4 03 5. #
$$a ildas b \in \text{Spa.} \land u = b - a ildas b = a + u = u + a #$$

4 03 6. *
$$a \land b \land c \in \text{Spa} \land \land u = b - a \land \land v = c - b \land u + v = c - a \land \land \cdot v = b - c \land \land \cdot u + (-v) = u - v :$$

$$u \in \text{Cmt Vco} \rightarrow \cdot u - u = 0_{\text{Cmt Vco}} = \text{Cmt O}_{\text{vco}} *$$

4 03 7. *
$$a
ightharpoonup b \in \operatorname{Spa} \cdot \land \cdot \gamma \in \operatorname{Rea} \cdot \land \cdot u = b - a$$

$$\Rightarrow \cdot \gamma u = ((1 - \gamma)a + \gamma b) - a *$$

4 04 2. *
$$\gamma \in \text{Rea} \cdot \wedge \cdot a \land b \in \text{Rea}^3 \cdot \wedge \cdot a = \lceil \alpha_1 \cdot \alpha_2 \cdot \alpha_3 \rceil \cdot \wedge \cdot b = \lceil \beta_1 \cdot \beta_2 \cdot \beta_3 \rceil \cdot$$

$$\rightarrow : \gamma a = \lceil \gamma \alpha_1 \cdot \gamma \alpha_2 \cdot \gamma \alpha_3 \rceil \cdot \wedge \cdot a + b = \lceil \alpha_1 + \beta_1 \cdot \alpha_2 + \beta_2 \cdot \alpha_3 + \beta_3 \rceil \cdot$$

$$\wedge \cdot a - b = \lceil \alpha_1 - \beta_1 \cdot \alpha_2 - \beta_2 \cdot \alpha_3 - \beta_3 \rceil \cdot \wedge \cdot \lceil 0 \cdot 0 \cdot 0 \rceil = 0_3 \cdot$$

$$\wedge \cdot \text{aIb} = \alpha_1 \beta_1 + \alpha_2 \beta_2 + \alpha_3 \beta_3 \cdot \wedge \cdot a \perp b \cdot \leftrightarrow \text{aIb} = 0 \cdot$$

$$\wedge \cdot |a| = \sqrt{a \cdot 1} a \cdot \wedge \cdot \text{Dst } ab \cdot = |a - b| \cdot$$

$$\wedge \cdot \wedge \varphi : \langle ab \cdot = \varphi \cdot \leftrightarrow \cdot 0 \leq \varphi \leq \pi \cdot \wedge \cdot \cos \varphi = \text{aIb} / |a| |b| *$$

'aIb' means the scalar product of the numerical vectors a and b; $\not \lt$ is the symbol of the angle-function.

¹⁾ We have forgotten to define the absolute value of a number.

4 04 3. *
$$\gamma \in \text{Rea.} \land \cdot a \in \text{Rea3.} \land \cdot |a| = 1$$

 $\Rightarrow : \text{Pla}_* a \gamma := {}^*x : x \in \text{Rea3.} \land \cdot a \text{I} x = \gamma :$
 $a \land b \in \text{Rea3.} \land \cdot |a| = 1$
 $\Rightarrow : \text{Ret}_* ab = {}^*x : x \in \text{Rea3.} \land : \forall \tau : \tau \in \text{Rea.} \land \cdot x = \tau a + b :$
 $p \in \text{Pla}_* : \leftrightarrow : \forall {}^*a . \gamma : p = . \text{Pla}_* a \gamma :$
 $r \in \text{Ret}_* : \leftrightarrow : \forall {}^*a . b : r = . \text{Ret}_* ab :$
 $\text{Pla}_* a \gamma : \bot . \text{Pla}_* a' \gamma' : \leftrightarrow : a \text{I} a' : = 0 :$
 $\text{Ret}_* ab : \bot . \text{Ret}_* a' b' : \leftrightarrow : a' = 0 :$
 $\text{Pla}_* a \gamma : \bot . \text{Ret}_* a' b' : \leftrightarrow : a' = a . \forall a' = -a *$

4 04 4.
$$\# f \in \text{Coo} \longleftrightarrow \text{!} f \in \bigstar \curvearrowright \bigstar \land \land \land f : \text{Spa} = \text{Rea}^3$$
.
 $\land \land \land \land \land a \land b \land a \curvearrowright b \in \text{Spa} \rightarrow \text{!} \text{Dst } ab = \text{!} \text{Cmt} \cdot \text{Dst } .fa .fb \text{!}$

$$\forall f . f \in \text{Coo} \#$$

The word written Coo means "coordinate system", or rather "coordinatization".

4 04 5. *
$$f \in \text{Coo}$$
. \rightarrow : $\land \land \land a \cdot b \Rightarrow a \Rightarrow a \Rightarrow b \Rightarrow fa = fb$
 $\land : \land p \cdot p \in \text{Pla} \Rightarrow : \land : f \mid p \in \text{Pla} \Rightarrow :$
 $\land : \land r \cdot r \in \text{Ret} \Rightarrow : \land : f \mid r \in \text{Ret} \Rightarrow :$
 $\land : \land \land r \cdot s \Rightarrow : \land : f \mid r \Rightarrow : \land : f \mid s \Rightarrow :$

4 04 6. #
$$f \in \text{Coo} \cdot \land \cdot a \land b \in \text{Spa} \cdot \land \cdot \alpha \land \beta \in \text{Rea} \cdot \land \cdot \alpha + \beta = 1$$

$$\rightarrow f \cdot \alpha a + \beta b = \cdot \alpha f a + \beta f b : \land : f \cdot a - b = \cdot f a - f b #$$

4 04 7. # Cmt
$$\alpha \times$$
 Cmt $\beta =$ Cmt² $\alpha \times \beta$ #

4 04 8. *
$$u imes v \in \operatorname{Cmt} \operatorname{Veo} \land \land \gamma \in \operatorname{Rea} \land \land f \in \operatorname{Coo}:$$

$$\to f \cdot u + v := \cdot f u + f v : \land : f \cdot \alpha u := \cdot \alpha f u :$$

$$\land \cdot |u| = \operatorname{Cmt} |f u| \cdot \land \cdot u I v := : \operatorname{Cmt}^2 \cdot f u \cdot I \cdot f v :$$

$$\land : \not \preceq u v := \cdot \not \preceq \cdot f u \cdot f v :$$

4 04 9.
$$*Coo = .Coo^+ \cup Coo^- \cdot \wedge \cdot \lceil \rceil = .Coo^+ \cap Coo^- \cdot \wedge \cdot \varphi \in Rfc$$
.
 $\wedge \cdot \varphi \in Rfc$.
 $\rightarrow : f \in Coo^+ \cdot \rightarrow .f \varphi \in Coo^- \cdot \wedge \cdot f \in Coo^- \cdot \rightarrow .f \varphi \in Coo^+ \mid Coo^+ \neq \lceil \rceil \cdot \wedge \cdot Coo^- \neq \lceil \rceil *$

Coo+ and Coo- mean "positive" and "negative" coordinate systems (lefthand and righthand). The distinction between Coo+ and Coo- is purely verbal. At present we cannot do more than we have done, in spite of the efforts of a good many textbook writers. In 4 24 6 we shall fasten one element of Coo+ to real space. From that moment onwards Coo+ and Coo- will be materially distinguishable.

4 05 1.
$$*u \in \text{Vco} \rightarrow \text{Cmt } u \in \text{Cmt Vco}$$
:
 $u \in \text{Vco} \rightarrow u \in u$:
 $\text{Cmt } u = \text{Cmt } v \rightarrow u = v *$

4 05 2. *
$$u ext{ } \sim v \in \text{Vco.} \land \cdot \gamma \in \text{Rea} \cdot \rightarrow \text{: Cmt.} u + v \cdot = \text{: Cmt.} u + \text{Cmt.} v :$$

$$\land \cdot \text{Cmt.} \gamma u = \gamma \text{ Cmt.} u \cdot \land : \text{Cmt.} u \cdot \text{I. Cmt.} v \cdot = \text{Cmt}^2 \cdot u \text{Iv.} :$$

$$\land \cdot \swarrow uv \cdot = : \swarrow \cdot \text{Cmt.} u \cdot \text{Cmt.} v \cdot \land \cdot | \text{Cmt.} u | = \text{Cmt.} |u| *$$

4 05 3. *
$$u \in V_{CO} \rightarrow u = 0_{V_{CO}}$$
 *

4 06 1.
$$\# A \in Lap a_1 a_2 a_3$$

$$\leftrightarrow : A \subset \operatorname{Spa} . \land . a_1 \curvearrowright a_2 \curvearrowright a_3 \in \operatorname{Pos}.$$

$$\land ! \lor f : f \in \operatorname{Coo}^+ . \land : \land x : x \in A.$$

$$\leftrightarrow : \land i : i \in 1 \cup 2 \cup 3. \rightarrow . 0 \leq \omega_i f x \leq a_i I$$

$$A \in \operatorname{Lap} . \leftrightarrow : \lor \lceil a_1 . a_2 . a_3 \rceil . A \in \operatorname{Lap} a_1 a_2 a_3 #$$

The word written Lap (fL lapis = stone) means "rectangular parallelepiped"; a_1 , a_2 , a_3 mean the lengths of its edges.

4 06 2. * Vol. Lap
$$a_1 a_2 a_3 = \text{Cmt}^3 a_1 \times a_2 \times a_3 =$$

4 06 3. # Cmt³
$$\gamma$$
.+.Cmt³ δ = • Cmt³. γ + δ : Etc #

4 06 4.
$$* \lor B \cdot B \in \text{Lap.} \land .A \subset B \cdot \rightarrow !\alpha = \text{Vol } A \cdot \leftrightarrow ! \land \lceil M \cdot \varphi \rceil$$
:
$$M = : ^*X \cdot \text{Car } X \in \text{Num.} \land .X \subset \text{Lap.} \land .A \subset \bigcup X ^* \land ^*$$

$$\varphi = : ^*X \cdot \Sigma .x \in X \cdot \text{Vol } x \mid M :$$

$$\rightarrow `\alpha = .\inf ^* \varphi *$$

'Vol' means volume. We gave the definition of exterior content.

4 06 5.
$$\# \pi = 100 \times (1 - 1/11 + 1/101 - 1/111 + \text{Etc})$$

= 11,0010010000111111101101010001000111100111001 Etc:
 $A \in \text{.Bul } \alpha \rho \cdot \rightarrow \text{.Vol } A = 100/11 \times \pi \rho^{11} \#$

4 07 1.
$$\# x \in .$$
 Hom \cup Bes $\cdot \to \cdot \text{Vol } x = .$ Vol Loc x :
 $x \in \text{Hom } \wedge \cdot \wedge \text{Act } x = .$ Sec 10^{11001} :
 $\to \cdot \text{Vol } x = \text{Cca. Cmt}^3 \ 10^{1101}$:
 $x \in \text{Hom } \wedge \cdot \wedge \text{Act } x \ge .$ Sec 101×10^{11011} :
 $\to \cdot \text{Vol } x = \text{Cca. Cmt}^3 \ 10^{10000}$:
 $\forall x : x \in \text{Bes. } \wedge \cdot \text{Nnc } x \text{ Ext. } \wedge \cdot \text{Vol } x = \text{Cca. Cmt}^3 \ 11 \times 10^{11001}$:
 $\neg \forall x : x \in \text{Bes. } \wedge \cdot \text{Nnc } x \text{ Ext. } \wedge \cdot \text{Vol } x > \text{Cmt}^3 \ 11 \times 10^{11001}$:
 $\text{Car } \uparrow x : x \in \text{Bes. } \wedge \cdot \text{Nnc } x \text{ Ext. } \wedge \cdot \text{Vol } x \ge .$ Cmt $^3 \ 10^{10000}$
 $\therefore |\text{Car } \uparrow x : x \in \text{Bes. } \wedge \cdot \text{Nnc } x \text{ Ext. } \wedge \cdot \text{Vol } x < .$ Cmt $^3 \ 10^{10000}$:
 $\text{Mul} < . 1 \ \#$

Volume of humans and animals. — Perhaps the receiver will arrive at some crude estimation of our length unit on the strength of these data and the next.

4 07 2. $* \lor B \cdot B \in \text{Lap.} \land A \subseteq B$: $\rightarrow \text{!Lon } A \cdot = \text{!sup.} ^*x \cdot \lor ^{\vdash}a \cdot b ^{\lnot} : a \stackrel{\wedge}{\sim} b \in A \cdot \land \cdot x = \text{.Dst } ab *$

The word written Lon (fL longitudo=length) means "length, diameter".

- 4 07 3. $x \in .$ Hom \cup Bes $\cdot \to \cdot$ Lon x = . Lon Loc x : $x \in .$ Hom $. \wedge \cdot$ Act x = . Sec $10^{11001} : \to \cdot$ Lon x = . Cca . Cmt $10^{110} :$ $x \in .$ Hom $. \wedge \cdot$ Act $x \ge .$ Sec $101 \times 10^{11011} : \to$ \cdot Lon x = . Cca . Cmt $101 \times 10^{1011} :$ $\forall x : x \in .$ Bes $. \wedge .$ Nnc x = . Ext $. \wedge \cdot$ Lon x = . Cca . Cmt $10^{1100} :$ Car $\cdot \wedge x : x \in .$ Bes $. \wedge .$ Nnc x = . Ext $. \wedge \cdot$ Lon x = . Cmt $10^{1100} :$. Car $\cdot \wedge x : x \in .$ Bes $. \wedge .$ Nnc x = . Cmt $101 \times 10^{1011} :$. Mul < . 1 #
- 4 08 1. * $Ha \operatorname{Inq} Hb$ * $t + h \operatorname{Usq} t' + h' \cdot Hd \operatorname{Ani} \cdot \operatorname{Utr} \cdot tt' Hc \operatorname{Inq} Hd p : \land \cdot h \neq h'$ $\rightarrow \cdot t \operatorname{Loe} Hd \neq t' \operatorname{Loe} Hd \cdot \lor \cdot t \operatorname{Loe} Hc \neq t' \operatorname{Loe} Hc = t' \operatorname{Loe} Hc = t'$
- 4 08 2. * $t \operatorname{Dst} XY = : \operatorname{Dst} \cdot t \cdot \operatorname{Loc} X \cdot t \cdot \operatorname{Loc} Y *$
- 4 08 3. * $t X \text{ Mov.} \leftrightarrow : \land \alpha : \alpha \in \text{Pos.} \rightarrow : \lor \ulcorner t' \cdot t'' \urcorner :$ $t - \text{Sec } \alpha \leq t' \leq t \leq t'' \leq t + \text{Sec } \alpha :$ $\land : t \text{Loc } X \neq t' \text{Loc } X : \land : t \text{Loc } X \neq t'' \text{Loc } X *$

The word written Mov (fL movetur = moves) means the neutre verb "moves".

4 08 4. $t_1 t_2 X \text{ Mov.} \leftrightarrow : \land t \cdot t_1 < t < t_2. \rightarrow .tx \text{ Mov.}$ $t \text{ Usq } t' \cdot a \text{ Usd } a' \cdot X \text{ Mov.}$ $\leftrightarrow \cdot tt' X \text{ Mov.} \land .t \text{ Loc } X = a \cdot \land .t' \text{ Loc } X = a' \text{ }$

The word written Usd (fL usque ad=up to) is the spatial analogue of the temporal Usq.

- 4 08 5. $*t \operatorname{Usq} t' \cdot a \operatorname{Usd} a' \cdot X \operatorname{Mov} \cdot \wedge \cdot t' \operatorname{Usq} t'' \cdot a' \operatorname{Usd} a'' \cdot X \operatorname{Mov} :$ $\rightarrow \cdot t \operatorname{Usq} t'' \cdot a \operatorname{Usd} a'' \cdot X \operatorname{Mov} *$
- 4 09 1. # t₁Ha Inq Hb ' Nnc Ha Vul: Sec 10 . Pst Nnc · Hb Apu Ha:

 Hb Inq Ha .? = Apu:

 Ha Inq Hb ' Nnc Ha Vul: Sec 10 . Pst Nnc · Loc Hb .= Cca . Loc Ha:

 t₁ Loc Hb . Usd . t₁ Loc Ha . Hb Mov:

 Ha Inq Ha: Nnc Loc Hb .= Cca . Loc Ha · Hb Apu Ha

 PAN . Hb . Usd . Apu Ha . Hb Mov #

The word written Apu (fL apud=at) means (the local) "at, near". There is a notable difference between the use of Apu in the first and in the last sentence. This transition is akin to that in the use of words like MAN (see 3 20 6).

4 09 2. * $Ha \operatorname{Inq} Hb$ 'Nnc $Ha \operatorname{Vul}$: Sec 10 . Pst Nnc · $Hb \operatorname{Apu} Ha$: $Hb \operatorname{Inq} Ha$: — Pot Etc: Qia · Nnc Dst Ha Hb . > Nim '
Pot: Sec 1000 . Pst Nnc · $Hb \operatorname{Apu} Ha$ *

4 09 3. * t₁Ha Inq Hb 'Nnc Ha Vul : Sec 10 . Pst Nnc · Hb Apu Ha :

Hb Inq Ha : Nnc Hb Vul 'p . PAN Ha Vul p ·

Tan : ¬ Nnc Hb Vul · Nnc Hb Mov '

Sed : Nnc Hb Vul · t₁ Loc Ha . Usd . t₁ Loc Hb . Ha Mov *

4 09 4. # Ha Inq Hb: Nnc Ha Vul: Nnc · Nnc Loc Hb. Usd. Apu Hc. Hb Mov:

Nne · Nne Apu Hc . Usd . Apu Ha . Hb riangleq Hc Mov | Hd Inq Hd : PAN Ha Vul · PPN . Apu Hc . Hb Inq Hc Etc | Usd . Apu Hc · Hb Mov |

Apu Hc. Usd. Apu Ha: ${}^{\checkmark}w$. w Uni Hb Hc ${}^{\checkmark}Mov$ I Ha Inq Ha: Nnc ${}^{\diamond}Hb$ Apu Ha. ${}^{\wedge}$. Hc Apu Ha †

'PPN. Apu Hc. Hb Inq Hc Etc' is a paradigm which shows how to say that a certain event happens at a certain place. The general pattern would be

tu Fit p,

where t is a temporal and u a spatial adjunct. More examples could be given.

4 10 1. p and q in the next program texts will represent definite noises.

 $+ t_0 Ha \operatorname{Inq} Ba \mathfrak{p}$:

t₁Usd . Apu $Ha \cdot Ba \operatorname{Mov}^{t_2}$:

 $Hb \operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$

Hc Inq Hb: tot1 Ha Vul. PPN Ba Apu Ha

Pau Pst. Ha Vul Etc: Ba Ani · Utr. PAN Ha Vul Etc

 $t_1 \operatorname{Usq} t_2 \cdot t_1 \operatorname{Loc} Ba \cdot \operatorname{Usd} \cdot t_1 \operatorname{Loc} Ha \cdot Ba \operatorname{Mov} *$

Whistling for one's dog.

4 10 2. $(t_1 < t_2 < t_3 < t_4)$ * ${}^{t_1}Ha \operatorname{Inq} Ba \mathfrak{P}^{t_1} I$ * $Ba \operatorname{Inq} Ha \mathfrak{q}^{t_2} I$ * $Hb \operatorname{Inq} Hc \cdot ? x \cdot t_1 t_3 \operatorname{Fit} x I$ * $Hc \operatorname{Inq} Hb : t_1 t_2 Ha \operatorname{Vul} \cdot t_2 \cdot \operatorname{Usq} \cdot \operatorname{Unq} \cdot \operatorname{Pau} \operatorname{Pst} \cdot t_2 :$ * $t_1 \operatorname{Loe} Ba \cdot \operatorname{Usd} \cdot t_1 \operatorname{Loe} Ha \cdot Ba \operatorname{Mov} :$ * $\operatorname{Sed} : \neg t_1 t_3 Ba \operatorname{Vul} \cdot {}^{\checkmark}q \cdot \operatorname{Nnc} Ha \operatorname{Vul} q \cdot \operatorname{Erg} \cdot \operatorname{Nnc} Ha \operatorname{Mov} I$

```
Hb Inq Hc '? u: Nnc ·t<sub>1</sub> Loc Ha . Usd u · Ha Mov I
Hc Inq Hb ' Nnc Ha Vul · PPN . Ha Apu Ba :
```

Erg · Nnc . Adv Ba . Ha Mov I

 $Hb \operatorname{Inq} Hc \cdot ? = \operatorname{Adv} I^{t_{\bullet}}$

Hc Inq Hb : Nnc Ha Vul · PPN Dst Ha Ba . < . Nnc Dst Ha Ba . .

Erg . Nes . $\wedge t : t = \text{Cea. Nnc} \rightarrow : \forall A :$

 $A \subset : Cvx \cdot t \operatorname{Loc} Ha \cdot \cup \cdot t \operatorname{Loc} Ba$

 \land : Cca $t \cdot t$ Loc Ha . Usd $A \cdot Ha$ Mov |

Hb Inq Hc · Utr . Nnc Ha Apu Ba I

Hc Inq Hb . Fal : Qia : Cca t4: Ba Ani : Utr · Nnc . Adv Ba . Ha Mov ·

^ 'Vul . ¬ PPN Ha Apu Ba:

Erg · Pau Pst t4. Usq Nnc. Dev Ha. Ba Mov I

 $Hb \operatorname{Inq} Hc.? = \operatorname{Dev} I$

Hc Inq Hb .. Nnc Ba Vul . — UPN Ha Apu Ba: Eti:

Nne Ba Vul · PPN Dst Ha Ba. > . Nne Dst Ha Ba

Erg | Nes : $\land t : t = \text{Cea} \cdot \text{Nnc} \rightarrow : \lor \ulcorner a \cdot b \cdot c \urcorner$

 $a \in \operatorname{Spa} . \land . b \in \operatorname{Loc} Ba . \land . c \in \operatorname{Loc} Ha . \land : b \in \cdot \operatorname{Cvx} . a \cup c : \land :$

 $\operatorname{Cea} t \cdot t \operatorname{Loc} Ba \cdot \operatorname{Usd} a \cdot \operatorname{Ha} \operatorname{Mov} \mathbf{I}$

 $\operatorname{Nnc} \operatorname{Dev} Ha \operatorname{Ba} \operatorname{Mov} *$

The words written Adv (fL adversus = towards) and Dev (fL deversus = away) mean "towards" and "away" respectively.

4 10 3. # $Ha \operatorname{Inq} Ha : \operatorname{Car}^{x} x \in \operatorname{Bes} . \land . \operatorname{Nnc} x \operatorname{Ext} . \land :$

 $\forall p \cdot \text{Pot. Nnc } x \text{ Ani } p$:

 $\wedge : \forall p : \text{Pot.Nnc} x \text{ Vul } p : \wedge : \text{Pot.Nnc} x \text{ Mov} :$

 $Mul > Car : x \cdot x \in Hom \cdot \land Nnc x Ext *$

4 10 4. $* t_1 Ha Inq Hb . Ha$:

Hb Inq Ha. Hbta:

 $Hc \operatorname{Inq} Hd \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$

 $Hd \operatorname{Inq} Hc : \operatorname{Cea} t_1 : Ha \operatorname{Ani} : ? = \operatorname{Loc} Hb :$

 $Argamma : Ha Vul \cdot Nnc Hb Ani \cdot ? = Loc Ha :$

Pau Pst. t₁ · Obv Ha Hb Mov :

 $Hc \operatorname{Inq} Hd$.? = Obv:

Hd Inq Hc Pau Pst t1: Loc Ha. Adv. Loc Hb · Ha Mov:

Loc Hb. Adv. Loc Ha · Hb Mov #

The word written Obv (fL obviam = going to meet) means "going to meet".

4 10 5. $\# t_1 t_1 \operatorname{Loc} Ba \cdot \operatorname{Usd} t_1 \operatorname{Loc} Ha \cdot Ba \operatorname{Mov}^{t_1} I$

 $Hb \operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x$

Hc Inq Hb · PAN Ba Vul . PPN Ba Apu Ha I

Hb Inq Hc Cur I

```
Ha Inq Hb: Qia: PAN Ba Sci:
                          ? y' Nne: Adv Ha. Ba Mov \rightarrow Pau Pst. Ha Inq Hb y I
                     Hb \operatorname{Inq} Hc. ? y \operatorname{Etc} I
                    Hc\operatorname{Inq} Hb. Etc Ha\operatorname{Inq} Hb\operatorname{Ben} *
4 11 1. # t1 pt2 1
                    Hb \operatorname{Inq} Hc \cdot ?x \cdot t_1 t_2 \operatorname{Fit} x
                    Hc \operatorname{Inq} Hb : \forall y \operatorname{Put} Ra \operatorname{PAN} Ha \operatorname{Vul} :
                              Nne Usq PPN · Nne Loe y . Usd . Nne Loe Ha \cdot y Mov :
                              \operatorname{Mod} \cdot \operatorname{Nne} Ha \operatorname{Inq} y \mathfrak{p}:
                          Qqm : \neg Pot : \forall t \cdot Cea \ t \ Ra \ Vul \cdot Nne \ Ra \ Mov
                          Qia: Ra \in \text{Res. Sed} \cdot Ra \notin . \text{Hom} \cup \text{Bes}:
                          \neg t_1 t_2 Ha \operatorname{Sci} \cdot \operatorname{Utr} \cdot Ra \in \operatorname{Res} :
                          \operatorname{Tan}:\operatorname{Cca} \operatorname{t}_1 \operatorname{Ha} \operatorname{Inq} \operatorname{Ha} \cdot \operatorname{Ra} \in \operatorname{Bes} \cdot \operatorname{Qqm} \operatorname{Fal}^{\perp}
                          \operatorname{Qqm} \cdot \operatorname{Loc} Ra \cdot \operatorname{CSpa} : \operatorname{Tan} \cdot Ra \notin \operatorname{Hom} \cup \operatorname{Bes} :
                          \neg \operatorname{Pot} \cdot \vee \lceil u \cdot p \rceil \cdot u \operatorname{Ra} \operatorname{Vul} p:
                         y \in \text{Res.} \rightarrow : \neg \text{Pot} \cdot y \text{Inq.} \lor .y \text{Sci.} \lor .y \text{Ani.} \lor .y \text{Vull}
                    t. Hb Inq Hc: Nnc · PAN Loc Ra. Usd. PAN Loc Ha · Ra Mov I
                    Hc \operatorname{Inq} Hb . A \operatorname{Inq} Ha \operatorname{Mov} \operatorname{Ceat}_1. Hd \operatorname{Apu} Ra
                         Pau Pst. t<sub>2</sub> · Hd Ani 'Utr: Nnc Ha Vul·
                              Nnc. Usd. Apu Ha. Ra Mov:
                          \forall D ' D Uni Hd Ra . \land: t_3 Usq t_4 ' t_3 Loc D . Usd . t_3 Loc Ha ' D Mov :
                         t_3 \operatorname{Usq} t_4 \cdot \operatorname{Loc} D := \operatorname{Loc} Hd \cdot \cup \operatorname{Loc} Ra
                         t_3 t_4 D  Ext: Pau Pst. t_4 \cdot t_4  Loc Hd \cdot Usd \cdot t_3  Loc Hd \cdot Hd  Mov:
                         Tan . — Nnc Ra Mov \vdots t_3 t_4 Hd Cau . Ra Mov \vdots
                         y \in \text{Res.} \rightarrow \neg \text{Pot.} y \text{ Vul.} y \text{ Mov.}
                              Sed Pot \forall x : x \in \text{Hom} \land x \text{Vul} \cdot y \text{Mov}
                         x \in \text{Hom} . \land . y \in \text{Res}
                              \rightarrow Pot t: x \text{ Vul} . Nnc y \text{ Mov} \rightarrow Pau Pst . y \text{ Mov} +
```

The word written Res (fL res = thing) means "concrete things". Ra is the name of a special thing. 'Cau' (fL causa = cause) means (the verb) "causes".

With respect to variables controlled by it Cau behaves as do words like Inq. (Oblique speech.)

A person can cause a thing to move.

```
4 11 2. * Ra 
ightharpoonup Rb 
ightharpoonup Rc 
ightharpoonup Rd 
ightharpoonup Etc 
ightharpoonup Res . 
ightharpoonup
```

4 11 3. * $Ha \operatorname{Inq} Hb$ `Nne $Ha \operatorname{Vul}: \operatorname{PPN} \cdot \check{\ \ } x$. Nne $x \operatorname{Apu} Hb \cdot \operatorname{Apu} Ha$: ${}^{t_1}\operatorname{Cea} \cdot t_1 \operatorname{Loe} Rb \cdot \operatorname{Usd} \cdot \operatorname{Cea} \cdot t_1 \operatorname{Loe} Ha \cdot \check{\ \ \ } y \cdot y \operatorname{Uni} Hb Rb \cdot \operatorname{Mov}^{t_a}$: $Hc \operatorname{Inq} Hd \cdot ? z \cdot t_1 t_2 \operatorname{Fit} z$:

 $Hd \operatorname{Inq} Hc \cdot t_1 t_2 Hb \operatorname{Cau} \cdot Rb \operatorname{Mov} \cdot \operatorname{Mod} \cdot Hb \operatorname{Fer} Rb :$ $t_1 t_2 \cdot \operatorname{Cea} \cdot t_1 \operatorname{Loc} Rb \cdot \operatorname{Usd} \cdot \operatorname{Cea} \cdot t_1 \operatorname{Loc} Ha \cdot Hb \operatorname{Fer} Rb *$ The word written Fer $(fL \text{ fert} = \operatorname{carries})$ means (the verb) "carries".

4 11 4. $Ha \operatorname{Inq} Hb$ 'Nne $Ha \operatorname{Vul}: \operatorname{PPN} \cdot \check{}_{x}$. Nne $x \operatorname{Apu} Hb \cdot \operatorname{Apu} Ha$: $^{t_{1}}\operatorname{Cea} \cdot t_{1} \operatorname{Loe} Rc \cdot \operatorname{Usd} \cdot \operatorname{Cea} t_{1} \cdot \operatorname{Loe} Ha \cdot Rc \operatorname{Mov}^{t_{2}}$: $Hc \operatorname{Inq} Hd \cdot ?z \cdot t_{1} t_{2} \operatorname{Fit} z$: $Hd \operatorname{Inq} Hc \cdot t_{1} t_{2} Hb \operatorname{Cau} \cdot Rc \operatorname{Mov} \cdot \operatorname{Tan} : \neg \operatorname{Mod} \cdot Hb \operatorname{Fer} Rc \cdot \operatorname{Sed} \cdot \operatorname{Mod} \cdot Hb \operatorname{Iac} Rc \cdot \operatorname{Qia} \cdot \neg t_{1} t_{2} Hb \operatorname{Mov} *$ The word written $\operatorname{Iac} (fL \operatorname{iactat} = \operatorname{throws})$ means "throws, pushes".

- 4 11 5. * $Ha \operatorname{Inq} Hb$ 'Nne $Ha \operatorname{Vul}$: Nne · Apu Hb . Usd . Apu $Ha \cdot Hb$ Fer Rd! $Hb \operatorname{Inq} Ha$: Pot 'Qia: $\operatorname{Vol} Rd \cdot \operatorname{Vol} Rd \cdot \operatorname{Mul} > \cdot \operatorname{Nim}$ 'Pot: ' $D \cdot D \operatorname{Uni} \cdot 11^e \operatorname{Hom} \cdot \operatorname{Fer} Rd = \cdot \operatorname{Vol} Rd \cdot \operatorname{Mul} > \cdot \operatorname{Nim} Rd \cdot \operatorname{Vol} Rd \cdot \operatorname{Mul} > \cdot \operatorname{Nim} Rd \cdot \operatorname{Mul} > \cdot \operatorname{Mul} = \cdot \operatorname{Mul} Rd \cdot \operatorname{M$
- 4 11 6. * t_1 Loc Rf . Usd . t_1 Loc $Ha \cdot Rf$ Mov t_2 : $Hc \operatorname{Inq} Hd \cdot ? z \cdot t_1 t_2$ Fit z : $Hd \operatorname{Inq} Hc : t_1 Hb \operatorname{Vul} : \operatorname{Nnc} \cdot \operatorname{Apu} Hb$. Usd . Apu $He \cdot Rf$ Mov \cdot $\wedge \cdot t_1 \cdot \operatorname{Adv} He \cdot Hb \operatorname{Iac} Rf :$ Tan $t_2 Ha \operatorname{Cau} : \neg \operatorname{Unq} \operatorname{PPN} \cdot \operatorname{Qqm} \operatorname{Nnc} \cdot Rf \operatorname{Mov} *$

The Nnc after Cau is bound to approximately t_2 , not to the moment when Hd is speaking. See the remark in 4 11 1.

4 11 7. * $^{t_1}t_1 \operatorname{Loc} Rf \cdot \operatorname{Usd} \cdot t_1 \operatorname{Loc} He \cdot Rf \operatorname{Mov}^{t_2}$: $Hc \operatorname{Inq} Hd \cdot ?z \cdot t_1 t_2 \operatorname{Fit} z :$ $Hd \operatorname{Inq} Hc \cdot t_1 Hb \operatorname{Vul} : \operatorname{Nnc} \cdot \operatorname{Apu} Hb \cdot \operatorname{Usd} \cdot \operatorname{Apu} He \cdot Rf \operatorname{Mov} \cdot$ $\wedge \cdot t_1 \cdot \operatorname{Adv} He \cdot Hb \operatorname{Iac} Rf :$ $Hc \operatorname{Inq} Hd \cdot \operatorname{Cur} : \neg \operatorname{PAN} Ha \operatorname{Cau} \cdot \neg \operatorname{Nnc} Rf \operatorname{Mov} :$ $Hd \operatorname{Inq} Hc \cdot \operatorname{Qqm} \cdot \operatorname{PAN} Ha \operatorname{Vul} \cdot \operatorname{Nnc} Ha \operatorname{Cau} \operatorname{Etc} :$ $\operatorname{Tan} \cdot \wedge t : t_1 \le t \le t_2 \cdot \rightarrow \cdot t \operatorname{Dst} Ha Rf \cdot > \operatorname{Nim} *$

4 11 8. $*^{t_1}Ha\ Hb\ Lud_6^{t_4}$ $Hc\ Inq\ Hd\cdot?\ x.\ t_1\ t_4\ Fit\ x$ $Hd\ Inq\ Hc\colon t_1\ t_4\ Ha\ Hb\ Lud_6$:

Deb : t_1 . Adv Hb. $Ha\ Iac\ Rg$ $t_2\ Hb\ Cau\cdot \neg\ Pau\ Pst\ t_2$. $Rg\ Mov$ $t_3\ Adv\ Ha$. $Hb\ Iac\ Rg$ $t_4\ Hb\ Cau\cdot \neg\ Pau\ Pst\ t_4$. $Rg\ Mov$ $t_2\colon \neg\ .Hb\ Cau\ Etc\cdot \rightarrow Pau\ Pst\cdot Hb\ Dat\ Ha$. Den 10 $t_4\colon \neg\ .Ha\ Cau\ Etc\cdot \rightarrow Pau\ Pst\cdot Ha\ Dat\ Hb$. Den 10 *Playing ball.

4 12 1. * Dst.
$$t_1$$
 Loc x . t_2 Loc x · = Cmt δ : \wedge . t_2 - t_1 = Sec τ . \wedge . τ > 0 · \Rightarrow : t_1 Usq t_2 Cel Med x Mov. = · Cmt/Sec. δ/τ : $\delta \in \text{Rea.} \wedge \cdot \tau \in \text{Pos.} \rightarrow : \text{Cmt.} \delta/\text{Sec.} \tau = \cdot \text{Cmt.}/\text{Sec.} \delta/\tau$: $\alpha \Rightarrow \tau \in \text{Rea.} \rightarrow \cdot \text{Cmt.}/\text{Sec.} \alpha \times \text{Sec.} \tau = \cdot \text{Cmt.} \alpha\tau$: $\alpha \Rightarrow \beta \in \text{Rea.} \rightarrow \cdot \text{Cmt.}/\text{Sec.} \alpha = \cdot \text{Cmt.}/\text{Sec.} \beta$: $\leftarrow \cdot \alpha = \beta$: Etc. *

The word written Cel Med (fL celeritas media = mean velocity) means "mean velocity".

4 12 2. #
$$\tau \sim \nu \in \text{Pos.} \rightarrow \cdot \text{Sec}^{-1} \nu \times \text{Sec } \tau . = \nu \tau$$
:

$$\text{Sec}^{-1} \nu = \text{Sec}^{-1} \nu_1 . \longleftrightarrow . \nu = \nu_1$$
:
Etc #

The word written Osc (fL oscillat = oscillates) means "oscillates".

4 13 2. * Aml·
$$t_1$$
 Usq t_2 . a . z Osc $_{\tau\alpha\nu}$: = a .

Fre· t_1 Usq t_2 . a . z Osc $_{\tau\alpha\nu}$: = ν .

Pha· t_1 Usq t_2 . a . z Osc $_{\tau\alpha\nu}$: = τ .

Eql· t_1 Usq t_2 . a . z Osc $_{\tau\alpha\nu}$: = a .

The words written Aml (fL amplitudo), Fre (fL frequentia), Pha, Eql (fL aequilibrium) mean "amplitude", "frequency", "phase", and "equilibrium" respectively.

4 13 3. #
$$t \cdot z \operatorname{Osc} \longleftrightarrow \bigvee \lceil a \cdot t_1 \cdot t_2 \cdot \tau \cdot \alpha \cdot \nu \rceil : t_1 \leq t < t_2 \cdot \wedge \cdot t_1 \operatorname{Usq} t_2 \cdot a \cdot z \operatorname{Osc}_{\tau \alpha \nu} #$$

And, Fre, Pha, Eql are defined as in 4132.

4 13 4. #
$$R \in \text{Res.} \rightarrow \text{i.} t_1 \text{Usq} t_2 R \text{Vib}_{\nu} \cdot \leftrightarrow \text{i.}$$

$$\forall A_1 : \text{Car } A \cdot \text{Mul} > \cdot 1 \cdot \wedge \cdot R \text{Uni } A \cdot \wedge \text{i.}$$

$$\land z_1 \cdot z \in A \cdot \rightarrow \text{i.} z \text{ Par } R \cdot \wedge \cdot \text{Vol } z \cdot \text{Mul} < \cdot \text{Cmt}^3 1 \cdot \\
 \land \vdots \land t \cdot t_1 \leq t \leq t_2 \cdot \rightarrow \cdot t \cdot z \text{ Osc} \cdot \wedge \cdot \text{Fre} \cdot t \cdot z \text{ Osc} := \nu \cdot \\
 \land \land \text{Aml} \cdot t \cdot z \text{ Osc} \cdot \text{Mul} < \cdot \text{Lon } R \mid \\
\text{Exg i: } Ra \in \text{Res.} \land \cdot \nu > \text{Sec}^{-1} 0 \cdot \wedge \cdot t_0 \in \text{Tem.} \\
 \land \cdot \cdot \vee f \cdot f \in \text{Coo.} \land \vdots \land z \cdot z \text{ Par } Ra \cdot \\
 \rightarrow \vdots \lor \lceil x \cdot \alpha \rceil \cdot \alpha > 0 \cdot \wedge \cdot \wedge t : t Ra \text{ Ext.} \rightarrow \cdot \land f \mid t \text{ Loc } z \cdot = \text{Cca.} \\
 \lceil x + \alpha \cos 10 \pi \nu (t - t_0) \cdot 0 \cdot 0 \rceil \mid \cdot \\
 \rightarrow \cdot \land t : t Ra \text{ Ext.} \rightarrow \cdot t Ra \text{ Vib}_{\nu} \mid$$

```
Exg: Ra \in \text{Res.} \land l \land \varepsilon \land \gamma \land \alpha \in \text{Pos.} \land \varepsilon < \text{Vld.}
\land t_0 \land t_1 \in \text{Tem.} \land .10 \pi \nu (t_1 - t_0) \in \text{Num.} \land \iota :
\lor f \cdot f \in \text{Coo.} \land . \land f \mid t_0 \text{ Loc } Ra = : \uparrow x \cdot x = \lceil \xi_1 \cdot \xi_2 \cdot \xi_3 \rceil \cdot \land \cdot
0 \leq \xi_1 \leq l \cdot \land . \mid \xi_2 \mid \leq \varepsilon \cdot \land . \mid \xi_3 \mid \leq \varepsilon \mid \land \iota .
\land z \cdot z \text{ Par } Ra \cdot \Rightarrow : \lor y : \land t \cdot t \in \text{Tem.} \Rightarrow :
t_0 \leq t - \text{Sec } \gamma y \leq t_1 \cdot \Rightarrow \cdot \land f \mid t \text{ Loc } z = \text{Cca.}
\lceil y + \alpha \sin 10 \pi \nu (t - t_0 - \text{Sec } \gamma y) \cdot 0 \cdot 0 \rceil :
\land : \neg \cdot t_0 \leq t - \text{Sec } \gamma y \leq t_1 \cdot \Rightarrow \cdot \land f \mid t \text{ Loc } z = \text{Cca.} \lceil y \cdot 0 \cdot 0 \rceil \iota :
\Rightarrow : \land t \cdot t_0 \leq t \leq t_0 + \text{Sec } \gamma l \cdot \Rightarrow : \lor w \cdot w \text{ Par } Ra \cdot \land \cdot \text{Cca } t w \text{ Vib}_{\nu} *
```

The word written Vib (fL vibrat = vibrates) means "vibrates". Examples of stationary and progressive waves have been given.

```
4 14 1. # Ha \operatorname{Inq} Ha : X \in \operatorname{Hom} . \land . R \in \operatorname{Res} . \land `
t X \operatorname{Cau} : t \operatorname{Usq} t' . a \operatorname{Usd} a' . R \operatorname{Mov} \cdot \operatorname{Mod} X \operatorname{Iac} :
\rightarrow \cdot \operatorname{Cel} \operatorname{Med} R \operatorname{Mov} . \operatorname{Mul} < . \operatorname{Cmt} / \operatorname{Sec} 10^{1111} #
```

4 14 2. * $Ha \operatorname{Inq} Hb : \operatorname{Cur} X \in \operatorname{Hom} \cup \operatorname{Bes} \to \operatorname{Pot} Loc X \neq t' \operatorname{Loc} X :$ $Hb \operatorname{Inq} Ha : \operatorname{Qia} : \operatorname{Pot} X \operatorname{Vul} t' X \operatorname{Mov} :$ $Ha \operatorname{Inq} Hb : \operatorname{Cur} R \in \operatorname{Res} \to \operatorname{Pot} t \operatorname{Loc} R \neq t' \operatorname{Loc} R :$ $Hb \operatorname{Inq} Ha : \operatorname{Qia} \cdot \operatorname{Pot} t' R \operatorname{Mov} *$

The last text is to prepare the way for the next, enabling the receiver to understand the next 'Cur' in the sense of physically explaining causality.

```
4 14 3. * Ha \operatorname{Inq} Hb \cdot \operatorname{Pon} X \wedge Y \in \operatorname{Hom} A
                        Cur Pot t X Inq Y p \rightarrow Pau Pst:
                            Y \text{ Ani } \cdot \text{Utr. PAN } X \text{ Inq } Y p
                        \mathbf{Qqm} \cdot t \operatorname{Dst} XY \cdot \mathbf{Mul} > \mathbf{Cmt} 0 \mathbf{I}
                   Hb \operatorname{Inq} Ha \cdot \operatorname{Qia} \cdot \cdot \vee w \cdot \operatorname{Put} w \cdot \wedge t : t \in \operatorname{Tem}.
                        \rightarrow Pot : Cca t. X Cau 'Nnc. Apu X. Usd. Apu Y. w Mov
                            \sim : w \text{ Fer } \stackrel{\vee}{\cdot} x \cdot \text{Nnc } X \text{ Inq } Yx :
                            \triangle : Y \text{ Ani } \cdot \text{Utr } \cdot \text{Nnc } w \text{ Apu } Y \text{ I}
                   Ha \operatorname{Inq} Hb \cdot ? = w I
                   Hb \operatorname{Inq} Ha : \operatorname{Pon} : \operatorname{Cel} \operatorname{Med} : w \operatorname{Mov} : \operatorname{Mul} < \cdot \operatorname{Cmt/Sec} 10^{1111}
                        Erg w \in \text{Res.} \land : \text{Pot. Apu } X \cdot \text{Usd. Apu } Y \cdot w \text{ Mov.}
                        Pon . Cel Med . w Mov · = Cea . Cmt/Sec 10^{1111} .
                        \neg \lor A : A \in \text{Res.} \land \text{Pot Apu } X \cdot \text{Usd. Apu } Y \cdot A \text{ Mov}
                            Mod: Cel Med \cdot A Mov = Cea \cdot Cmt/Sec 10^{1111}
                        Pau Pst. t \cdot R Uul ' \wedge: Pau Pst. t \cdot Y Ani · Utr. Nnc R Uul I
                    Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot w = \operatorname{Aer} \cdot \vee \cdot w \operatorname{Par} \operatorname{Aer} \mathsf{I}
```

```
Hb \operatorname{Inq} Ha : \operatorname{Fal} : \operatorname{Qia} : \neg \operatorname{Pot} : v \operatorname{Par} \operatorname{Aer} \cdot \wedge \cdot t v \operatorname{Fer} \cdot x \operatorname{Etc} \cdot \operatorname{Mod}
         t \text{ Cel Med } v \text{ Mov} = \text{ Cel Med } w \text{ Mov } I
 Ha \operatorname{Ing} Hb : \operatorname{Utr} \cdot t \operatorname{Loc} w \cdot \subset \cdot t \operatorname{Loc} \operatorname{Aer} \mathsf{I}
 Hb \operatorname{Ing} Ha ' \operatorname{Ver} \cdot \cdot \vee M | \operatorname{Put} M \cdot \operatorname{Car} M > \operatorname{Vld} \cdot \wedge \vdots
          \land Z: Z \in Hom \cup Bes \rightarrow Car: \uparrow z \cdot z \in M \land z \land pu Z > Vld \land \uparrow
          \land z : z \in M \rightarrow z \operatorname{Par} \operatorname{Aer} \land \operatorname{Vol} z < \operatorname{Vld} \land
          \land \lceil z \cdot z' \cdot t \rceil \lceil z \land z' \in M \cdot \land \cdot z \operatorname{Apu} z' \cdot \land \cdot t \in \operatorname{Tem} 
                   \rightarrow : t \cdot z \text{ Mov } \rightarrow \text{Pau Pst. } z' \text{ Mov } :
         Pon z_1 
ightharpoonup z_2 
ightharpoonup Etc 
ightharpoonup z_N 
ig
                 z_1 \operatorname{Apu} X \cdot \wedge \cdot z_2 \operatorname{Apu} z_1 \cdot \wedge \cdot z_3 \operatorname{Apu} z_2 \cdot \wedge \cdot \operatorname{Etc} \cdot \wedge \cdot Y \operatorname{Apu} z_N \cdot \wedge
                   Loc z_1 . \cup .Loc z_2 . \cup .Etc . \cup .Loc z_N \cdot \subset
                 : Cvx \cdot Loc X \cdot \cup .Loc Y :.
         \operatorname{Erg} : t : z_i \operatorname{Mov} \to \operatorname{Pau} \operatorname{Pst} : z_{i+1} \operatorname{Mov} : \wedge i
                   \forall h : \land t : t \in \text{Tem} . \land : t : X \text{ Cau } . t z_1 \text{ Mov} \rightarrow h \text{ Pst} :
                 z<sub>N</sub> Mov. ~ Y Ani. Utr Etc ::
         Erg :: Pot: t Aer Uul. A. t X Cau. Nnc Aer Uul. A.
                   \forall \alpha .. \land \lceil X . Y . t \rceil . X \land Y \in \text{Hom} . \land . t \in \text{Tem} \rightarrow \exists
                           \forall h: h = \alpha \times t \operatorname{Dst} XY . \wedge :
                                   t X \operatorname{Cau} . Nnc Aer Uul \rightarrow h \operatorname{Pst} Pot : Y \operatorname{Ani}
                                            Utr. Nnc Aer Uul I
 Ha \operatorname{Inq} Hb \cdot ? = w \mathsf{I}
{}^{t_{\bullet}}Hb \operatorname{Ing} Ha \cdot w = {}^{t_{\bullet}}'t {}^{t_{\bullet}}v {}^{t_{\bullet}}v \operatorname{Uni} : {}^{t_{\bullet}}z \cdot z \in M . \land .tz \operatorname{Mov} \operatorname{Tem} :
         t \operatorname{Loc} w = t \operatorname{Loc} v v \operatorname{Uni} z z \in M \wedge tz \operatorname{Mov} t
         w \notin \text{Res. Sed}: w \in . \text{Und Aer.} \subset \text{Und}
         u \in \text{Und Aer.} \rightarrow : \text{Cel Med.} u \text{ Mov.} = \text{Cea. Cmt/Sec } 10^{1111}:
         Eti : \forall T : T \in \text{Res.} \land \cdot \text{Pot.} T \text{ Uul.} \land : \land u : u \in \text{Und } T
                 \rightarrow: Cel Med. u Mov. = Cea. Cmt/Sec 10^{10011} |
 Ha \operatorname{Ing} Hb : ? p \cdot R \in \operatorname{Res} A \cdot \operatorname{Pot} R \operatorname{Uul} A z \operatorname{Par} R :
         \rightarrow: Pot · z Mov . Mod p I
 Hb \operatorname{Inq} Ha : \operatorname{Exg} {}^{\bullet} Ra \in \operatorname{Res.} \wedge \cdot \operatorname{Pot.} Ra \operatorname{Uul} {}^{\bullet} \wedge \cdot z \operatorname{Par} R :
         \rightarrow: Pot z Osc. Mod {last part of 4 13 4}
 Ha \operatorname{Inq} Hb: ? u \cdot \operatorname{Cea} t_0 \cdot Ra \operatorname{Uul} \cdot \operatorname{Mod} \{ \operatorname{last} \operatorname{part} \operatorname{of} 4 \cdot 13 \cdot 4 \}.
         \rightarrow : u \in \text{Und } Ra . \land \cdot \text{Cea } t_0 . u \text{ Mov } I
Hb \operatorname{Inq} Ha : \wedge t : t_0 \leq t \leq t_0 + \operatorname{Sec} \gamma l.
                 \rightarrow '' f t \operatorname{Loc} u = : {}^{\uparrow}x \cdot x = {}^{\lceil}\xi_1 \cdot \xi_2 \cdot \xi_3 {}^{\rceil} \cdot \wedge \cdot 0 \leq \operatorname{Sec}^{-1}(t-t_1)/\gamma
                   \leq \xi_1 \leq \operatorname{Sec}^{-1}(t-t_0)/\gamma \leq l \cdot \wedge \cdot \mid \xi_2 \mid \leq \varepsilon \cdot \wedge \cdot \mid \xi_3 \mid \leq \varepsilon
         \rightarrow u \in U \operatorname{nd} Ra \cdot \wedge \cdot \operatorname{Cea} t_0 \cdot u \operatorname{Mov} \cdot \wedge
         {}^{\bullet}Cca t_0. Cel Med {}^{\bullet} u Mov := . Cmt/Sec \gamma I
 Ha \operatorname{Inq} Hb : \operatorname{Pon} : \operatorname{Cel} \operatorname{Med} \cdot w \operatorname{Mov} \cdot
         Mul > Cmt/Sec 10^{10011} Cur \cdot Pot w Mov I
```

The word written Uul (fL undulat=undulates) means "undulates" in the original sense such as Huygens took it (no underlying oscillation being supposed). The word written Und (fL unda=wave) means "wave". (This notion cannot be expressed by Uulio.) The words written Aer and Ath (fL aer = air, and aether = ether) mean "air" and "field" respectively. The precise meaning of these notions may remain unknown to the reveicer. The only thing that matters is the existence of some media that admit progressing waves with certain specified velocities.

'c' is a constant, the velocity of light in vacuo.

'{...}' stands for a quotation or repetition of the program text indicated between the braces.

In the next paragraph we shall account for the purpose of the last talk.

4 14 4. Introducing length is as difficult as introducing time was easy. In the beginning of the present chapter a word Cmt appeared. The receiver could guess that it means some length unit, but no indication was given of its value. It has been fixed by the text given in the last paragraph. This goal has been reached after long detours. Most of the developments of this chapter have merely been steps on this lengthy way.

Of course there are other possible contexts for presenting our length unit. Natural yardstick of astronomic size are the curvature radius of the universe, which is still extremely uncertain, and the parsec (=1296000/2 π times the distance between sun and earth), which is much better known, though not with great precision, but which can be significant only for a receiver inside our solar system. There are many microscopic yardsticks, but as long as we have not composed a vocabulary of chemistry or atomic physics, we cannot present them. As far as I can see, the only exception is the

phenomenological Rydberg constant R of hydrogen (or some other hydrogen-like element). It has the dimension of the inverse of a length, and it is one of the best-known physical constants (if not the best of all), with a precision of 10^{-7} . We could introduce it by its historical pre-Bohrian definition. We could say that there is a line spectrum with the reciprocal wave-lengths

$$\frac{1}{\lambda} = R\left(\frac{1}{m^2} - \frac{1}{n^2}\right); m < n, \text{ integers.}$$

But this cannot be done until something has been said about light. At this stage of the program the most striking feature of light, through which it could be characterized, is its tremendous velocity. On the other hand, as soon as we can say what light is, we can also present a reasonably satisfactory yardstick: the light-second, i.e. the distance light travels in the time-unit (known by mere presentation).

My original idea was to define velocity of light merely as the upper bound of the velocities of corpuscular inertial motions. Yet serious objections may be raised against this procedure. It would not be wise to start with relativistic mechanics. In the beginning we shall be committed to classical mechanics. But in classical inertial systems an upper bound for velocities would be an intolerable stumbling-block.

So I conceived the idea of introducing the velocity of light by a less physical and more behaviouristic procedure. Light waves, and more generally, electromagnetic waves are a means of communication. Fortunately communicating is a well-known phenomenon in our program. Start and finish are marked by the word Inq and Ani respectively. In the last paragraph Ha has asked the question what the intermediate link is. The answer is that there is something propagating from the speaker to the person who becomes aware of the spoken communication. Three cases have been dealt with by Hb, the first two cases as a background for the third, the only one that matters. The first case is that of mere convection; this is a rather slow method. Wavepropagation is much quicker. There is an undulating medium, characterized by sonic velocities. Finally there is wave propagation through an non-corporal medium (∉ Res) which is characterized by a much higher velocity (that of light). If we should state that there is no quicker means of communication, we should give one more piece of information. This information would not be fundamental, and it would be harmless, because at this stage of our program mechanics is still in statu nascendi, and because that statement would not exclude any quicker motion, as long as it cannot serve as a means of communication.1)

We have quoted the numerical value of the velocity of light with the greatest precision now available.²) Of course it is not unlikely that the mere appearance of a number that is to mean the numerical value of some velocity and that bears a high degree of precision, would make the receiver assume that this number is to mean the numerical value of the velocity of light. Then our behaviouristic introduction would be redundant. But it is better not to rely on this possibility. Moreover it has been our policy to give multiple tautologous information.

As soon as the receiver has recognized that the numerical constant c means the velocity of light, he may use his knowledge of our time-unit in order to find out our length-unit. The precision of the result will mainly depend on five components: 1° and 2° the precision of time-measurements at the transmitter and the receiver, 3° and 4° the precision of measurements of the velocity of light at the transmitter and the receiver, 5° the precision of the receiver's information about the relative velocity of transmitter and receiver.

As to the last source of error we know that some knowledge about this relative velocity is needed for the elimination of the Ole Römer effect (see 2 01 2). If the program lasts long enough this elimination can be done on the strength of the knowledge of the mean relative velocity of transmitter and receiver. Inside our solar system this mean velocity is zero; if the receiver belongs to another solar system it is equal to the relative velocity of the two central bodies. A receiver who understands that the program is sent from our solar system can eliminate the unknown Ole Römer effect by means of an observed Doppler effect. The precision of this determination depends on that of wave-length measurements at the receiver's end.

The receiver may reduce the error from the first and second source by observing the program clock over a long period. In any case a lower bound for the possible relative error in the receiver's

¹⁾ From our text the receiver may get the impression that human bodies can generate electromagnetic communication waves. This misunderstanding would be fairly harmless. Later on we can remove it.

²⁾ This is a sore point. The most trustworthy values of fundamental physical constants quoted by the most trustworthy authors (from unquoted sources) at about the same time differ significantly from each other; even the values given in one and the same work at different places can differ.

knowledge about our length-unit will be given by the relative error of that value for the velocity of light we have inserted in our program. In fact the velocity of light seems to be known with a precision of three in one million, and it is with this precision that we can present our length-unit.

But afterwards this precision may be increased (and we could have done so even if it had been much worse) by presenting Rydberg's constant for hydrogen, which is known with a precision of one in ten millions. As we cannot yet convey in a more definite way what hydrogen is, it is possible that the receiver may interpret "hydrogen" in our context as "ionized helium", or as the name of some other hydrogen-like element. This can be prevented by a list of the Rydberg constants of all hydrogen-like elements or rather by a list of the Rydberg constants for the light elements and their limit for infinite nuclear mass. In these texts we have confined ourselves to hydrogen (see 4 14 6).

4 14 5. Continuation of 4 14 3:

$Ha \operatorname{Inq} Hb$ Pot: $t \operatorname{Pst} \cdot \operatorname{Nnc} Z \operatorname{Cpu} : ? = \operatorname{Cmt} #$

4 14 6. Further continuation of 4 14 3:

This is a description of the hydrogen spectrum. 'Ato H' means the genus "hydrogen atom".

4 18 1. The following text is a continuation of that of 4 11 5.

```
* Ha \operatorname{Inq} Hb ' \operatorname{Nnc} Ha \operatorname{Vul}:

\operatorname{Nnc} \cdot \operatorname{Apu} Hb . \operatorname{Usd} \cdot \operatorname{Apu} Ha \cdot Hb \operatorname{Fer} Rh |

Hb \operatorname{Inq} Ha \cdot - \operatorname{Pot} |

Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot \operatorname{Vol} Rh . > \operatorname{Nim} |
```

```
Hb \operatorname{Inq} Ha : \operatorname{Fal} : \operatorname{Vol} Rh .= \operatorname{Cea} .\operatorname{Cmt}^3 11 \times 10^{1010} .
\operatorname{Tan} \cdot \operatorname{Mas} Rh .= \operatorname{Cea} .\operatorname{Gra} 11011 \times 10^{1011} :
\operatorname{Mas} Rh .> \operatorname{Nim}
\operatorname{Tan} \cdot \wedge D : D \operatorname{Uni} .10^6 \operatorname{Hom} \cdot \to \cdot \operatorname{Pot} .t D \operatorname{Fer} Rh :
\operatorname{Dns} Rh .= \operatorname{Cea} \cdot \operatorname{Gra} 11011 \times 10^{1011} ./.\operatorname{Cmt}^3 11 \times 10^{1010} .
= \operatorname{Cea} .\operatorname{Gra}/\operatorname{Cmt}^3 10010,1011 !
Ha \operatorname{Inq} Hb : \operatorname{Utr} : \operatorname{Dif} \cdot \operatorname{Nne} Hb \operatorname{Cau} : \operatorname{Nne} \cdot \operatorname{Apu} Hb .\operatorname{Usd} .
\operatorname{Apu} Ha \cdot Ri \operatorname{Mov} := : \operatorname{Dif} \cdot \operatorname{Nne} .\operatorname{Etc} .Rj \operatorname{Mov} !
Hb \operatorname{Inq} Ha : \operatorname{Qqm} \cdot \operatorname{Vol} Ri .= .\operatorname{Vol} Rj :
\operatorname{Tan} : \operatorname{Dif} .\operatorname{Nne} Hb \operatorname{Iae} Ri \cdot > \cdot \operatorname{Dif} .\operatorname{Nne} Hb \operatorname{Iae} Rj :
\operatorname{Alpha} : \operatorname{Cea} .\operatorname{Cmt}^3 10010, \operatorname{Cea} : \operatorname{Cmt}^3 11 \times 10^{1010} .
\operatorname{Cea} .\operatorname{Gra} : \operatorname{Cmt}^3 10010, \operatorname{Cea} : \operatorname{Cmt}^3 11 \times 10^{1010} .
\operatorname{Cea} .\operatorname{Gra} : \operatorname{Cmt}^3 10 \times 10^{1011} .
\operatorname{Cea} .\operatorname{Cmt}^3 10 \times 10^{1011} .
\operatorname{Cea} .\operatorname{Cmt}^3 11 \times 10^{1011} .
\operatorname{Cea} .\operatorname{Cmt}^3 11 \times 10^{1010} .
\operatorname{Cea} .\operatorname{Cea} .\operatorname{Cmt}^3 11 \times 10^{1010} .
\operatorname{Cea} .\operatorname{Cmt}^3 11 \times 10^{1010} .
\operatorname{Cea}
```

The words written Mas, Gra, Dns mean "mass", "gramme", "density" respectively. 'Mas' has been introduced by a behaviouristic definition. A crude estimation of the mass-unit would be possible on the strength of the data of the last talk.

'Gra' behaves syntactically like 'Sec', 'Cmt' and so on.

```
4 18 2. * yc = \operatorname{Gra} x . \rightarrow \cdot x \in \operatorname{Rea} . \land . x \ge 0:

\operatorname{Gra} \in {}^{\uparrow}x \cdot x \in \operatorname{Rea} . \land . x \ge 0 : \frown . {}^{\uparrow}\operatorname{Mas}:

\operatorname{Gra} . a + b \cdot = .\operatorname{Gra} a + \operatorname{Gra} b:

\operatorname{Gra} . a - b \cdot = .\operatorname{Gra} a - \operatorname{Gra} b:

\alpha \ge 0 . \rightarrow .\operatorname{Gra} \alpha a = \alpha \operatorname{Gra} a:

\operatorname{Gra} a . \le .\operatorname{Gra} b \cdot \leftrightarrow . a \le b:

\operatorname{Gra} a . = .\operatorname{Gra} b \cdot \leftrightarrow . a = b:

\operatorname{Gra} a . = \operatorname{Cca} .\operatorname{Gra} b \cdot \leftrightarrow \cdot a . = \operatorname{Cca} .b:

\operatorname{Gra} a / \operatorname{Cmt}^3 b . = \cdot \operatorname{Gra} / \operatorname{Cmt}^3 . a / b:

\operatorname{Gra} a \times \operatorname{Cmt} / \operatorname{Sec} . b \cdot = \cdot \operatorname{Gra} \operatorname{Cmt} / \operatorname{Sec} . a \times b:

\operatorname{Etc} *
```

```
4 18 3. * Vol x \neq \text{Cmt}^3 0 . \rightarrow · Dns x . = . Mas x/\text{Vol }x : x \in . Hom \cup Bes · \rightarrow · Dns x . = Cca . Gra/Cmt^3 1 *
```

```
4 18 4. * z \operatorname{Uni} xy . \wedge \cdot \operatorname{Mas} x . = .\operatorname{Gra} \alpha \cdot \wedge \cdot \operatorname{Mas} y . = .\operatorname{Gra} \beta : \rightarrow :
\operatorname{Mas} z = \cdot \operatorname{Mas} x . + .\operatorname{Mas} y \cdot = \cdot \operatorname{Gra} \alpha . + .\operatorname{Gra} \beta = \cdot \operatorname{Gra} (\alpha + \beta)
z \operatorname{Uni} xy . \wedge \cdot \operatorname{Mas} x . = .\operatorname{Gra} \alpha \cdot \wedge \cdot \operatorname{Mas} z . = .\operatorname{Gra} \gamma : \rightarrow :
\operatorname{Mas} y = \cdot \operatorname{Mas} z . - .\operatorname{Mas} x \cdot = \cdot \operatorname{Gra} \gamma . - .\operatorname{Gra} \alpha \cdot = .\operatorname{Gra} (\gamma - \alpha) *
```

```
4 18 5. \# x \in \operatorname{Res} \cup \operatorname{Hom} \cup \operatorname{Bes} \longrightarrow \operatorname{ht} t \in \operatorname{Tem} \wedge \operatorname{ht} \operatorname{Ext} \longrightarrow \operatorname{ho} \times \operatorname{as} > 0 \cdot \operatorname{ho} t \operatorname{Mas} x = \operatorname{Gra} x \#
```

```
4 18 6. \# x \in \text{Res} \cup \text{Hom} \cup \text{Bes} \rightarrow \land t \ t \in \text{Tem} \land tx \text{ Ext}
                                                    \rightarrow: \forall y \cdot y = .t \operatorname{Cor} x:
                                          t \operatorname{Cor} x = t \operatorname{Cor} y \longleftrightarrow x = y
                                          t \operatorname{Loc} x = \operatorname{Loc} t \operatorname{Cor} x:
                                          t \operatorname{Vol} x = \operatorname{Vol} t \operatorname{Cor} x
                                          t \operatorname{Lon} x = \operatorname{Lon} t \operatorname{Cor} x
                                          t \operatorname{Mas} x = \cdot \operatorname{Mas} t \operatorname{Cor} x:
                                          t \operatorname{Dns} x = \operatorname{Dns} t \operatorname{Cor} x 
                                           The word written Cor (fL \text{ corpus} = \text{body}) means "body".
                                  Though it has been introduced in a mere formal way, the receiver
                                  may guess that the bodies of men, animals and things may depend
                                  on time. In 4 07 1 the volume of humans as a function of time has
                                  been communicated. An air wave may be considered as a thing
                                  (\in Res) with a varying body:
4 18 7. A comment on 4 14 3, especially on the part from t<sub>3</sub> to t<sub>4</sub>:
                                  ^* Hc\operatorname{Inq} Hd : \operatorname{Qqm} : \operatorname{PAN} Hb\operatorname{Inq} Ha * \operatorname{Cel}\operatorname{Med} . w\operatorname{Mov} *
                                                    = 'Cca.Cmt/Sec 10<sup>1111</sup>: \rightarrow w \notin \text{Res.} \land w \neq \text{Aer.} \land -w \text{ Par Aer}
                                                    Tan: Nnc Hc Inq Hd: w \in \text{Res}.
                                                             \wedge \wedge t : tw \to t \text{Cor } w \cdot \text{Par Aer} .
                                                   Qqm t_3 t_4 Hb \operatorname{Inq} Ha \cdot w = \cdot \forall t \operatorname{Etc} \cdot
                                                   Tan : Nne Hc \operatorname{Inq} Hd: \operatorname{Cor} w = \operatorname{Cor} Y v v \operatorname{Uni} z z \in M.
                                                              \wedge tz \text{ Mov}
                                                   \forall A : A \in \text{Res.} \land \text{Pot} : \forall \lceil t_1, t_2 \rceil \text{ Put} \lceil t_1, t_2 \rceil
                                                   t_1 \operatorname{Usq} t_2 \cdot \operatorname{Apu} X \cdot \operatorname{Usd} \cdot \operatorname{Apu} Y \cdot A \operatorname{Mov}:
                                                   \wedge \cdot \text{Dst } XY/(t_2-t_1) = \text{Cea. Cmt/Sec } 10^{1111}
                                                   Sed \land \land \vdash t \cdot t' \dashv : t_1 \leq t < t' \leq t_1 \cdot \rightarrow \cdot t \operatorname{Cor} A \cdot \neq \cdot t' \operatorname{Cor} A = t' \operatorname{Co
4 18 8.
                                         Ha \operatorname{Inq} Ha
                                          x \in \text{Hom} \rightarrow \text{Ini} x \text{Ext} - \text{Ini} \text{Cor} x \text{Ext} =
                                                   Cca . Sec 11 \times 10^{10111}:
                                          \forall x : x \in \text{Bes.} \land : \text{Ini.} x \text{ Ext.} - : \text{Ini.} \text{ Cor } x \text{ . Ext.} > \text{Sec } 0 :
                                          x \in \operatorname{Hom} A \to V \setminus Y \cdot z \cap Y \cdot z \in \operatorname{Hom} A \cdot y = A \cdot \operatorname{Mat} x \cdot A \cdot z = A \cdot \operatorname{Pat} x
                                          \forall x : x \in \text{Bes.} \land ` \lor [y : z] : y \land z \in \text{Bes.} \land ` y = . \text{Mat } x : \land ` z = . \text{Pat } x :
                                          x \in \text{Hom} . \rightarrow : \land t : \text{Ini} \cdot \text{Cor} \ x \cdot \text{Ext} : \text{Ant} : t : \text{Ant} : \text{Ini} \cdot x \cdot \text{Ext}
                                                    \rightarrow: t \operatorname{Cor} x \cdot \operatorname{Par} \cdot t \operatorname{Cor} \cdot \operatorname{Mat} x:
                                           \forall x : x \in \text{Bes.} \land \land \land t \text{. Etc.}
                                          x \in \text{Hom} \cdot \wedge : s = \text{Ini} \cdot \text{Cor} x \cdot \text{Ext}
                                                    \rightarrow : \vee \lceil u \cdot v \rceil s \operatorname{Cor} u \cdot \operatorname{Par} \cdot s \operatorname{Cor} \cdot \operatorname{Mat} x :
                                                              Λ' Pau Ant. s · Cor v : Par : Pau Ant. s · Cor . Pat x ·
```

 $\wedge : s \operatorname{Cor} x \cdot \operatorname{Uni} \cdot s \operatorname{Cor} u \cdot s \operatorname{Cor} v :$

```
\forall x : x \in \text{Bes.} \land . \text{Etc.}
\text{Hom} = \text{Hom Fem.} \cup . \text{Hom Msc.}
\text{Hom Fem } \cap \text{Hom Msc.} = \lceil \rceil
\text{Car.} \land x \cdot \text{Nnc.} x \text{ Ext.} \land . x \in \text{Hom Fem.}
\text{Pau>} \land \text{Car.} \land x \cdot \text{Nnc.} x \text{ Ext.} \land . x \in \text{Hom Msc.}
y = \text{Mat.} x \cdot \land . y \in \text{Hom.} \rightarrow . y \in \text{Hom Fem.} \land :
y = \text{Pat.} x \cdot \land . y \in \text{Hom.} \rightarrow . y \in \text{Hom Msc.}
```

 $x \in \text{Hom} \cup \text{Bes} : \rightarrow : \text{Fin. Cor } x \cdot \text{Pst. Fin } x \neq$

The words written Mat (fL mater=mother), Pat (fL pater=father), Fem (fL femininus=female), Msc (fL masculinus=male) mean "mother", "father", "female", "male" respectively. The last paragraph is somewhat premature. Its notions will not be used in this volume.

4 19 1.
$$* \forall a \land \land t \cdot t' \exists t_1 \leq t \leq t' \leq t_2 \rightarrow t \operatorname{Usq} t' \operatorname{Cel} \operatorname{Med} x = a :$$

 $\rightarrow : \land \land t \cdot t' \exists t_1 \leq t \leq t' \leq t_2 \to : \operatorname{Cel} x \cdot t' = \operatorname{Cel} x \cdot t' *$

The word written Cel (fL celeritas = velocity) means "velocity". Here it has been introduced for uniform motions only.

4 19 2.
$$* \land \ulcorner t . t' \urcorner : t_1 < t \le t' < t_2.$$

$$\rightarrow ` \operatorname{Cel} x . t := \cdot \operatorname{Cel} x . t' : \land : \operatorname{Cel} y . t := \cdot \operatorname{Cel} y . t' :$$

$$\rightarrow : \land t : t_1 < t < t_2. \rightarrow : \operatorname{Cel} x . t := \cdot \operatorname{Cel} y . t' :$$

$$\leftrightarrow ` \land t : t_1 \le t \le t' \le t_2. \rightarrow : t \operatorname{Dst} xy . = . t' \operatorname{Dst} xy .$$

- 4 19 3. * $\tau \in \text{Rea.} \land . \tau \neq 0 . \land . a \in \text{Vco}$ $\rightarrow : \text{Cmt/Sec.} a/\tau := .\text{Cmt } a/\text{Sec } \tau . \in : \text{Cmt/Sec.} \text{Vco}$: $\tau \in \text{Rea.} \land . b \in \text{Vco} : \rightarrow . \text{Sec } \tau \times \text{Cmt/Sec } b . = \text{Cmt } \tau b$: $\mu \in \text{Pos.} \land . b \in \text{Vco} : \rightarrow . \text{Gra } \mu \times \text{Cmt/Sec } b . = \text{Gra Cmt/Sec } \mu b$: Etc *
- 4 19 4. $*v \in \operatorname{Cmt/Sec} . \operatorname{Vco} : \leftrightarrow ` \lor w : w \in \operatorname{Vco} . \land `v = .\operatorname{Cmt/Sec} w .$ $v \in \operatorname{Gra} \operatorname{Cmt/Sec} . \operatorname{Vco} : \leftrightarrow ` \lor w : w \in \operatorname{Vco} . \land `v = .\operatorname{Gra} \operatorname{Cmt/Sec} w .$ Etc *
- 4 19 5. * $a \sim b \in \text{Veo} . \land . \gamma \in \text{Rea} \cdot \rightarrow \text{`Cmt/Sec } a \cdot = \cdot \text{Cmt/Sec } b \cdot \leftrightarrow$ $.a = b : \land : \text{Cmt/Sec } a + \text{Cmt/Sec } b \cdot = \cdot \text{Cmt/Sec } . a + b : \land$ $\cdot \gamma \text{ Cmt/Sec } a \cdot = \cdot \text{Cmt/Sec } \gamma a \cdot \land$ $: \text{Cmt/Sec } a \cdot = \cdot \text{Cmt/Sec } b \cdot \leftrightarrow \cdot a \cdot = \text{Cca } . b : \land$ $: \text{Cmt/Sec } a \cdot = \text{Cca } . \text{Cmt/Sec } b \cdot \leftrightarrow \cdot a \cdot = \text{Cca } . b : \land$ $\cdot \text{Etc} *$
- 4 19 6. $*f \in \text{Coo} \rightarrow `a \in \text{Veo} \rightarrow :f \cdot \text{Cmt/Sec} \ a \cdot = \cdot f \cdot \text{Gra} \ \text{Gra} \ a \cdot = \cdot f \cdot \text{Gra} \ a \cdot = \cdot f \cdot$

```
4 19 7. * \wedge t t_1 < t < t_2 \rightarrow: \operatorname{Cel} x \cdot t = v \rightarrow v \in \operatorname{Cmt/Sec} \cdot \operatorname{Veo}:
                               \rightarrow v = (t' \operatorname{Loc} x - t \operatorname{Loc} x)/(t-t')
                                      \wedge \cdot t \operatorname{Loc} x = \cdot t_1 \operatorname{Loc} x + (t - t_1) v *
4 19 8. * \wedge \lceil tt' \rceil \cdot t_1 \leq t < t' \leq t_2. \rightarrow : t \operatorname{Usq} t'. Cel Med x := \operatorname{Cmt} \gamma:
                               \rightarrow ' \land t: t_1 < t < t_2 \rightarrow ' |\operatorname{Cel} x \cdot t| = \operatorname{Cmt} \gamma #
4 20 1. In the following texts p is a substitute for the noise of a collision.
                    * t1 p t2 :
                         Ha \operatorname{Inq} Hb \cdot ?z \cdot t_1 t_2 \operatorname{Fit} z
                         Hb \operatorname{Inq} Ha \cdot \vee \lceil x \cdot y \rceil \cdot \operatorname{Put} \lceil x \cdot y \rceil \cdot x \in \operatorname{Res} \cdot \wedge \cdot y \in \operatorname{Res} \cdot \wedge \cdot t_1 xy \operatorname{Ce} u:
                         Ha \operatorname{Inq} Hb. ? = \operatorname{Ccu}:
                         Hb \operatorname{Inq} Ha : \operatorname{Cea} t_1 : y \operatorname{Mov} \cdot \wedge \cdot - x \operatorname{Mov} \cdot \wedge \cdot t_1 y \operatorname{Apu} x:
                               t < t_1 < t'. \rightarrow: Cel y \cdot t' \cdot = \cdot - Cel y \cdot t \cdot \neq \cdot Cmt/Sec <math>0_{Vco}: \land
                                t_1 x \operatorname{Cau} : \operatorname{Cel} y \cdot t' \neq \operatorname{Cel} y \cdot t'
                                      \wedge t_1 x \operatorname{Mut} : \operatorname{Cel} y \cdot \operatorname{Cea} t_1 : \operatorname{Mod} \cdot - a \operatorname{Ilo} a :
                         Ha Inq Hb. Cur Pot:
                         Hb \operatorname{Inq} Ha : \operatorname{Qia} \cdot \operatorname{Mas} x \cdot \operatorname{Mul} > \cdot \operatorname{Mas} y *
                           The word written Ccu (fL concutiunt = collide) means "collide".
4 20 2. # t<sub>1</sub>p l
                         Ha \operatorname{Inq} Hb \operatorname{Utr} : \vee \operatorname{r} x \cdot y \operatorname{r} \cdot x \curvearrowright y \in \operatorname{Res} \cdot \wedge \cdot t_1 xy \operatorname{Ccu} I
                         Hb \operatorname{Inq} Ha \cdot \operatorname{Ver} \cdot t_1 x \operatorname{Apu} y 
                         Ha \operatorname{Inq} Hb 'Utr · Mas x . Mul > . Mas y |
                         Hb \operatorname{Ing} Ha \cdot \operatorname{Fal} \operatorname{Mas} x = \operatorname{Mas} y \cdot \operatorname{Erg} \cdot t < t_1 < t' \cdot \rightarrow
                               \operatorname{Cel} x \cdot t' = \operatorname{Cel} y \cdot t : \land : \operatorname{Cel} y \cdot t' : = \operatorname{Cel} x \cdot t : \land
                               t_1 x \operatorname{Cau} : \operatorname{Cel} y \cdot t' \neq \operatorname{Cel} y \cdot t \wedge
                                t_1 y \operatorname{Cau} : \operatorname{Cel} x \cdot t' \cdot \neq \cdot \operatorname{Cel} x \cdot t' \wedge : t_1 xy \operatorname{Ceu} \cdot \operatorname{Cau} \cdot \operatorname{Etc} *
4 20 3. # t<sub>1</sub> p l
                         Ha \operatorname{Inq} Hb \operatorname{Utr} \vee \lceil x \cdot y \rceil \operatorname{Put} \lceil x \cdot y \rceil \cdot x \rightsquigarrow y \in \operatorname{Res} \wedge \operatorname{tr} xy \operatorname{Ceu} \mathsf{I}
                          Hb \operatorname{Inq} Ha \operatorname{Ver} I
                         Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot \operatorname{Mas} x := \operatorname{Mas} y
                         Hb \operatorname{Inq} Ha \longrightarrow \operatorname{Nes} : \operatorname{Sed}^{t_1} : t < t_1 < t' \longrightarrow :
                               \operatorname{Cel} x \cdot t \cdot = a : \wedge : \operatorname{Cel} x \cdot t' = a' : \wedge : \operatorname{Cel} y \cdot t = b : \wedge : \operatorname{Cel} y \cdot t' = b' : \wedge
                               \operatorname{Mas} x = \mu \cdot \wedge \cdot \operatorname{Mas} y = \nu \rightarrow \cdot \mu a + \nu b = \mu a' + \nu b' \cdot \wedge \cdot
                                     a' \neq a. A. \mu a Ia + \nu b Ib = \mu a' Ia' + \nu b' Ib'
                                \operatorname{Erg} \cdot \mu a - \mu a' = \nu b' - \nu b \cdot \wedge
                                     \mu a \mathbf{I} a - \mu a' \mathbf{I} a' = \nu b' \mathbf{I} b' - \nu b \mathbf{I} b
                               ^{t_{3}}Pon: \bigvee \lceil \alpha . \beta . e \rceil \cdot Put e \cdot \alpha \land \beta \in Rea. \land .
                                      e \in \text{Cmt/Sec Vco} \land \land a = \alpha e \land \land b = \beta e
                                Erg: \forall \lceil \alpha' \cdot \beta' \rceil \cdot \alpha' = \alpha' e \cdot \land \cdot b' = \beta' e
```

4 20 4. $# Ha \operatorname{Inq} Ha$

```
\begin{aligned} &\operatorname{Erg^{t_*}} \cdot a + a' = b + b' \cdot \wedge \cdot \\ &a' = [\mu a + \nu b + \nu (b - a)] / (\mu + \nu) \cdot \wedge \cdot \\ &b' = [\mu a + \nu b + \mu (a - b)] / (\mu + \nu)^{t_*} \vdots \\ &\operatorname{Nnc} Hb \operatorname{Mut} \overset{\checkmark}{\gamma} p \cdot t_2 t_5 Hb \operatorname{Inq} Ha p \\ &\overset{\cdot}{\operatorname{Mod}} : \operatorname{Sin} \overset{\checkmark}{\gamma} q \cdot t_3 t_4 Hb \operatorname{Inq} Ha q \vdots \\ &\to : \overset{\checkmark}{r} \cdot \operatorname{Nnc} Hb \operatorname{Inq} Ha r \cdot \in \operatorname{Ver} \overset{\checkmark}{\sim} \in \operatorname{Lex} \operatorname{Nat} \vdots \\ &t'' Hb \operatorname{Sci} : \overset{?}{\circ} = \operatorname{Cel} x \cdot t \cdot \overset{\checkmark}{\sim} \cdot \operatorname{Cel} y \cdot t \cdot \\ &\to : \operatorname{Pot} \overset{\star}{t}'' Hb \operatorname{Cpu} : \overset{?}{\circ} = \operatorname{Cel} x \cdot t' \cdot \overset{\checkmark}{\sim} \cdot \operatorname{Cel} y \cdot t' \cdot \\ &Ha \operatorname{Inq} Hb \overset{\cdot}{\tau} x \overset{\checkmark}{\tau} y \in \operatorname{Res} \cdot \to : \operatorname{Nnc} Ha \operatorname{Sci} \cdot \overset{?}{\tau} = \overset{\checkmark}{\tau} \varrho \cdot \\ &\operatorname{Mas} y = \varrho \operatorname{Mas} x \overset{\#}{\tau} \end{aligned}
```

Hb has formulated the law of elastic collision. Ha states that this law implies an unambiguous definition of mass ratio. The receiver can conclude that 'Mas' means something that is proportional with mass. The mass unit, however, remains unknown.

Lex Nat (fL lex naturalis=law of nature) means "law of nature".

It might be remarked that an ideal elastic collision cannot produce any noise. This is true. Nevertheless I think if there were some danger in showing the noise of collision, it is more than counterbalanced by the obvious gain of demonstrativeness.

```
Ctr x \in \text{Cvx.Loc} x!

Mas x = \mu \land A. Mas y = v \land Az Uni xy \Leftrightarrow Ax = \mu \land Ax. Mas y = v \land Az Uni xy \Leftrightarrow Ax = \mu \land Ax. (\mu + v) Ctr x = \mu \land Ax. (\mu + v) Ctr x = \mu \land Ax. (\mu + v) Cel Ctr x + v \land Ax.

Ipu x = Ax Ax x \land Ax Cel Ctr x = Ax.

Ipu x \in Ax Cra Cmt/Sec. Vco:

x \land Ax = Ax Cra Cmt/Sec. Vco:

x \land Ax
```

 $t < t_1 < t' \land t_1 xy \text{ Ccu} \rightarrow z \text{ Uni } xy$

 \rightarrow : Kin z.t' = 'Kin z.t' #

The word written Ctr (fL centrum=centre) means "mass centre". The words written Ipu and Kin mean impulse and kinetic energy respectively.

It is understood that in 4 20 1-4 20 3 we dealt with mass points, whereas in 4 20 4 the objects may be bodies with positive extensions. This could be stressed by means of a comment.

4 21 1. # Via $x \cdot t = .t$ Loe Ctr x #

The word written Via (fL via = path) means the path function of some body.

4 21 2.
$$* \operatorname{Cel} x \cdot t \cdot = : \frac{d}{dt} \cdot \operatorname{Via} x \cdot t :$$

$$= : \lim^{*} h \to 0 \cdot \operatorname{Via} x \cdot t + h \cdot - \operatorname{Via} x \cdot t : /h :$$

$$\operatorname{Cel} x \cdot t \cdot \in \cdot \operatorname{Cmt}/\operatorname{Sec} \cdot \operatorname{Vco} *$$

4 21 3. *
$$\operatorname{Acc} x \cdot t \cdot = : \frac{d}{dt} \cdot \operatorname{Cel} x \cdot t :$$

$$= : \lim^{\cdot} h \to 0 \cdot \operatorname{Cel} x \cdot t + h \cdot - \operatorname{Cel} x \cdot t : /h :$$

$$\operatorname{Acc} x \cdot t \cdot \in \cdot \operatorname{Cmt}/\operatorname{Sec}^{2} \cdot \operatorname{Vco} *$$

The word written Acc (fL acceleratio = acceleration) means "acceleration".

The use of the traditional notation of calculus is rather provisional (see 1 35 4).

4 22 1. $\# \div \in \text{Res}$:

Loc † . ∈ Cca . Bul:

 $\operatorname{Loc} := \operatorname{Cca} : \operatorname{Bul} a \varrho : \to a := \operatorname{Cca} : \operatorname{Ctr} : \circ \circ \varrho = \operatorname{Rad} :$

= Cca \cdot Cmt 101×10^{11011} :

 $Dns \pm . = Cca . Gra/Cmt^3 101,1 *$

The word written (the astronomical symbol for "earth") means "earth".

- 4 22 2. * Loc $\, \pm \, \text{Loc} \, \pm \, \text{Cea. Bul} \, a \varrho \, \cdot \, \rightarrow \, \cdot \, \text{Loc} \, \text{Sfi} \, \pm \, \cdot \, = \, \text{Cea. Sph} \, a \varrho \, \, *$ The word written Sfi (fL superficies = surface) means "surface".
- 4 22 3. * $x \in Hom \cup Bes \rightarrow Dst.x.Sfi + < Vld *$
- 4 22 4. $\mathfrak p$ is again a substitute for a noise of some collision.

 $Ha \operatorname{Inq} Hb : ?x \cdot \operatorname{Cea} t_1 \cdot \operatorname{Fit} x$

 $Hb \operatorname{Inq} Ha \lor y \cdot \operatorname{Put} y \cdot y \in \operatorname{Res.} \land \cdot t_1 y \circ \operatorname{Ccu}$:

 $t. Pau Ant.t_1 \rightarrow ty Mov. \land ty Cad$

 $Ha\operatorname{Inq} Hb$. $? = \operatorname{Cad} \mathsf{I}$

 $Hb \operatorname{Inq} Ha: ty \operatorname{Cad} . \longleftrightarrow t . \operatorname{Adv} \operatorname{Sfi} _{\bigcirc} . y \operatorname{Mov} I$

 $Ha\operatorname{Inq} Hb: ?z \cdot \operatorname{Ceat}_{1} \cdot z\operatorname{Cau} \cdot \operatorname{Nnc} y\operatorname{Cad} I$

Ha Inq Hb Cur I

 $Hb\operatorname{Inq} Ha: \neg \cdot \operatorname{Dns} y \cdot < \operatorname{Vld} : \wedge \cdot t \cdot \operatorname{Dst} \cdot y \cdot \operatorname{Sfi} \dot{\neg} :$

 $Mul < Rad + \rightarrow Nes ty Cad I$

Ha Inq Hb Cur I

 $Hb \operatorname{Inq} Ha: {}^{\vee}p \cdot \operatorname{PAN} Ha \operatorname{Inq} Hb p \cdot \in \operatorname{Lex} \operatorname{Nat} I$

 $Ha\operatorname{Inq} Hb$. $? = \operatorname{Cel} y$

```
Hb \operatorname{Inq} Ha \cdot \operatorname{Pon} : \operatorname{Dst} \cdot y \cdot \operatorname{Sfi} \cdot \operatorname{Mul} < \cdot \operatorname{Rad} \cdot :
            \wedge \cdot t_0 \leq t \leq t_1 \cdot \wedge \cdot ty \operatorname{Cad} \cdot
            \wedge : t_0 \operatorname{Loc} \operatorname{Ctr} : - : t_0 \operatorname{Loc} y := t_0 : \operatorname{Dst} : \operatorname{Ctr} : y : \times e
      Erg : \operatorname{Cel} y \cdot t \cdot - : \operatorname{Cel} y \cdot t_0 = g(t - t_0) e : \land :
            Via y \cdot t = Via y \cdot t_0 + (t - t_0) \cdot Cel y \cdot t_0 + 0, 1 \times g(t - t_0)^{10} e
            \wedge \cdot t \operatorname{Acc} y = ge
            \forall p . \text{PAN } Hb \text{ Inq } Ha \cdot \in \text{Lex Nat I}
Ha \operatorname{Inq} Hb \cdot ? = g 
Hb \operatorname{Inq} Ha \cdot g := \operatorname{Cea} \cdot \operatorname{Cmt}/\operatorname{Sec}^2 1111010101 
Ha Inq Hb: Cur 'PAN Hb Inq Ha: Pon · Etc. ty Cad. Etc |
Hb \operatorname{Inq} Ha = \operatorname{Nes} : \operatorname{Erg} : \operatorname{Nne} Hb \operatorname{Mut}^{\rightarrow} p \cdot \operatorname{PAN} Hb \operatorname{Inq} Ha p
      Mod: Pot. ty Cad · Ilo. ty Cad |
Ha \operatorname{Inq} Hb \operatorname{Utr} : t \uparrow \operatorname{Cau} \cdot t \operatorname{Acc} y = ge \operatorname{I}
Hb \operatorname{Inq} Ha \operatorname{Ver} I
Ha \operatorname{Inq} Hb : \operatorname{Utr} \cdot \operatorname{Mod} \cdot \operatorname{Fer} \sim \operatorname{Iac} \mathbf{I}
Hb \operatorname{Inq} Ha \longrightarrow \operatorname{Mod} \operatorname{Fer} A \longrightarrow \operatorname{Mod} \operatorname{Iac} 
      Pon tz Ext. \land .z \in \text{Res} \leftrightarrow .z = x \sim y : \land .x \neq y.
            \wedge Via x \cdot t - \text{Via } y \cdot t \cdot = \cdot t \text{ Dst } xy \times e^{-t}
     \operatorname{Erg} tx \operatorname{Cau} : \operatorname{Acc} y \cdot t = x \operatorname{Acc} y \cdot t \wedge
           ty \operatorname{Cau} : \operatorname{Acc} x \cdot t := y \operatorname{Acc} x \cdot t :
                  \wedge: \forall \gamma Put \gamma
                 x \operatorname{Acc} y \cdot t = \gamma \operatorname{Mas} x/(\operatorname{Dst} xy)^{10} \cdot \times e : \wedge :
                 y \operatorname{Acc} x \cdot t = -\gamma \operatorname{Mas} y/(\operatorname{Dst} xy)^{10} \times e \cdot \wedge :
                 \gamma = \frac{\text{Cmt}^3}{\text{Sec}^2 \text{Gra.} 100011110,01} \times 10^{-100000}: A
                  Err._{\nu} = Etc._{c} = Cca.Cmt^{3}/Sec^{2}Gra 10^{-100010}
Ha \operatorname{Inq} Hb \operatorname{Cur} I
Hb \operatorname{Ing} Ha : {}^{\vee}p \cdot \operatorname{PAN} Hb \operatorname{Ing} Ha \cdot \in \operatorname{Lex} \operatorname{Nat} I
Ha \operatorname{Inq} Hb \cdot \operatorname{Nnc} Ha \operatorname{Sei} ? = \operatorname{Gra} *
```

The law of free fall and the universal attraction law. Cad (fL cadit = falls) means "falls". As stated by Ha, the receiver can now express our mass unit by his own unit, if he compares our value of the universal gravitation constant γ with his. Unfortunately the precision of the value of γ we can communicate is rather poor. His knowledge about our mass-unit will be as poor, but we can increase the precision by indicating the mass of the hydrogen atom. This is still far from satisfactory. For the present we must content ourselves with this result. If we want to give the receiver the opportunity to acquire more precise knowledge about our mass-unit, we ought to introduce chemical substances and a few physical notions such as pressure and temperature.

```
4 22 6. Continuation of the text of 4 22 4.
```

* $Ha \operatorname{Inq} Hb \operatorname{:} \operatorname{Pot} : \operatorname{Car} : ^{x} \cdot tx \operatorname{Ext} . \wedge .x \in \operatorname{Res} : > 10 \operatorname{I}$ $Hb \operatorname{Inq} Ha : \operatorname{Acc} x . t : = : \mathcal{L} \cdot z \in \operatorname{Res} .z \neq x . \wedge .tz \operatorname{Ext} \cdot z \operatorname{Acc} x . t :$ $\wedge z : z \neq x . \rightarrow : \operatorname{Dst} xz . < \operatorname{Vld} : \rightarrow : \operatorname{Acc} x . t : = \operatorname{Cca} .\operatorname{Cmt}/\operatorname{Sec}^{2} \operatorname{O}_{\operatorname{Veo}} : \wedge :$ $\wedge z : z \neq x . \wedge .z \neq y : \rightarrow : \operatorname{Dst} xz . < \operatorname{Vld} : \rightarrow :$ $\operatorname{Acc} x . t : = \operatorname{Cca} : y \operatorname{Acc} x . t :$

4 23 1. *
$$p > 0$$
. $\land . e \ge 0$. $\land . p \curvearrowright e \in \text{Rea} \cdot \rightarrow \vdots$
 $\text{Ssc}_* pe . = : \uparrow x : x \in \text{Rea}^3 . \land \land \land \lceil \xi_1 . \xi_2 . \xi_3 \rceil : x = \lceil \xi_1 . \xi_2 . \xi_3 \rceil .$
 $\rightarrow : \xi_3 = 0$. $\land . \sqrt{\xi_1^{10} + \xi_2^{10}} = p + e \xi_1 .$
 $p > 0$. $\land . e \ge 0 : \rightarrow : z \in . \text{Ssc} pe : \land : y \in . \text{Foc} z :$
 $\leftrightarrow : z \subset \text{Spa} . \land : \forall f : f \in \text{Coo} . \land \land f : z : = . \text{Ssc}_* pe : \land : fy : = \lceil 0 . 0 . 0 \rceil :$
 $z \in \text{Ssc} . \leftrightarrow : \forall \lceil p . e \rceil : z \in . \text{Ssc} pe *$

The word written Ssc p e (fL sectio semiconi) means "ellipses or parabolae or semi-hyperbolae with parameter p and eccentricity e". (The eccentricity is 1 in the case of the parabola; if Ssc is an ellipse with the axes a, b, and a > b, then $e = \sqrt{a^2 - b^2}/a < 1$, and $p = b^2/a$; in the case of the hyperbola with "real" axis a and "imaginary" axis b, we have $e = \sqrt{a^2 + b^2}/a > 1$, and $p = b^2/a$.)

The word written Foc means "focus".

$4\ 23\ 2. \quad \# Ha \operatorname{Inq} Hb$

Pon: $A = : \text{Mas } x . \text{Mul} > . \text{Mas } y \cdot \wedge \cdot \text{Cel } x = \text{Cmt}/\text{Sec } 0_{\text{Vco}}.$ $\wedge : \neg \forall t . t \text{ Dst } xy = \text{Cmt } 0 : \wedge \cdot \text{Acc } y . = .x \text{ Acc } y \cdot \wedge \cdot \neg \forall u : u \in \text{Ret } . \wedge \cdot \wedge \text{Via } y . \in u \cdot \cdot \text{Sec } y \cdot \neg \forall u : u \in \text{Ret } . \wedge \cdot \wedge \text{Via } y . \in \text{Foc } \wedge \text{Via } y \cdot \cdot \wedge \neg \forall u : u \in \text{Acc } y \cdot \neg \forall u : u \in \text{Acc } y \cdot \neg \forall u : u \in \text{Acc } y \cdot \neg \forall u : u \in \text{Ret } . \wedge \cdot \wedge \neg \forall u : u \in \text{Acc } y \cdot \neg \forall u : u \in \text{A$

The word written Rev (fL revolutio) means "period of revolution".

4 23 3. *
$$tx \operatorname{Rot} ad . \leftrightarrow \iota a . = \operatorname{Cea} . \operatorname{Ctr} x . \wedge . d \in \operatorname{Sec}^{-1} \operatorname{Veo} .$$

$$\wedge ! \lor f : f \in \operatorname{Coo}^{+} . \wedge . f a = 0_{3} . \wedge . f d = \lceil 0 . 0 . \operatorname{Sec} \mid d \mid \rceil .$$

$$\wedge ! \land \lceil y . \zeta_{1} . \zeta_{2} . \zeta_{3} \rceil " y \operatorname{Par} x . \wedge \cdot f . t \operatorname{Loe} \operatorname{Ctr} y = \lceil \zeta_{1} . \zeta_{2} . \zeta_{3} \rceil :$$

$$\rightarrow "f \cdot \operatorname{Cel} y . t := \lceil -\operatorname{Sec} \mid d \mid \zeta_{2} . \operatorname{Sec} \mid d \mid \zeta_{1} . 0 \rceil "$$
The word written Rot (fL rotat=rotates) means "rotates".

d is the vector of angular velocity. 4 23 4. # $t_1 \operatorname{Usq} t_2 x \operatorname{Rot} a d \leftrightarrow : \land t \cdot t_1 \leq t \leq t_2 \rightarrow . t x \operatorname{Rot} a d =$

4 23 5. *
$$t_1 \operatorname{Usq} t_2 x \operatorname{Rot} a d . \wedge . f \in \operatorname{Coo}^+ . \wedge . f a = 0_3$$
.
 $\wedge . f d = \lceil 0 . 0 . \operatorname{Sec} \mid d \mid \rceil . \wedge . y \operatorname{Par} x . \wedge \lceil f \cdot \operatorname{Via} y . t_1 := \lceil \zeta_1 . \zeta_2 . \zeta_3 \rceil \vdots$
 $\Rightarrow \lceil f \cdot \operatorname{Via} y . t_2 := \lceil \zeta_1 \cos \mid d \mid (t_2 - t_1) - \zeta_2 \sin \mid d \mid (t_2 - t_1) . \zeta_3 \rceil = \zeta_1 \sin \mid d \mid (t_2 - t_1) + \zeta_2 \cos \mid d \mid (t_2 - t_1) . \zeta_3 \rceil = \zeta_1 \sin \mid d \mid (t_2 - t_1) + \zeta_2 \cos \mid d \mid (t_2 - t_1) . \zeta_3 \rceil = \zeta_1 \sin \mid d \mid (t_2 - t_1) + \zeta_2 \cos \mid d \mid (t_2 - t_1) . \zeta_3 \rceil = \zeta_1 \cos \mid d \mid (t_2 - t_1) . \zeta_3 \mid (t_2 - t_1) . \zeta_3 \mid d \mid (t_2 - t_1) . \zeta_3 \mid (t_$

```
4 23 6. # t_1 \operatorname{Usq} t_2 \cdot x \operatorname{Rot} a d \cdot \wedge \cdot t_2 - t_1 = 10 \pi / |d| \cdot \wedge \cdot y \operatorname{Par} x:

\rightarrow \operatorname{Via} y \cdot t_2 \cdot = \operatorname{Via} y \cdot t_1 : \wedge \cdot \operatorname{Rtt} x = t_2 - t_1 = 10 \pi / |d| \cdot \wedge \cdot y \operatorname{Par} x

The word written Rtt means "period of (uniform) rotation".
```

4 23 7. # $d/|d| = \cdot \operatorname{Axs} \cdot tx \operatorname{Rot} a d : \wedge : a = \cdot \operatorname{Ctr} \cdot tx \operatorname{Rot} a d #$ 'Axs' and 'Ctr' mean axis and centre of the rotation. 1)

4 24 1.
$$\begin{align*}{l} \begin{align*}{l} \b$$

The word written Ste (fL stella=star) means "celestial body". A few near celestial bodies have been mentioned. The usual astronomical names have been used. As i means the asteroids. Components of double and triple star systems are distinguished by the letters A, B, C.

```
4 24 2. * x \in \text{Ste} . \to \text{:} \text{Loc} \, x \in \text{Spa} . \land \cdot \text{Loc} \, x . \text{Cea} \in . \text{Bul} \cdot \land : \text{Loc} \, x . = \text{Cea} . \text{Bul} \, a \varrho \cdot \to \cdot a . = \text{Cea} . \text{Ctr} \, x \cdot \land \cdot \varrho . = \text{Cea} . \text{Rad} \, x \cdot \land \cdot \text{Sfi} \, x . = \text{Cea} . \text{Sph} \, a \, \varrho ^{\#}
```

```
4 24 3. * Mas \odot .= Cca . Gra 11 × 10<sup>1101101</sup>:

Mas \d .= Cca . Gra 101 × 10<sup>1011010</sup>:

Mas \d .= Cca . Gra 10<sup>10101010</sup>:

x \in \text{Aoi} .→ · Mas x .< . Gra 10<sup>1010000</sup>:

Mas Sirius A .= Cca . 10,1 × Mas \odot:

Mas Procyon A .= Cca . 1,1 × Mas \odot:

Mas \alpha Centauri A .= Cca . 0,11 × Mas \odot:

Mas \alpha Centauri B .= Cca . 0,11 × Mas \odot:

Mas \alpha Centauri C . Mul < . Mas \odot:

Rad \bigcirc .= Cca . Cmt 10<sup>100100</sup>:

Rad \bigcirc .= Cca . Cmt 10011 × 10<sup>11001</sup>:

Rad \bigcirc .= Cca . Cmt 101,0011 × 10<sup>11001</sup>:

Rtt \bigcirc .= Cca . Sec 10<sup>10101</sup>:

Rtt \bigcirc .= Cca . Sec 10<sup>10101</sup>:
```

¹⁾ The use of the same word Ctr as in 4204 is not correct.

Numerical data on the celestial bodies mentioned. The order of magnitude of masses, radii, and distances will be an indication for any receiver that 'Ste' means "celestial body".

We have restricted ourselves to a few instances. More numerical data of this kind should be given, especially on the planets. A receiver within our solar system might then possibly identify some bodies. A receiver outside our solar system will guess that the smaller bodies are members of our solar system.

```
* Pon: Cel \bigcirc . t = \text{Cmt/Sec } 0_{\text{Veo}}
    ^{\circ} Via x \cdot \text{Cca} \in \text{Ssc} \cdot \wedge : 1^{\circ} Foc ^{\circ} Via x \cdot = \text{Cca} \cdot \text{Ctr} \odot : \wedge : .
            A = {}^{"}x {}^{"}a \cdot \vee {}^{"}p \cdot e^{"} : {}^{"}Via x \cdot Cca \in Ssc pe \cdot \wedge \cdot p = a(1-e^2)
                \Diamond \cup \Diamond \cup \mathbf{Etc} \cup \Psi \cup \mathbf{P} | \Lambda
            E =   Yx Ye :    Yp  Yia x . Cca \in   Ssc pe  Etc  
                \land x "x Etc. \rightarrow: A'x = .Ax/A \uparrow "\land "R'x = .\text{Rev }x/\text{Rev }\uparrow. \rightarrow:
                A = 1101,10011 \times 10^{101000} \cdot \land \cdot
                \text{Rev}_{\bullet}^{+} = \text{Sec} \, 1111000011011000100010101, 100000001 \, \text{Etc}_{\bullet}^{+}
                A' \ \ = 0.0110001100011 \ \text{Etc} \cdot \land \cdot
                E \, \forall . = 0.0011011100110 \, \text{Etc} \cdot \land \cdot
                R' \ \ = 0.001111101100001 \ \text{Etc} \cdot \land \cdot
                A' \circ . = 0.1011100100101 Etc · \wedge ·
                E \circ ... = 0.0000000111000 \,\mathrm{Etc} \cdot \wedge ...
                R' \circ = 0.100111010111111 \text{ Etc.} \land \bullet
                E \div . = 0.00000100010001 Etc *
  and so on.
```

The Kepler motion of planets, the astronomical unit $(A \, \pm)$, the sideral year (Rev \pm), mean distances from sun (A'x), mean periods (R'x), and eccentricities (Ex). Other elements should be added, yet the data of $4\,24\,3-4\,24\,4$ will suffice to identify the various bodies within our solar system.

4 24 5. * Pon: Cel \pm . $t = \text{Cmt/Sec } 0_{\text{Vco}}$

Erg 'Via (. Cca \in Ssc · \wedge : 1 Foc · Via (. = Cca . Ctr \uparrow :

 $\land \cdot \text{Rev} (\cdot = \text{Sec } 100100,0000010100001111011011001 \times 10^{10000 + 100000}))$

The motion of the moon. In the same way satellites of other planets could be dealt with.

4 24 6. $\# f \in t_0 \operatorname{Coo} \operatorname{Ago} \longleftrightarrow i$

't₀ Coo Aqo' means the equatorial coordinate system at (the moment) t₀. Here, and in 4 24 7-4 24 9 t₀ is a substitute for something that means the moment of the vernal equinox of the year 1950.

The equatorial coordinate system has been defined by the following properties:

The fixed stars are approximately at rest in it.

The centre of the Earth is the origin.

The first axis is directed towards the first point of Aries.¹) At t₀ the Sun moves in the direction of increasing second and third coordinates.

The third axis is the directed rotation axis of the Earth.

The conditions are redundant. They yield the supplementary information that the annual (apparent) revolution of the Sun and the daily revolution of the Earth take place in the same sense.

By the condition $f \in \text{Coo}^+$, which seems to be redundant, we have fixed the meaning of Coo^{+2}).

Of course it does not matter which coordinate system we use in our communications.

 $* \operatorname{Car} \cdot t_0 \operatorname{Coo} \operatorname{Aqo} \cdot = 1 *$

¹⁾ The true equinox of March 21, which is slightly different from the mean equinox of January 1, as used in catalogues.

²⁾ Coo+ cannot be defined by means of the right hand rule of electromagnetics without referring to geographical data (the terrestrial north pole). In our program text we offered cosmographical data. Any oriented pair of stars with their velocity vectors are sufficient to settle space orientation. Nowadays thanks to the non-invariance of parity we could define Coo+ without any reference to geographical or cosmographical data.

4 24 7. #
$$f \in t_0 \operatorname{Coo} \operatorname{Aqo} \cdot \wedge \cdot fx = \lceil \xi_1 \cdot \xi_2 \cdot \xi_3 \rceil \cdot \wedge \cdot \xi_1^2 + \xi_2^2 = \varrho^2 \cdot \wedge \cdot \varrho > 0 \longrightarrow \xi_1 = \varrho \operatorname{cos} \operatorname{Ras} x \cdot \wedge \cdot \xi_2 = \varrho \operatorname{sin} \operatorname{Ras} x \cdot \wedge \cdot \xi_3 = \varrho \operatorname{tg} \operatorname{Del} x \cdot \wedge \cdot 0 \le \operatorname{Ras} x < 2\pi \cdot \wedge \cdot - \frac{1}{2}\pi < \operatorname{Del} x < \frac{1}{2}\pi = 0$$

The words written Ras (fL rectascensio) and Dcl (fL declinatio) mean "right ascension" and "declination" respectively.

4 24 8. *
$$x \in \text{Spa.} \land f \in t_0 \text{Coo Aqo} \land (\omega_1 f x)^2 + (\omega_2 f x)^2 \neq 0 : \rightarrow :$$

$$x_{dcl} = \lceil -\sin \text{Ras } x \cdot \cos \text{Ras } x \cdot 0 \rceil \cdot \land \cdot$$

$$x_{ras} = \lceil -\cos \text{Ras } x \sin \text{Del } x \cdot -\sin \text{Ras } x \sin \text{Del } x \cdot \cos \text{Del } x \rceil \cdot \land \cdot$$

$$x_{rad} = (f x) / |f x| *$$

 x_{acl} and x_{ras} mean the tangential directions at x to the declination parallel and the hour circle respectively, taken in the sense of increasing right ascension and declination respectively (the directions called eastwards and northwards in the jargon of the astronomer who tacitly assumes himself to observe the southern sky).

The meaning of x_{rad} is evident.

```
* x_{dcl} \perp x_{ras} \cdot \wedge \cdot x_{dcl} \perp x_{rad} \cdot \wedge \cdot x_{ras} \perp x_{ras} \cdot |x_{dcl}| = |x_{ras}| = |x_{rad}| = 1 *
```

4 24 9. * Pon: Cel
$$\div$$
 . to \cdot = . Cmt/Sec 0_{Vco} : \wedge . \wedge . Set \cdot . \wedge :

Via s . to \cdot = x : \wedge . Del $x = \frac{1}{2}\pi \delta$. \wedge . Ras $x = 2\pi \lambda$. \wedge . $\alpha > 0$. \wedge . f · Cel s . to : = $10^{10000} \times [\alpha (\cos 2\pi \mu x_{ras} + \sin 2\pi \mu x_{del}) + \beta x_{rad}]$:

Erg

```
s = \alpha Centauri. \rightarrow: \lambda = 0,1001101111000. \land.

\delta = -0,1010110001101. \land.

\alpha = 100011. \land.

\mu = 0,110010000. \land.

\beta = -100110. \land:

s = \text{Munich } 15040. \rightarrow: \lambda = 0,10111111100011. \land.

\delta = 0,0000110011110. \land
```

$$\delta = 0.0000110011110 . \land .$$
 $\alpha = 10001001 . \land .$
 $\mu = 0.1111111010 . \land .$
 $\beta = -10100101 : \land :$

$$s = \text{Wolf } 359. \rightarrow \lambda = 0.0111010001010. \land.$$

$$\delta = 0.0001010011111. \land.$$

$$\alpha = 1010010. \land.$$

$$\mu = 0.101001110. \land.$$

$$\beta = 10100: \land:$$

```
s = \text{Luyten } 726 - 8 \rightarrow \lambda = 0.0001000100101 . \Lambda
                              \delta = -0.0011001111010. A.
                              \alpha = 101010. A.
                             \mu = 0.001000010 . \land .
                             \beta = 100100 : \land :
s = \text{Lalande } 21185. \rightarrow \lambda = 0.0111010101110. \land .
                                \delta = 0.0110011100000.
                               \alpha = 1011001.
                               \mu = 0.100001010. \wedge.
                               \beta = 10000011. A:
s = \text{Sirius} \rightarrow \lambda = 0.0100011110100 \land \lambda
                   \delta = -0.00101111111011.
                   \alpha = 11000.
                   \mu = 0.100100010.
                   \beta = -1100 . A:
s = \text{Procyon} \rightarrow \lambda = 0.0101000100110 \wedge \lambda
                       \delta = 0.0001001000010. A.
                       \alpha = 111111.
                       \mu = 0.100110000. \wedge.
                       \beta = -101.
s = \text{Altair.} \rightarrow \lambda = 0.1101001101000. \land
                    \delta = 0.0001100001000. \wedge.
                    \alpha = 11000. A.
                   \mu = 0.001001101.
                   \beta = -101000 *
```

 $2 \pi \lambda$ and $\frac{1}{2} \pi \delta$ are the right ascension and the declination of the star, α and β are the numerical values (in 2^{16} Cmt/Sec) of its cross and radial velocities, μ is the so-called position angle of its proper motion. λ and δ together with the distances (in 4 24 3) determine the position, and the last three data determine the velocity in space. On the strength of this information the receiver will be able to identify the objects mentioned and the position of the sender. The information is highly redundant.

```
4 25 1. * Ha \operatorname{Inq} Hb \cdot \operatorname{Utr} \cdot \operatorname{Cel} \odot .= .\operatorname{Cmt}/\operatorname{Sec} 0_{\operatorname{Vco}}:

\operatorname{Utr} \operatorname{Cel} \circ .= .\operatorname{Cmt}/\operatorname{Sec} 0_{\operatorname{Vco}}:

Hb \operatorname{Inq} Ha \cdot \operatorname{Acc} x .= \operatorname{Cca} .\operatorname{Cmt}/\operatorname{Sec}^2 0_{\operatorname{Vco}} \cdot \wedge .y \in \operatorname{Hom} . \wedge .t \in \operatorname{Tem} :

\rightarrow :\operatorname{Nnc} Hb \operatorname{Ccd} \cdot \operatorname{Cca} t .y \operatorname{Inq} Hb :\operatorname{Pon} \cdot \operatorname{Cel} x .= .\operatorname{Cmt}/\operatorname{Sec} 0_{\operatorname{Vco}} \cdot Ha \operatorname{Inq} Hb :\operatorname{Nnc} Hc \operatorname{Inq} Hb \cdot \operatorname{Cel} x = a . \wedge .\operatorname{Cel} y = b :

\wedge :\operatorname{Nnc} Hd \operatorname{Inq} Hb \cdot \operatorname{Cel} x .= .\operatorname{Cmt}/\operatorname{Sec} 0_{\operatorname{Vco}} \cdot A :

\rightarrow :\operatorname{Nes} \cdot \operatorname{Nnc} Hd \operatorname{Inq} Hb \cdot \operatorname{Cel} y = b - a :
```

```
Hb Inq Ha '|a|  |a| | |b-a| . Mul  |c| | |a| | |a|
```

4 26 2. *
$$t_1 \operatorname{Usq} t_2 X \operatorname{Mut} \operatorname{For} x \leftrightarrow t_1 \operatorname{Usq} t_2 X \operatorname{Cau} \cdot t_1 \operatorname{For} x \neq t_2 \operatorname{For} x = t_2 \operatorname{Fo$$

4 26 3.
$$x \in \text{Rig.} \leftrightarrow \text{!} \land \land \land t. X \Rightarrow \text{!} \land \forall y \cdot y = \text{Cor } x : \land \land \text{Dif.} \land \text{Cea } t. X \text{ Mut For } x : > \text{Vld I}$$

$$x \in \text{Lqd.} \leftrightarrow \text{!} \land \land \land t. X \Rightarrow \text{!} X \in \text{Hom.} \land . t. x \Rightarrow \text{Ext.}$$

$$\rightarrow \text{!} \lor y \cdot y = \text{Cor } x : \land \text{!} \text{Dif.} \land \text{Cea } t. X \text{ Mut For } x : < \text{Vld.}$$

$$\land \cdot \text{Dif.} \land \text{Cea } t. X \text{ Mut Vol } x : > \text{Vld.} \text{I}$$

$$x \in \text{Gas.} \leftrightarrow \text{!} \land \land \land t. X \Rightarrow \text{!} X \in \text{Hom.} \land . t. x \Rightarrow \text{Ext.}$$

$$\rightarrow \text{!} \lor y \cdot y = \text{Cor } x : \land \text{!} \text{Dif.} \land \text{Cea } t. X \text{ Mut Vol } x : < \text{Vld.} \text{I}$$

$$\text{Exg.} \land \text{Aer} \subseteq \text{Gas.} *$$

The words written Rig (fL rigidus = rigid), Lqd (fL liquidus = liquid), and Gas mean "rigid", "liquid" and "gaseous" respectively.

4 26 4. *
$$x \in .$$
 Rig $\cup Lqd \cdot \rightarrow \cdot Dns x .$ Mul $> .$ Gra/Cmt³ $10^{-1000} : x \in Gas . \rightarrow \cdot Dns x .$ Mul $< .$ Gra/Cmt³ $10^{-1000} *$

```
4 27 1.  \begin{array}{c} t_1 & \\ Ha \operatorname{Inq} Hb \overset{!}{:}? = : \operatorname{Dur} \overset{\vee}{\cdot} x . \operatorname{PAN} \operatorname{Fit} x \text{ } \\ Hb \operatorname{Inq} Ha . \operatorname{Sec} 100 \text{ } \\ Ha \operatorname{Inq} Hb \overset{!}{:}? u \overset{!}{:} \operatorname{PAN} Hb \operatorname{Cog} \overset{!}{:}? = : \operatorname{Dur} \overset{\vee}{\cdot} x . \operatorname{Nnc} \operatorname{Fit} x \overset{!}{\cdot} \operatorname{Mod} u \text{ } \\ Hb \operatorname{Inq} Ha \overset{!}{:} t_1 t_2 Hb \operatorname{Met} \overset{!}{:}? = : \operatorname{Dur} \overset{\vee}{\cdot} x . \operatorname{Nnc} \operatorname{Fit} x \text{ } \\ Ha \operatorname{Inq} Hb . \overset{!}{:}? \operatorname{Met} \text{ } \\ Hb \operatorname{Inq} Ha \overset{!}{:} \operatorname{PAN} Hb \operatorname{Ani} \overset{!}{:} \lceil z_1 z_2 \rceil \overset{!}{\cdot} \operatorname{Fre} z_1 = \mathfrak{w} . \wedge . \operatorname{Fre} z_2 = \mathfrak{w} . \\ \wedge : \operatorname{Ini} z_1 = \cdot \operatorname{Ini} \overset{\vee}{\cdot} x \operatorname{Etc} : \wedge : \operatorname{Ini} z_2 = \cdot \operatorname{Fin} \overset{\vee}{\cdot} x \operatorname{Etc} \overset{!}{\cdot} \\ \wedge : \operatorname{PAN} Hb \operatorname{Cpu} \overset{!}{:}? = \cdot \operatorname{Ini} z_2 - \operatorname{Ini} z_1 \overset{!}{\cdot} . \end{array}
```

```
w \operatorname{Uni} \cdot {}^{\uparrow}z \cdot \operatorname{Fre} z = \omega : \rightarrow \cdot w \in \cdot \operatorname{Sca} \operatorname{Sec} 1 :
\operatorname{Ecc} \dots \cdot {}^{\downarrow}y \cdot \operatorname{PAN} \operatorname{Fit} y \cdot \in \cdot \operatorname{Sca} \operatorname{Sec} 10^{-100} :
\vee A : w \operatorname{Uni} A \cdot \wedge \cdot \operatorname{Car} A \cdot \operatorname{Mul} > \cdot 1 :
\wedge : \wedge i : i \in \operatorname{Num} \cdot \wedge \cdot i < \operatorname{Car} A \cdot \wedge \cdot x (i-1)^{\operatorname{um}} \operatorname{Par} w \cdot \wedge \cdot y i^{\operatorname{um}} \operatorname{Par} :
\rightarrow : \operatorname{Ini} x \cdot \operatorname{Sec} 10^{-100} \cdot \operatorname{Ant} \cdot \operatorname{Ini} y :
\rightarrow \cdot w \in \operatorname{Sca} \operatorname{Sec} 10^{-100} :
K \in \operatorname{Hor} \cdot \wedge \cdot t K \operatorname{Ext} :
\leftrightarrow : \vee {}^{\vdash}w \cdot \sigma^{\lnot} : w \in \cdot \operatorname{Sca} \operatorname{Sec} \sigma \cdot \wedge : \operatorname{Pot} \cdot \operatorname{Cca} t \cdot K \operatorname{Cau} w :
\operatorname{PAN} Hb \operatorname{Met} : ? = : \operatorname{Dur} \cdot {}^{\checkmark}x \cdot \operatorname{Nnc} \operatorname{Fit} x \cdot \operatorname{Per} \cdot 1^{\epsilon} \operatorname{Hor} *
```

The word written Met (fL metitur = measures) means the verb "measures". The word written Sca (fL scala = scale) means "scales"; after 'Sca' the unit of the scales has been mentioned. The word written Hor (fL horologium = clock) means instruments that can produce time-scales. 'Met... Per. 1^{ϵ} Hor' means "measuring by means of a clock".

The time-signal in the beginning lasts 4 seconds. The sequence of time signals after 'Ecc' consists of peeps with intervals of a sixteenth of a second.

The w refers to the program clock (see 2024).

```
4 27 2. * Ha \operatorname{Ing} Hb:? = \operatorname{Dst} \operatorname{Ab} = a \operatorname{Ab} \in Ra \cdot A \cdot \operatorname{Dst} ab = \operatorname{Lon} Ra
                      Hb \operatorname{Ing} Ha : \operatorname{Dst} \cdot {}^{\vee}a {}^{\vee}b \operatorname{Etc} \cdot = \operatorname{Cmt} 11001001 I
                      Ha \operatorname{Inq} Hb : ? u \cdot \operatorname{PAN} Hb \operatorname{Cog} \cdot ? = \operatorname{Dst} \operatorname{Etc} \cdot \operatorname{Mod} u
                      Hb \operatorname{Inq} Ha \operatorname{PAN} Hb \operatorname{Met} : ? = \operatorname{Dst} \operatorname{Etc} : \operatorname{Per} \cdot 1^{\epsilon} \operatorname{Rgu} i
                      Ha \operatorname{Inq} Hb \cdot ? = \operatorname{Rgu} I
                      Hb \operatorname{Inq} Ha :: w \in \operatorname{Rig} . \wedge : \forall y \cdot \operatorname{Put} y \cdot y \operatorname{Par} w . \wedge .
                                 y \in \operatorname{Sca} \operatorname{Cmt} 1 \to w \in \operatorname{Rgu}
                           Exg. Pon A = : ^{\uparrow}x \cdot x \in \text{Rea}^3.
                                      \land . 0 \le \omega_1 x \le 10^{1001} . \land . 0 \le \omega_2 x \le 10 . \land . 0 \le \omega_3 x \le 10
                                 Pon n \in \mathbb{N} n \in \mathbb{N} n \in \mathbb{N} n \in \mathbb{N} n \in \mathbb{N}
                                   \rightarrow B_n = : ^x x : x \in \text{Rea}^3. A. n - 10^{-100} \le \omega_1 x \le n + 10^{-100}.
                                   \land .0 \le \omega_1 x \le 10^{1001} . \land .0 \le \omega_2 x \le 10 . \land .10 \le \omega_3 x \le 10 + 10^{-100}
                                 Pon C = : ^{\uparrow}x \cdot \forall n \cdot x \in B_n
                           \wedge "Yv " f | \operatorname{Loc} v = B_n : n^{\operatorname{um}} \operatorname{Par} y :
                                 y \in \operatorname{Sca} \operatorname{Cmt} 1.w
                      Ha \operatorname{Inq} Hb : u \cdot \operatorname{Pot} : \operatorname{Nnc} Hb \operatorname{Met} \cdot \operatorname{Per} \cdot 1^{\epsilon} \operatorname{Rgu} \cdot \operatorname{Mod} u
                      Hb Inq Ha t \cdot X Ani Utr : Nnc a Apu \forall v \cdot v \cdot n^{um} Par
                                           \forall y: y \in \operatorname{Sca} \operatorname{Cmt} 1.w:
                                     \wedge Nnc b Apu \forall v : v : n^{um} \operatorname{Par} \forall y \operatorname{Etc}
                                \rightarrow "X Sci: Utr · Dst a b . = Cca . Cmt |m-n|::
```

The word written Rgu (fL regula=ruler), means "ruler, measuring-staff". 'Sca' now appears in connection with a length-unit.

The word written Tdo means "theodolite". We have dropped its description which would be similar to that of the ruler. (Roughly speaking it is a thin rigid cylinder with a hole around its axis and a ring of teeth on its convex surface serving as an angular scale.)

4 27 3. * Dst.x.Sfi ♂ · Mul < . Rad ♂ : ∧: Dst.y.Sfi ♂ · Mul < . Rad ♂ · →: x Ifa y. ↔ · Loc x ⊂ Cvx. Loc y. Loc Ctr ♂ · ↔ . y Sua x *

The words written Ifa (fL infra = below) and Sua (fL supra = above) mean "below" and "above" respectively.

```
4 27 4. * w \operatorname{Apu.Sfi}_{\overline{c}} \cdot \wedge . w \in \operatorname{Rig.} \wedge . z \operatorname{Par} w . \wedge . \operatorname{Vol} z < \operatorname{Vld.}
\wedge \cdot \wedge t : t w \operatorname{Ext.} \rightarrow \cdot - \operatorname{Pot.} t z \operatorname{Mov}^*
\wedge \cdot \wedge \cdot t . v : v \operatorname{Par} w . \wedge . v \neq z \cdot \rightarrow \cdot \operatorname{Pot.} t v \operatorname{Mov}^*
\rightarrow : - \cdot \vee v \cdot v \operatorname{Par} w . \wedge . t v \operatorname{Mov}^*
\leftrightarrow \cdot t \cdot \operatorname{Ctr} w . \operatorname{Ifa} z : \vee : t \cdot \operatorname{Ctr} w . \operatorname{Sua} z *
```

```
4 27 5. # Ha \operatorname{Inq} Hb \cdot ? = .\operatorname{Mas} Ra \mid
Hb \operatorname{Inq} Ha \cdot \operatorname{Gra} 11 \times 10^{111} \mid
Ha \operatorname{Inq} Hb \cdot ? u : \operatorname{PAN} Hb \operatorname{Cog} \cdot ? = .\operatorname{Mas} Ra \cdot \operatorname{Mod} u \mid
Hb \operatorname{Inq} Ha : \operatorname{PAN} Hb \operatorname{Met} \cdot ? = .\operatorname{Mas} Ra \cdot \operatorname{Per} \cdot 1^{\epsilon} \operatorname{Lib} \mid
Ha \operatorname{Inq} Hb \cdot ? = \operatorname{Lib} \mid
Hb \operatorname{Inq} Ha \cdot \operatorname{Pon} : w \in \operatorname{Rig} \cdot \wedge : w \cdot \operatorname{Apu} \cdot \operatorname{Sfi} t :
\wedge \cdot z_0 \sim z_1 \sim z_2 \sim z_3 \operatorname{Par} w \cdot \wedge \cdot \operatorname{Vol} z_0 \sim z_1 \sim z_2 \sim z_3 < \operatorname{Vld} \cdot
\wedge : t \operatorname{Loc} z_0 \cdot \subset \cdot \operatorname{Cvx} \cdot t \operatorname{Loc} z_1 \cdot t \operatorname{Loc} z_2 : \wedge :
t \operatorname{Loc} z_0 \cdot \subset \cdot \operatorname{Cvx} \cdot t \operatorname{Loc} z_3 \cdot t \operatorname{Loc} \operatorname{Ctr} w :
```

```
\land : \land \vdash r \cdot s \vdash r \land s \in \text{Ret} \cdot \land \cdot t \text{Loc} z_0 \land t \text{Loc} z_1 \land t \text{Loc} z_2 \cdot \text{Cea} \subseteq \cdot r \cdot
           \wedge \cdot t \operatorname{Loc} z_0 \wedge t \operatorname{Loc} z_3 \cdot \operatorname{Cea} \subseteq s:
           \rightarrow r \text{Cca} \mid s
     \land \lor p : p \in \text{Pla} . \land \cdot \text{Ctr} \, \exists . \in p \cdot \land \lor \land t : t \, w \, \text{Ext}.
           \rightarrow · Loc z_0 \sim z_1 \sim z_2 \sim z_3 . Coa \subset p:
           \wedge Pot: \forall t \cdot t w \text{ Ext. } \wedge \cdot t z_0 \text{ Mov}
           \wedge ' \wedge t : t w \to \mathbf{Ext} \longrightarrow \mathbf{Pot} \cdot t z_1 \wedge z_2 \wedge z_3 \to \mathbf{Mov}
\operatorname{Erg} \cdot w \in \operatorname{Lib} \cdot
PAN Hb Met . ? = Mas Ra . Per w.
    \leftrightarrow_1 \lor \lceil u.z^{1}_1.z^{1}_2 \rceil \exists u \operatorname{Uni} w \operatorname{Ra} \operatorname{Rb} \land \cdot \operatorname{Ra} \operatorname{Apu} z^{1}_1. \land \cdot \operatorname{Rb} \operatorname{Apu} z^{1}_2. \land \cdot
                 z^{1} Par w \cdot \wedge \cdot z^{1} Par w \cdot
           \wedge : \wedge t : t u \to : Loc z_1 : C \cdot Cvx : t Loc z_1 : t Loc z_0 :
                  \wedge : t \operatorname{Loc} z^{1}_{2} \cdot \operatorname{C} \cdot \operatorname{Cvx} \cdot t \operatorname{Loc} z_{2} \cdot t \operatorname{Loc} z_{0} \cdot \wedge :
           PAN Hb Cau · Nnc · u Ext \sim z_3 Sua z_0: \wedge
           PAN Hb \operatorname{Cog} : \operatorname{Utr} \cdot \operatorname{Mas} Ra \times \operatorname{Dst} z_0 z_1^1.
                   = \operatorname{Mas} Rb \times \operatorname{Dst} z_0 z_2^1 \wedge :
           PAN Hb Met \cdot? = \lceil \text{Dst } z_0 z_1 \cdot \text{Dst } z_0 z_2 \rceil \cdot \land :
           MAN Hb Cog.? = Mas Rb : \land :
           PAN Hb Cpu. ? = Mas Ra *
```

The word written Lib (fL libra = balance) means "balance". 'Met Per . 1 Lib' means "weighs".

```
4 28 1. # Ha Ing Hb? = . Dst Hb Ra
                   t1 Hb Inq Ha. Cmt 1101111 × 1010010
                   Ha \operatorname{Inq} Hb \cdot ? u \cdot \operatorname{PAN} Hb \operatorname{Cog} \operatorname{Etc} \operatorname{Mod} u
                    Hb \operatorname{Inq} Ha : t_1 Hb \operatorname{Met} : ? = \operatorname{Dst} Hb Ra : \operatorname{Per} \cdot 1^{\bullet} \operatorname{Und} \operatorname{Ath} :
                        \vee \lceil w \cdot w' \cdot t \cdot t' \rceil \quad w \rightsquigarrow w' \in \text{Und Ath}
                             \wedge : t Hb \operatorname{Cau} \cdot \operatorname{Cea} t \cdot \operatorname{Adv} Ra \cdot w \operatorname{Mov} : \wedge \cdot t' w \operatorname{Apu} Ra.
                             \wedge : t' Ra Cau \cdot Cca t' \cdot Adv Hb \cdot w' Mov
                        Sec 10^{-1001} Pst. Hb Cau Etc. \forall w' Etc. Apu Hb:
                        \operatorname{Erg} \cdot t_1 \operatorname{Dst} Hb \operatorname{Ra} \cdot = \operatorname{Cea} \cdot \operatorname{Sec} 0, 1 \times 10^{-1001} \times c
                              =Cmt 11011111 \times 10<sup>10010</sup>
                    Ha \operatorname{Ing} Hb \cdot \operatorname{Utr} \cdot \operatorname{Nnc} Ra \operatorname{Mov} \cdot
                    Hb Inq Ha · Nnc. Dev Hb. Ra Mov ·
                         Qia 'PAN Loc Ra . C: Cvx · Loc Hb . Nnc Loc Ra ' \wedge:
                             Sec 10^{110} . Ant Nnc · Dst Hb Ra = Cmt 1101111 \times 10^{10010}:
                             \land \cdot \text{Nnc Dst } Hb \ Ra := \cdot \text{Cmt } 11011111 \times 10^{10011}.
                         Erg · Nnc Cel Med Ra = \text{Cea} \cdot \text{Cmt/Sec} = 11011111 \times 10^{1100}
```

'Met...Per.1 Und Ath' means "measuring by means of radar".

```
4 28 2. # Hb \operatorname{Inq} Ha \operatorname{I}
Pon : X 	ildar Y \in \operatorname{Hom} . \land : \lor \ulcorner a . b \urcorner : \land t . t . = \operatorname{Cca} . t_1 \cdot
\rightarrow : \operatorname{Cel} X . t \cdot = a : \land : \operatorname{Cel} Y . t \cdot = b : '
Erg : : \operatorname{Cca} t_1 . X \operatorname{Inq} X . t = \operatorname{Nnc} .
```

 \rightarrow 'Cel X . t ' = . Cmt/Sec 0_{Vco} : \land : Cel Y . t ' = y ' \rightarrow 1' $\lor x$ | . |x| = |y| .

As we wish to initiate the receiver into relativity, we are committed to a careful treatment of the time variables. In 'Cca $t_1 \cdot X$ Inq $X \cdot t = \text{Nnc Etc}$ ' the variable t_1 relates to Hb's clock, because Hb is the direct user of t_1 . The variable t, however, though belonging to Hb's discourse and bound through quantification over by Hb, is to be considered as being used by X and relating to X's clock. Finally 'Nnc' belongs to X in any respect. In 'Cca $t_1 \cdot X$ Ccd 'Cca $t \cdot Y$ Inq $Y \cdot t' = \text{Nnc Etc}$ ' new time variables 't'' and 'Nnc' occur; they relate to Y's clock. Many such examples can be found in the following paragraphs.

$4\ 28\ 3$. # $Hb \operatorname{Inq} Ha I$

Pon : $X ildar Y \in \text{Hom} . \land ! \lor \lceil a . b \rceil : \land t . t . = \text{Cca} . t_1 .$ $\to `\text{Cel } X . t \cdot = a : \land : \text{Cel } Y . t \cdot = b : .$ Erg : $\text{Cca} t_1 . X \text{ Inq } X : t = \text{Nnc} .$ $\to `\text{Cel } X . t \cdot = .\text{Cmt/Sec } 0_{\text{Vco}} : \land : \text{Cel } Y . t \cdot = y : .$ $\land `\text{Nnc. Dev } X . Y \text{ Mov } ! \to !$ $\text{Ncs. } \text{Cca} t_1 . X \text{ Ccd. } \text{Nnc} \text{ Hb Inq } \text{Ha} : \text{Cel Rel } X Y . t_1 \cdot = |y| *$

The word written 'Cel Rel' means relative velocity. If X and Y approach, it will be considered to be negative. In order to say this, we should repeat the last program text with 'Adv' instead

of 'Dev' and -|y| instead of |y|.

4 28 4. # $Hb \operatorname{Inq} Ha I$

```
\forall \rho : \text{Put } \rho : \wedge \lceil \beta : \gamma \rceil : \beta \curvearrowright \gamma \in \text{Pos.}
                           \rightarrow: Cca t_1. X Inq X. U \in \text{Sca Sec } \beta:
                           \rightarrow Nes: Cea t_1. X Ced · Nnc Y Inq Y. U' \in \text{Sea Sec } \beta \rho.
                                      \land : \operatorname{Cca} t_1 : Y \operatorname{Inq} Y : U' \in \operatorname{Sca} \operatorname{Sec} \gamma
                                               \rightarrow 'Nes: Cea t_1. Y Ced · Nne X Inq X. U'' \in \text{Sea Sec } \gamma \rho:
                           Eti Cea t_1 \cdot X Inq X \cdot U \in Sea Sec \beta.
                                     \rightarrow Nes: Cea t_1. X Ced · Nec X Inq X. U'' \in \text{Sea Sec } \beta \rho^{10}. \wedge
                  \operatorname{Cea} t_1 \cdot X \operatorname{Inq} X \cdot U \in \operatorname{Sea} \operatorname{Sec} 1 \cdot \wedge \cdot w \wedge w_1 \in W \cdot \wedge \cdot w' \wedge w_1' \in W' \cdot \wedge :
                                     w_1 \operatorname{Apu} X \cdot \operatorname{Sec} 1 \cdot \operatorname{Pst} \cdot w \operatorname{Apu} X : \wedge :
                                    w_1 \operatorname{Apu} Y \cdot t \operatorname{Pst} \cdot w \operatorname{Apu} Y : \wedge :
                                    w_1' Apu Y \cdot t Pst · w' Apu Y : \wedge :
                                    w_1' Apu X \cdot t' Pst \cdot w' Apu X:
                           \rightarrow: Nes 'Cea t_1. X Ced: t' - Sec 1 = 10 (t - Sec 1).
                                \wedge \cdot c (t - \operatorname{Sec} 1) = vt \cdot \wedge \cdot t = \operatorname{Sec} c/(c - v) \cdot \wedge \cdot t' = \operatorname{Sec} (c + v)/(c - v)
                                     \operatorname{Erg} U' \in \operatorname{Sea} \operatorname{Sec} c/(c-v). A. U'' \in \operatorname{Sea} \operatorname{Sec} (c+v)/(c-v).
                                               \wedge \cdot \rho = \sqrt{(c+v)/(c-v)} \cdot \wedge
                                     • Nnc Y Inq Y . U' \in \text{Sea Sec } \sqrt{(c+v)/(c-v)} :
                 Erg: Cea t_1. X Inq X. U' \in Sea Sec 1.
                           \rightarrow 'Nes: Cea t_1. X Ced · Nne Y Inq Y.
                                     U' \in \text{Sea Sec } \sqrt{(c+v)/(c-v)}/(c/(c-v)) = \text{Sea Sec } \sqrt{1-v^{10}/c^{10}} + c^{10}/(c^{10}) = \text{Sea Sec }
          Relativistic time dilatation has been proved in the most
elementary and direct way.
                 Pon X \curvearrowright Y \in \text{Hom} \land \land \delta . \text{Mul} > . \text{Sec } 1 \cdot \land \exists
                           \forall \lceil a \cdot b \rceil : \land t : t_1 - \delta \leq t \leq t_1 + \delta . \rightarrow
                                     \operatorname{Cel} X \cdot t = a : \wedge : \operatorname{Cel} Y \cdot t = b : \wedge : \operatorname{Cel} \operatorname{Rel} X Y \cdot t = v : \wedge
```

Hb Inq Ha I

```
\alpha \in \text{Pos.} \land . V \in \text{Sea Sec} \alpha . \land : \land t \cdot t_1 - \delta \leq t \leq t + \delta . \rightarrow .
           t V Apu Y ..
Erg : Cea t_1 . X Vul . V \in Sea Sec 1 .
   \rightarrow Nes: Cca t_1. X Ccd. Nnc Y Vul. V \in \text{Sea Sec}/(1-v^{10}/c^{10}).
\land: Cca t_1. Y Vul. V \in Sca Sec 1.
   \rightarrow Ncs: Cca t_1. Y Ccd · Nnc X Vul. V \in \text{Sca Sec } 1/\sqrt{1-v^{10}/c^{10}}
```

This is only a paraphrase of the result of the preceding dialogue.

```
4 28 5. # Hb \operatorname{Inq} Ha I
```

```
Pon X 	ilda Y \in \text{Hom} . \land . \alpha \in \text{Pos} . \land . A \in \text{Sca Cmt} \alpha.
     \land \land \land n \land 1 + n \in \text{Num} \land \land n \leq N : \rightarrow z_n \land n^{\text{um}} \text{ Par } A.
     \wedge z_N. Ultum Par A \cdot \wedge \delta. Mul > Sec 1.
     \wedge : \vee \lceil a \cdot b \rceil : \wedge t \cdot t_1 - \delta \leq t \leq t_1 + \delta.
          \rightarrow Cel X \cdot t = a : \land : Cel Y \cdot t = b : \land : Cel Rel X \cdot Y \cdot t = v :
          \land .tz_0 \text{Apu } X . \land : t \text{Loe } X . \subset : \text{Cvx} \cdot t \text{Loe } A . \cup .t \text{Loe } Y : .
```

Erg : Cca t_1 : X Vul : $A \in$ Sca Cmt 1. \rightarrow 'Nes: Cea t_1 . X Ced

```
Nnc Y Vul . A \in \text{Sea Cmt } 1/1 - v^{10}/c^{10} ::
                       Qia · Pon · w \in \text{Und Ath}.
                           Apu Y . Usd . Apu z_N \cdot w Mov :
                           \land \land \land n \land n + 1 \in \text{Num} \land \land n \leq N \rightarrow w_n \in \text{Und Ath}
                               Apu z_n. Usd. Apu Y \cdot w_n Mov
                           \wedge : \wedge t \cdot t \ w \ \mathrm{Apu} \ z_n . \longleftrightarrow . \ t w_n \ \mathrm{Apu} \ z_n :
                           \wedge : \wedge n : z_n \operatorname{Par} A \longrightarrow : u_n = : w_n \operatorname{Apu} z_0 : \wedge : u'_n = : w_n \operatorname{Apu} Y :
                           \wedge : U \operatorname{Uni} \cdot u_1 \cdot u_2 \cdot \operatorname{Etc} : \wedge : U' \operatorname{Uni} \cdot u_1' \cdot u_2' \cdot \operatorname{Etc} :
                       Erg: Nes · Cea t_1 · X Ced · U \in Sea Cmt 10 c^{-1}:
                           Erg: Nes · Cea t_1. Y Ced. U' \in \text{Sea Cmt } 10 \ c^{-1} \sqrt{(c+v)/(c-v)}
                               \land Cca t_1. Y Inq Y: \beta \in \text{Pos}. \land A \in \text{Sca Cmt } \beta.
                                    \wedge \cdot w \operatorname{Apu} z_n \cdot t_n \operatorname{Pst} \cdot w \operatorname{Apu} Y \cdot \wedge \cdot
                                                                                                             1)
                                    w_n Apu Y \cdot t'_n Pst \cdot w Apu Y
                               \rightarrow Nes Cea t_1 Y Ced: t'_n = 10 t_n.
                               \wedge . Cmt \beta = (t_{n+1} - t_n)(c - v) . \wedge . Cmt 10 \beta = (t'_{n+1} - t'_n)(c - v)
                                    Erg. Cmt 10 \beta = 10 c^{-1} \sqrt{(c+v)/(c-v)} (c-v) \cdot \text{Erg.} \beta
                                         =\sqrt{1-v^{10}/c^{10}} #
                   A short proof of Lorentz contraction.
4 28 6. # Hb \operatorname{Inq} Ha
                      Pon X \curvearrowright Y \curvearrowright Z \in \text{Hom} \land \delta Mul > \text{Sec 1}
                           \land: Cel Y \cdot t \cdot = b: \land: Cel Z \cdot t \cdot = c: \land: Cel Rel X \cdot Y \cdot t \cdot = x:
                           \wedge: Cel Rel YZ \cdot t = y : \wedge: Cel XZ \cdot t = z.
                           \wedge "t \operatorname{Loc} Y \cdot \subset : \operatorname{Cvx} \cdot t \operatorname{Loc} X \cdot \cup : t \operatorname{Loc} Z :
                      Erg. z = (x+y)/(1+xy/c^{10}):
                       Qia : Pon : W \subset Und Ath . \land . \alpha \land \beta \land \gamma \in Pos.
                           \land. U \in \operatorname{Sca} \operatorname{Sec} \alpha. \land. U' \in \operatorname{Sca} \operatorname{Sec} \beta. \land. U'' \in \operatorname{Sca} \operatorname{Sec} \gamma.
                           \land \vdots \land \ulcorner u . u' . u'' \urcorner \vdots u \operatorname{Par} U . \land . u' \operatorname{Par} U' . \land . u'' \operatorname{Par} U'' 
                              \leftrightarrow: \forall \lceil w \cdot t \cdot t'' \rceil " w \in W \cdot \land : t \operatorname{Usq} t'' \cdot \operatorname{Apu} X \cdot \operatorname{Usd} .
                                    ApuZ \cdot w Mov:
```

 \rightarrow Nes Cca t_1 X Ccd Nnc Y Vul $U' \in$ Sca Sec

Erg • Cca t_1 . X Vul • $U \in$ Sca Sec 1:

 \land Cca t_1 . Y Vul· $U' \in$ Sca Sec 1:

 $\wedge \cdot u = \cdot w \operatorname{Apu} X \cdot \wedge \cdot u' = \cdot w \operatorname{Apu} Y \cdot \wedge \cdot u'' = w \operatorname{Apu} Z \cdot \cdot$

 $\sqrt{(c+x)/(c-x)}$ \cdot Nnc Z Vul. $U'' \in \operatorname{Sea} \operatorname{Sec} \sqrt{(c+z)/(c-z)}$:

¹⁾ Inadvertently I introduced a t_n , which for n=1 does not mean the explicit t_1 .

```
SPACE, MOTION, MASS
                                                                                                CH. IV
                        \rightarrow 'Ncs: Cca t_1. Y Ccd · Nnc Z Vul. U'' \in Sca Sec
                                                                                    \sqrt{(c+y)/(c-y)}:
                    \land: Cca t_1. Y Vul· U' \in \text{Sca Sec } \frac{1}{(c+x)/(c-x)}:
                        \rightarrow 'Ncs: Cca t_1 Y Ccd · Nnc Z Vul. U'' \in Sca Sec
                           \sqrt{(c+x)/(c-x)} \times \sqrt{(c+y)/(c-y)}
                 Erg: Cca t_1 \cdot X Vul · U \in Sca Sec 1:
                     \rightarrow 'Nes: Cea t_1. X Ced · Nnc Z Vul. U'' \in Sea Sec
                       \sqrt{(c+x)/(c-x)} \times \sqrt{(c+y)/(c-y)}:
                 Erg. \sqrt{(c+z)/(c-z)} = \sqrt{(c+x)/(c-x)} \sqrt{(c+y)/(c-y)}:
                 Erg. z = (x+y)/(1+xy/c^{10}) *
              An easy proof of the relativistic addition-theorem of velocities.
4 28 7. After this behaviouristic introduction into relativity, Lorentz
           transformations should be derived. We omit those purely mathe-
           matical developments.
4 28 8. # Hb Inq Ha :
                 \forall \sigma : Put \sigma : \sigma \in \uparrow v : -c < v < c : \frown Pos \land i
                     \land \ulcorner X.Y.z.t.a.v.\mu \urcorner \exists X \land Y \in \text{Hom} . \land .z \in \text{Res} . \land .z \text{Apu} X.
                        \wedge : tX \text{ Vul}'
                                       \operatorname{Cel} z \cdot \operatorname{Nnc} \cdot = \operatorname{Cmt/Sec} 0_{\operatorname{Veo}} : \wedge :
                                       Cel Y. Nnc = a : \land :
                                       Cel Rel X Y \cdot Nnc \cdot = v \cdot \wedge \cdot
                                       Mas z = \mu:
```

 \rightarrow Ncs: $t \times Ccd$ Nnc $Y \times Vul$: Cel Y. Nnc = Cmt/Sec 0_{Vco} . \wedge . Mas $z = \mu \sigma(v)$ $Ha \operatorname{Inq} Hb \cdot ? = \sigma$ $Hb \operatorname{Inq} Ha \cdot \sigma(v) = 1/\sqrt{1-v^{10}/c^{10}}$ Ha Inq Hb Cur! $Hb \operatorname{Inq} Ha :$ Pon $x \land z \in \text{Res} \land \text{Nnc} Hb \text{ Vul}$ $\operatorname{Cel} Hb \cdot \operatorname{Nnc} := \operatorname{Cmt}/\operatorname{Sec} 0_{\operatorname{Veo}} : \wedge :$ $\operatorname{Cel} x \cdot \operatorname{Nnc} = -\operatorname{Cel} z \cdot \operatorname{Nnc} = -a : \wedge :$ Cel Rel Hbx. Nnc = Cel Rel Hbz. Nnc = $v: \land$ $\operatorname{Mas} x = \operatorname{Mas} z = \mu \, \sigma(v) \cdot \wedge \cdot$ $y \operatorname{Uni} xz$ Erg . Nes : Nne Hb Ced: $\operatorname{Cel} y \cdot \operatorname{Nnc} = \operatorname{Cmt}/\operatorname{Sec} 0_{\operatorname{Veo}} \cdot \wedge \cdot$ Mas $y = 10 \mu \sigma(v)$. \wedge Nnc Ha Vul: $\operatorname{Cel} x \cdot \operatorname{Nnc} \cdot = \operatorname{Cmt}/\operatorname{Sec} 0_{\operatorname{Veo}} : \wedge :$ Cel z. Nnc = $10 a/(1+v^{10}/c^{10})$: A: $Cel y . Nnc = a : \land .$

```
\begin{aligned} \operatorname{Mas} x &= \mu \cdot \wedge \cdot \\ \operatorname{Mas} z &= \mu \, \sigma \, (10 \, v / (1 + v^{10} / c^{10})) \cdot \wedge \cdot \\ \operatorname{Mas} y &= 10 \, \mu \, \sigma \, (v)^{10} \cdot \wedge \cdot \\ \operatorname{Mas} y &= \operatorname{Mas} x + \operatorname{Mas} z \cdot \wedge \cdot \\ \operatorname{Mas} y &\times \operatorname{Cel} y &= \operatorname{Mas} x \times \operatorname{Cel} x + \operatorname{Mas} z \times \operatorname{Cel} z \end{aligned} \operatorname{Erg} \ \ 10 \, \sigma \, (v)^{10} &= 1 + \sigma \, (10 \, v / (1 + v^{10} / c^{10})) \cdot \wedge \cdot \\ 10 \, \sigma \, (v)^{10} \, a &= \sigma \, (10 \, v / (1 + v^{10} / c^{10})) \times 10 \, a / (1 + v^{10} / c^{10}) : \\ \operatorname{Erg} : \sigma \, (v)^{10} \, (1 + v^{10} / c^{10}) &= 10 \, \sigma \, (v)^{10} - 1 \ . \end{aligned} \operatorname{Erg} : \sigma \, (v)^{10} \, (1 + v^{10} / c^{10}) &= 10 \, \sigma \, (v)^{10} - 1 \ . \end{aligned}
```

Proofs given in textbooks of relativistic mechanics for the mass increase formula appeared on closer inspection to be sham proofs. This led me to believe that a mechanical proof (not using electromagnetics or quantum theory) was impossible. In rational relativistics the movement after an elastic collision is indeterminate 1) — I thought this was the reason why the mass increase formula could not be proved. This is not correct because the scalar features of a collision are determinate, and mass is a scalar. A proof was found, and finally it was even possible to eliminate collision. The proof is extremely simple now. The only essential assumption we have made is that the classical procedure of calculating the mass and the velocity of the union of two material points holds true in every inertial system, in other words that the four-impulses are additive (we have in practicue sed it only in a very special case.)

We have omitted the proof that the function σ defined by $\sigma(v) = 1/\sqrt{1-v^{10}/c^{10}}$ satisfies the requirement that the relational system

$$\sum_{i=1}^{3} \mu_{i} \sigma(x_{i}) = 0, \quad \sum_{i=1}^{3} \mu_{i} \sigma(x_{i}) x_{i} = 0$$

remains valid if the velocities x_i are simultaneously replaced by $(x_i+v)/(1+x_i v/c^{10})$. This can easily be proved by using the transformation law of the expression (c+x)/(c-x) which gets the factor (c+v)/(c-v). From the given relations we can derive

$$\sum_{1}^{3} \mu_{i} \sqrt{c^{10} - x_{i}^{10}}^{-1} (c \pm x_{i}) = 0$$

which appears to multiply with $\sqrt{(c \pm v)/(c \mp v)}$.

¹⁾ In reality classical and relativistic mechanics do not diverge from each other in this respect. In classical mechanics the result of an elastic collision cannot be predicted from laws of conversation. But it can be as soon as we add the postulates: 1° the solution is unique, 2° it is Galileo-invariant, 3° by the collision the velocities are changed (in order to exclude the "trivial" solution). The case of relativity is not different. Galileo-invariance is superseded by Lorentz-invariance.

SUMMARY OF CHAP. I-IV

```
1011-1015 Ostensive numerals. >, <, =, +, -.
1021-1022 Algorithmic numerals.
1031-1032 Numerical relations.
1041-1044 Variables.
1050 - 1052 Punctuation.
1.061 - 1.069 \rightarrow.
1 07 1
1 07 2
              Representation of tautologies.
1081 - 1082 \neq .
1091 - 1093 v.
1101-1102 \geq \leq \leq
1111-1114 \wedge ...
1 12 1
              2
1 13 1 - 1 13 9 Negative numbers.
1141-1145 Multiplication.
1151-1157 Fractions.
1 16 1 - 1 16 3 Dyadic fractions.
              Periodic fractions. Etc.
1 16 4
1165-1167 Etc. Powers. Infinite series.
1171-1173 The first names of sets. Num. Int. \in.
1181-1184 ∧-quantification, as infinite conjunction. \notin.
1 185-1188 V-quantification, as infinite disjunction.
1191-1192 Further examples.
              Div. divides.
1 20 1
              Pri, prime numbers.
1 20 2
1211-1216 C. D by examples.
1211-1222 Examples of real (irrational) and complex numbers.
1 22 3 - 1 22 7 Rat, rational numbers.
1231-1233 Non-formal introduction of Rea (real numbers).
1234-1236 More formal definition of Rea.
1237-1239 Complex numbers.
1240-1241 The word Agg (set).
1 24 2 - 1 24 6 Examples.
              \lceil x \rceil, the set with one element x only.
1 24 7
1251-1254 C, D, =.
```

```
1 25 3
                 Complete induction.
1261 - 1263 \cup, \cap, \setminus
                 \lceil \rceil, the void set.
1 26 4
1 27 1 - 1 27 5 Car, cardinal number, by means of examples.
1 28 1 – 1 28 6 Pairs of elements, \lceil a \cdot b \rceil. The pair set \lceil A \cdot B \rceil.
1 28 7
                 Triples.
1 29 1
                 Simplified notation, \vee \lceil a \cdot b \cdot c \rceil.
                 Colloquial conjunctions and disjunctions \checkmark, \land.
1 29 2
1293-1295 \uparrow x \dots The set of all x with ...
                 The set of subsets.
1301
1\ 30\ 2-1\ 30\ 3 \bigcirc.
1304 - 1305 [].
1 31 1
                 \forall x \ldots, the x with \ldots
                 Function. A \sim B, the set of functions from A to B.
1321
                 \forall x \text{ Etc } A, Etc as a function of x, defined in A.
1322-1323 *\,^f, the set of values of the function f, and other notations.
1 32 4
                 Examples.
1 32 5
                 The universal projections \omega_1, \ \omega_2, \ \omega_3; \ \omega_i \lceil x_1 \cdot x_2 \cdot x_3 \rceil = x_i.
1 32 6
                 Composition of functions.
                 Formal definition of \operatorname{Car} A = \operatorname{Car} B.
1 32 7
                 Example.
1 32 8
                 f \mid A, restriction of f to A.
1 32 9
                 inf.
1 33 1
1 34 1
                 Agg ½ Ord, partially ordered set.
                 Agg Ord, ordered set.
                 Agg Ord Ded, Dedekind-ordered set.
1342
1 34 3
                 Gru, group.
                 Gru Abe, abelian group.
1344
1 34 5
                 Cam, field.
1 34 6
                 Cam Ord, ordered field.
                 Cam Ord Ded. Dedekind ordered field.
1 35 1
                 The category of some symbols.
1 35 2
                 The sum symbol.
1 35 3
                 Sequences.
1 35 4
                 Calculus.
1360-1365 Ver, true. Fal, false. Prp, proposition. Qus, question. Iud,
                 truth value. (Introduced by means of examples.)
1366-1369 Some logical relations.
```

2011-2012 Duration (Dur) and the time unit Sec are introduced by means of "time signals" accompanied by data specifying their length in seconds.

201- 201-	T' 1			
2013-2014	Likewise Nos, number of oscillations, and Fre, frequency, are introduced.			
201 =				
2 01 5	Relations between these functions. Operating concrete			
2 01 6	numbers.			
	Fin, end, and Ini, beginning, are shown.			
	Verbal descriptions of time signal events. Ecc, see. Ant, before, and Pst, after, are shown and described.			
2 02 3	Pre, the relation of immediate predecessor.			
2 02 4	A clock is installed and described. Mom n is the clocktime n.			
2 02 5	Tem, the "time axis". Various relations between time			
2000	functions.			
2026	The category of some functions.			
	Quoting events. t_1 Usq t_2 Fit p , from t_1 till t_2 p happened.			
	Explicit definition of Ant.			
2 03 9	Dating by means of Ant and Pst.			
2040 - 2046				
2050	speech).			
2 05 0	Naming and describing events.			
2000 2000	How behaviour can be shown. The actors Ha , Hb , Hc ,,			
3000-3000	the acts, Inq, says, the value judgements Ben, Mal, good,			
	bad. Ha Inq Hb p.			
3 01 1 - 3 01 3				
3014	Quoting linguistic acts. $t_1 t_2 Ha \text{ Inq } Hb p$.			
3 01 5	The problem of oblique speech.			
3016-3018				
3 02 0 - 3 02 4	Interrogative pronouns and adverbs are circumscribed by means of '?'.			
3 03 1	Enu, counting, is introduced by means of counting time			
3 03 1	signals.			
3 03 2	Counting prime numbers.			
3 03 3	$t_1 t_2$ Enu freed from the $t_1 t_2$.			
3 03 4	Ordinals, 1um, 2um,, Ultum (last).			
	Cpu, computing, introduced by means of examples.			
3 04 4	$t_1 t_2$ Cpu freed from $t_1 t_2$.			
3 04 5	Erg, thus.			
3 04 6	A cardinal has been obtained not by counting, but (Sed) by			
	computing.			
3 05 1	Dem, proving a theorem, has been shown.			
3 05 2	Put, namely, used to justify existentially-bound variables			
	as instances.			
3 05 3	Pon, put.			
5 50 6	, <u>F</u>			

3224

3225

200- 200-	Det and the season of Dem Suding introduced by seasoning		
3061-3062	Pet, searching, and Rep, finding, introduced by searching for a person who said something, for a certain prime number.		
2000	$t_1 t_2$ Pet, $t_1 t_2$ Rep freed from $t_1 t_2$.		
3 06 3			
3 07 1	Anl, analyzing a structure of events, by means of Par, part, Uni, union.		
3 07 2	Multiplicative numerals 1es, 2es,		
	Other examples of analyzing. Mod, in the manner. The		
2013-2019	interrogative adverb "how".		
3081-3083	Nomina actionis derived by means of the termination 10.		
3 09 1	The degree of exactness of analyses. Err, error.		
3101-3104	Rsp, responding, introduced by means of examples.		
3111-3114	Mut, changing. Sin, omitting. Add, adding. Ilo, in-		
	stead of.		
3 12 0 - 3 12 3	Cur, why. Qia, because. Ncs, necessary.		
3 13 0 - 3 13 9	Sci, knowing.		
3141-3147	Ani, perceiving.		
3148-3149	Knowing by inference.		
3 15 1	No perception of mathematical truth.		
3 15 2 3 15 5	Cog, recognizing.		
3 15 6	Oblique speech after Sci. Eti, even.		
3 16 1	Complicated patterns of knowing.		
3170-3176	Pseudo-interrogative mode. Inq!		
3181-3188	Examples of mispronouncing and misunderstanding. Itg,		
	understanding.		
3 18 9	Rejecting meaningless speech.		
3 19 1	Estimation of duration and frequency of time signals. Err		
	measures the observational error. Cca, about, nearly. Mul,		
	much. Pau, little.		
3 19 2 - 3 19 3	Estimation of irrational numbers, of cardinal numbers.		
3 19 4	Comparison of copies of noises.		
3 19 5	Estimation of times.		
3 20 1 – 3 20 2	Nnc, now.		
3206 - 3207	Dating by means of Pau Ant Nnc etc.		
3 21 1 - 3 21 4	Abbreviations for datings.		
3 22 1	An event of a certain duration cannot have taken place in		
	too short an interval. Pot, can.		
3 22 2	A person cannot pronounce two different things at the		
	same time.		
3 22 3	Future events cannot be perceived.		

Knowledge about past events may be acquired in different

Mathematical knowledge may be acquired in different ways.

3 22 6 – 3 22 8	There may be different reasons why somebody does not				
0.00	answer a question.				
3 23 1 - 3 23 3	Reasons why somebody can naswer a question, can prove				
	a theorem, can compute the solution of an equation.				
3 23 4	$Sei \leftrightarrow Pot Inq.$				
3 24 1 – 3 24 2	Imperceptible speaking.				
3 24 3	Events perceived by one person though not by another.				
3244 - 3246	Too short, < Nim. Too long, > Nim.				
3 25 1	Relations between Ncs and Pot.				
3 25 2	A punctuation convention,				
3 26 1	The fact that somebody did not witness a certain event				
	gives an estimation of his age. Aet, age.				
3 26 2	Ext, exists. t Aet x defined by $t - \cdot$ Ini. x Ext. There is a				
	$\operatorname{Fin} x \operatorname{Ext}$.				
3 26 3	Average mental development of humans.				
3 26 4	Animals cannot speak.				
3 26 5	Hom, human. Bes, animals.				
3 26 6 - 3 26 7	Demographic statistics of actual mankind.				
3 26 8	Number of animals.				
3 27 1	Short history of Fermat's theorem.				
3 28 1	n^{ϵ} Hom. n people. And so on. 1^{ϵ} as indefinite article.				
3 28 2 - 3 28 3	Ise, (speaking) to each other.				
3 28 4	Alt, alternating.				
	Reasons why a person does not answer.				
3 29 3	A person does not answer a question because he does not				
	wish to answer. Vul, wishes. The only witness of an event				
	can prevent other prople from knowing anything about the				
	event. Vul.				
3 29 4	A long talk. One of the actors complies with all wishes of				
0 20 2	another actor. Other persons are less obedient. Ccd,				
	allowing.				
3 29 5	Relations between Vul and Ccd.				
3 29 6	One actor asks another to relay a message to a third actor				
3 2 9 0	who cannot be reached by the first directly.				
3 29 7	· · · · · · · · · · · · · · · · · · ·				
J 28 1	An actor refuses to relay a message, but he agrees when he is asked more politely through another actor.				
3 29 8	The state of the s				
3 28 8	Understanding (Itg) defined as recognizing what the other				
2 00 0	want to say.				
3 29 9	Somebody knows something if he says it whenever he wants				
2 00 1	to say it.				
3 30 1	An actor agrees to give information about a certain event				
	on condition that the other will keep the information secret.				

214

	When asked about this event the other refuses to give the		
	information because he has promised secrecy. Pol, promising.		
	Sat, enough.		
3 30 3	Explicit definition of Pol.		
3 30 4	Two actors agree to exchange information. The obligations		
	are honoured.		
3 30 5	One of the actors at first refuses to honour his obligation,		
	but later does so.		
3 31 1	Retroactive and retarded implications.		
3 32 1	Deb, ought to, introduced by comments on earlier talks.		
3 32 2	Lic, is allowed, similarly introduced.		
3 32 3	Relations between Deb and Lic.		
3 32 4	Wishing a thing not allowed is bad.		
3 32 5	Talk related to 3 29 4. Ha tells Hc that a has happened		
	and asks Hc to tell other people that b has happened. Hc		
	though always obedient to Ha, refuses, because this is not		
	allowed. Afterwards it seems that really b had happened.		
	Hd does not believe this explanation.		
$3\ 32\ 6 - 3\ 32\ 8$	Prompting is not allowed.		
3 32 9	Hb refuses to give some information. Ha threatens he will		
	tell Hc that some time ago Hb has betrayed some secret to		
	Ha though Hb had promised Hc not to tell it. Hd states		
	that Ha is doing something that is not allowed.		
3 33 1 - 3 33 3	How to prove by circumstantial evidence that Ha wishes		
	something though he does not use the word "I wish". It		
	is possible that somebody asks a question, though he knows		
	the answer. It may be his intention to check whether the		
	other knows the answer or even only whether the other is		
	listening.		
3 33 4	Uttering one's wishes in a more circumstantial manner.		
3 34 1	This may be done for reasons of decency. Dec, Plt are two		
	levels of decency.		
3342 - 3345	Examples of polite speech.		
3346 - 3347	Examples of good breeding: not to listen to confidential		
	talks of other persons, not to tell about confidential talks		
	of others which one has happened to witness.		
3 34 8	Relations between Dec and Plt.		
3 34 9	Relations between Ncs, Deb, Plt, Dec, Lic, Pot.		
3 35 1 – 3 35 7	General reflections about modality.		
3 36 1	A talk about ethics. What is good and who are good people.		
3371 - 3372	Dif, degree of difficulty, is introduced by means of many		

examples of actions whose difficulty is compared.

3 37 3 - 3 37 4	>Vld, <vld, large,="" small.<="" th="" very=""></vld,>			
3 37 5	It is very difficult to know something about an event, if			
	the only witness refuses to tell anything about it. It is still			
	more difficult if nobody witnessed it.			
3 37 6	It is very difficult to forget something especially if somebody			
	asks you expressly to forget it.			
3 38 0	A more solid general background of human wishes is needed.			
3 38 1	A bet on solving a cubic equation. The winner refuses to			
	tell the general method of solution. The loser refuses to pay			
	up. An arbiter is called in. The winner tells his method, the			
	loser pays. Dat, giving. Den, penny.			
3 38 2	Pec, fortune. How one's fortune changes under the influence			
	of payments.			
3 38 3	This talk is related to 3 32 5 and 3 29 4. It appears that			
	the obedient servant is regularly paid by his master.			
	Another person is disobedient, because the master did not			
	pay him his due.			
3 39 1 – 3 39 3	Regular dialogues, e.g. one actor pronounces a number,			
	and the other pronounces its square. A third person states			
•	the law, Lex, of the dialogue.			
3 40 1	Two persons play a simple arithmetical game (adding a			
	number ≤ 8 ; he who reaches 32, wins).			
	Lud, playing, Vin, winning. Prd, losing. Ict, move (of a			
2 40 0	game).			
3 40 2	General definition of game.			
3 40 3	The same game played by other players. Good and bad			
	strategies. The game is unfair, because the beginner can			
2 40 4	always win. Ius, fair.			
3 40 4	Another game. One player chooses one of the numbers 1 2, 3, 4, and the other tries to guess it. His chance of winning			
	is only 1/4, but the stake of the first player is 3 against 1 of			
	the second. So the game is fair. Prb, probability.			
3 40 5 - 3 40 6	General notions about probability, expectation (Uti).			
3407	One player chooses one of the numbers 1, 2, 3, 4, and the			
3 40 7	other tries to guess it. If he hits the chosen number, the			
	first player pays as many pennies as the chosen number			
	has units. If he fails, nothing is paid. The players discuss			
	the stake. Minimax principle.			
3 40 8	A still more complicated two-person game.			
3 40 9	A three-person game. Aux, aiding. Soc, allies.			
3 41 1	A paradox: I do not say anything.			
3 41 2	The paradox of the liar.			

216 SUMMARY

4 01 1	Difference of place is recognized by delay of signals. Loc,		
	place. Spa, space, the union of all places. Dst, distance,		
	proportional with the time of delay. Cmt, the length unit.		
4 01 2	Cmt as a linear function from reals to concrete reals.		
4 01 3	Axioms of metric space.		
4 01 4	Special properties of the metrics of Euclidean space.		
4 01 5	Convex envelope, Cvx.		
4 01 6	Ret, straight lines.		
4 01 7	Pla, planes.		
4 01 8	//, parallel.		
4 02 1	Sph, spheres. Bul, solid spheres.		
4 02 2 - 4 02 3	Some geometrical properties of Euclidean space. \perp ,		
	orthogonal.		
4 02 4	Reflections, Rfc.		
4 03 1	Barycentre $\alpha a + \beta b$ of two points a and b with $\alpha + \beta = 1$.		
4 03 2	Vector $b-a$ of two points a , b , element of Cmt Vco, centi-		
	meter vector space.		
4 03 3	$a-b=c-d$ means $\frac{1}{2}(a+d)=\frac{1}{2}(b+c)$. Parallelogram. $\#$.		
4 03 4 - 4 03 7	Algebraic operations with points and vectors.		
4 04 1	$Rea^3 = \lceil Rea. Rea. Rea \rceil.$		
4 04 2	Algebraic operations in Rea3. alb, scalar product. \triangleleft ,		
	angle, \perp .		
4 04 3	Rct _* , Pla _* , straight lines and planes in Rea ³ .		
4 04 4	$f \in \text{Coo means a coordinatization of Spa by means of Rea}^3$.		
4 04 5	In a coordinatization Pla*, Rct*, \(\preceq\) and so on correspond		
	to Pla, Ret, ⊥.		
4 04 6	In a coordinatization $\alpha a + \beta b$, and $a - b$ are reflected by		
	formally the same expressions.		
4 04 7	Cmt ² .		
4 04 8	In a coordinatization operations with vectors are reflected		
	by formally the same operations.		
4 04 9	Positive and negative coordinate systems, Coo+, Coo		
$4\ 05\ 1 - 4\ 05\ 3$	Transition from Cmt Vco to Vco (vector space).		
4 06 1	Lap, rectangular parallelepiped.		
4 06 2	Its volume.		
4 06 3	Cmt ³ .		
4 06 4	General definition of volume.		
4 06 5	π . Volumes of solid spheres.		
4 07 1	Volumes of humans and animals.		
4 07 2	Lon, diameter, of sets.		
4 07 3	Diameters of humans and animals.		
4081-4085	Mov, moving Usd, from to		

4 09 1	One actor calls another. He moves. Apu, at, near.		
4 09 2	In too short a time a person cannot go from one place to		
	another.		
4 09 3	Hb refuses to walk to Ha. Hb wants Ha to walk to him.		
4 09 4	Ha wants Hb to fetch Hc and to walk with him to Ha .		
	They obey.		
4 10 1	Whistling for one's dog.		
4 10 2	The dog refuses. He persues the fleeing dog. Adv, towards.		
	Dev, away.		
4 10 3	Number of animals which can perceive, wish and move;		
	much larger than that of humans.		
4 10 4	One actor perceives another, he calls him, they go to meet.		
4 10 5	A dog joins his master in order to please him.		
4111	Ha was mistaken, he called to a thing, thinking that it was		
	a human or an animal. Things cannot move themselves.		
	Now this thing is moved by Hd towards Ha. Humans can		
	cause things to move. Res, things. Cau, causing.		
4 11 2	Things are in space.		
4113	Fer, carrying.		
4 11 4	Iac, throwing.		
4115	If an object is too large it cannot be carried by a single		
	person, though perhaps by more persons together.		
4116	An object thrown and intercepted.		
4117	An actor tries to intercept a thrown object, but does not		
	succeed because the distance is too large.		
4118	Playing ball.		
4 12 1	Mean velocity.		
4 12 2	Sec-1.		
4 13 1	Harmonic oscillation of a small particle. Osc, oscillating.		
4 13 2 - 4 13 3			
4.7.5	equilibrium of an oscillation.		
4 13 4	Vib, an object vibrates, if its parts oscillate. Examples for		
474-	stationary and progressive waves.		
4141	Movements of objects caused by humans are rather slow.		
4 14 2	Why can a human or an animal change its place? Because		
	he want to move. Why can a thing change its place? Because it moves.		
4140			
4 14 3	Why can a person receive a message sent from another place? There are three ways. The first, through mere convection of		
	the message, is rather slow. Wave propagation is much		
	quicker. There is an undulating medium, Aer, characterized		
	by sonic velocities. Finally there is wave propagation through		
	~, como, crocidos, a mais dicionos wave propagament and a more		

4 20 3

	an uncorporal medium, Ath, characterized by a much				
	higher velocity, that of light. Uul, undulating, Und, wave.				
4144	Exposition of the underlying ideas.				
4145	On the strength of 4 14 3 the length unit can be calculated.				
4 14 6	Hydrogen spectrum and Rydberg's constant. One of the				
	actors states that this yields more precise knowledge about				
	the length unit than the velocity of light can.				
4181					
	by a single person though its dimensions are not too large.				
	Its mass is too large. The larger the mass of an object, the				
	more difficult it is to move. Mas, mass. Gra, gramme.				
	Dns, density.				
4 18 2	Gra as a linear function. Gra/Cmt3. Gra Cmt/Sec. And so on.				
4 18 3	Dns $x = \text{Mas } x/\text{Vol } x$. Density of humans and animals.				
4184	Additivity of mass.				
4185	Humans, animals, things have a mass.				
4186	Humans, animals, things have a body, Cor, which may				
	depend on time. Place, volume, diameter, mass, density of				
	a body.				
4 18 7	A wave as a thing with a variable body.				
4188	The existence of a human body begins some time earlier				
	than that of the human itself. The same is true for some				
	animals. Mat, mother. Pat, father. Before the individual				
	existence of a human, its body is part of the body of its				
	mother. It has originated from a part of the body of its				
	mother and a part of the body of its father. Hom Msc,				
	males, Hom Fem, females. Number of human males and				
	females. Females can be mothers, males can be fathers.				
4 19 1	Cel, vector velocity for uniform motions.				
4 19 2	Equality of velocities.				
4 19 3	Cmt/Sec. Gra Cmt/Sec. And so on.				
4 19 4	Cmt/Sec Vco, space of velocity vectors.				
4 19 5	Cmt/Sec as a linear function.				
4 19 6	The image of Cmt/Sec Vco under coordinatization.				
4 19 7	Uniform motion with given velocity.				
4 19 8	Relation between Cel and Cel Med.				
4 20 1	Talk about an elastic collision of a large and a small body,				
	change of velocity. Ccu, colliding.				
4 20 2	Collision of bodies with equal masses.				

Elastic collision of arbitrary bodies. The law of elastic

collision. Lex Nat, law of nature. One of the actors states

that by this law the mass ratio is fixed.

SUMMARY 219

4 20 4	Ctr, mass centre. Ipu, impulse, Kin, kinetic energy.
4 21 1	Via, path.
4 21 2	Cel, velocity of arbitrary motions.
4 21 3	Acc, acceleration.
4 22 1	† earth. Its place is nearly a solid sphere. Rad, its radius.
	Its density.
4 22 2	Sfi, its surface.
4 22 3	Humans and animals are at the surface of the earth.
4 22 4	Talk about a falling body. Cad, falling. The law of free fall,
	a law of nature. The universal attraction law, a law of
	nature. The universal gravitation constant. One of the
	actors states that on the strength of these data he can
	compute the mass unit.
$4\ 22\ 5$	Mass of the hydrogen atom.
4 22 6	Addivity of gravitational accelerations.
4 23 1	Definition of conics. Ssc (half) conic. Foc, focus.
4 23 2	Conics as paths of Kepler motions. Rev, period of revolution.
4 23 3 – 4 23 5	Rot, rotating.
4 23 6	Rtt, period of uniform rotation.
4 23 7	Axs, axis of uniform rotation.
4 24 1	Names of sun, moon, planets and some fixed stars. Ste,
	celestial body.
4 24 2	Celestial bodies are nearly solid spheres. The centre is the
	mass centre. Surface. Radius.
4 24 3	Mass, radius, period of rotation and distances of some
	celestial bodies.
4244-4245	Kepler motion and elements of planets, and of the moon.
4 24 6 - 4 24 8	Equatorial coordinate system (Coo Aqo), right ascension
	(Ras), declination (Dcl), northwards, eastwards.
4 24 9	Coordinates and proper motion of some fixed stars.
4 25 1	Ha asks whether the sun or the earth are at rest. Hb allows
	any non-accelerated body to be said to be at rest. Ha tells
	how velocities transform if another body is taken to be at
	rest. Ha admits this to be nearly true for low velocities.
	Hb applies it to the velocity of light. Ha asserts that the
	velocity of light is always the same. Further explanations
	are postponed.
4 26 1	For, shape.
$4\ 26\ 2$	Change of shape.
4 26 3	Rig, solids - their shape cannot easily be changed. Lqd,
	liquids - their shape can easily be changed, but not their
	1 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1

volume. Gas, gases - their volume can easily be changed.

4 26 4	Density of solids, liquids, gases.		
4 27 1	One of the actors tells how time is measured. Met, measur-		
	ing. Sca, scale. Hor, clock. Met Per Hor, measuring		
	by means of a clock.		
4 27 2	One of the actors describes a measuring staff and measuring		
	distances by means of a measuring staff, and by trigono-		
	metrical means. Rgu, measuring staff. Tdo, theodolite.		
	MetPerRgu, measuring by means of a measuring staff.		
4 27 3	Ifa, below. Sua, above.		
4 27 4	Equilibrium of solid bodies.		
4 27 5	One of the actors describes a balance and weighing.		
4 28 1	One of the actors describes measuring distances and velo-		
	cities by means of radar. Met Per Und Ath, measur-		
	ing by means of radar.		
4 28 2	Hb says that if X and Y move uniformly and if X asserts		
	that he is at rest and that Y has the velocity y , X is comitted		
	to allowing that Y can assert that he is at rest and that X		
	has a velocity x which is absolutely equal to y .		
4 28 3	In this way the relative velocity, Cel Rel XY , of X and Y		
	is defined; if X and Y approach it will be considered to		
	be negative.		
4 28 4	Observation of moved clocks leads to a proof of relativistic		
	time dilatation.		
4 28 5	Observation of measuring staffs by means of radar leads		
	to a proof of Lorentz contraction.		
4 28 6	The relativistic addition theorem of velocities.		
4 28 7	Lorentz transformations.		
4 28 8	Relativistic mass increase formula.		

REGISTER

#	1 00 0	∉	1 18 2
>	1 01 1	∉ ∨	1 18 5
<	1 01 2	n^{ϵ}	3 28 1
_	1 01 3	C)	1 21 1
+	1 01 4	C	1 21 1
_	1 01 5	$oldsymbol{i'}{oldsymbol{i}}$	1 22 1
1	1 02 1	$oldsymbol{i}$	1 22 1
10	1 02 1	π	$1\ 22\ 2, 4\ 06\ 5$
11	1 02 1	$oldsymbol{e}$	$1 \ 22 \ 2$
a, b, c, \ldots		Г٦	1 24 7, 1 26 4
x, y, z, \ldots		U	1 26 1
(variables)	1041 - 1044	\cap	1 26 1
·,:)		\	1 26 1
i, I }	1 05 0	×	1 27 4
<i>,</i> ,	1 06 1, 3 31 1	Γ. ٦	1 28 2
\leftrightarrow	1 07 1	~	$1\ 29\ 2$
≠	1 08 1	^	1 29 2
	1 09 1	${}^{ au}\!x$	1 29 3
∨ ≦ ≧	1 10 1	\cap	1 30 2
_ ≥	1 10 1	ij	1 30 4
^	1 11 1	$\overset{\smile}{\forall}_x$	1 31 1, 3 01 3
š	1 12 1, 3 02 1, 3 01 2,	$^{\forall}x$	1 32 1
	3 01 3		1 32 1, 1 32 9
0	1 13 2	~	1 32 1, 1 32 2
-1	1 13 2	*	1 32 2
10	1 13 2	f^	1 32 3
×	1 14 1	.^ <i>f</i>	1 32 3
a/b	1 15 2	$\omega_1, \omega_2, \omega_3$	1 32 5
1,1	1 16 1	Σ	1 35 2
0,1	1 16 2	_	1 36 0, 1 36 6
1,101 101	1 16 3		2 01 1
a^n	1 16 6	a, b, c,	2 01 1
€	1 17 1	1	3 17 0
٨	1 18 1	//	4 01 9
		"	

	4 02 2	Bes	3 26 5
I	4 04 2, 4 04 8	Bul	4 02 1
⋠	4 04 2, 4 04 8	Cad	4 22 4
	4 03 1, 4 04 2	Cam	1 34 5
$0_{ m Vco}$	4 03 6	Cam Ord	1 34 6
0_3	4 04 2	Cam Ord Ded	1 34 6
c	4 14 3	Car	1 27 1, 1 32 7
t	4 22 1	Cau	4111
	4 24 1	Cca	3 19 1
¥	4 24 1	Ced	3 29 4
우	4 24 1	Ceu	4 20 1
⊙ \$\phi\$ 2\	4 24 1	Cel	4191, 4212
24	4 24 1	Cel Med	4 12 1
ħ	4 24 1	Cel Rel	4 28 3
₫	4 24 1	\mathbf{Cmt}	4 01 1
Ψ	4 24 1	Cmt Vco	4 03 2
В	4 24 1	Cog	3 15 2
(4 24 1	Com	1 23 1
A	4210 4220	Coo	4 04 4
Acc	4 21 3, 4 22 6	Coo+, Coo-	4 04 9
Add	3 11 3	Coo Aqo	4 24 6
Adv	4 10 2	Cor	4186
Aer	4 14 3	Cpu	3 04 1
Aet	3 26 1	$\overline{\operatorname{Ctr}}$	4 20 4, 4 23 7
Agg	1 24 1	Cur	3 12 0
Agg Ord	1 34 1	Cvx	4 01 5
Agg Ord Ded		ъ.	2.00
$Agg \frac{1}{2} Ord$	1 34 1	Dat	3 38 1
Alt	3 28 4	Del	4 24 7
Aml	4 13 2	del	4 24 8
Ani	3 14 1	Deb	3 32 1
Anl	3 07 1	Dec	3 34 1
Ant	2 02 1, 2 03 8, 2 03 9,	Dem	3 05 1
A -:	3 31 1	Den	3 38 1
Aoi	4 24 1	Dev	4 10 2
Apu	4 09 1	Dif	3 37 1
Ath	4 14 3	Div	1 20 1
Ato H	4 14 6	Dns	4 18 1
Aux	3 40 9	Dst	4 01 1
Axs	4 23 7	Dur	2 01 1
Ba, Bb, Bc, \dots	3 26 5	Eee	2 01 7
Ben	3 01 2	Enu	3 03 1

\mathbf{Eql}	4 13 2	Iud	1 36 0
Erg	3 04 5	\mathbf{Ius}	3 40 3
Err	3 09 1, 3 19 1	***	4
es	3 07 2	Kin	4 20 4
\mathbf{Etc}	1 16 4	Lap	4 06 1
$\mathbf{E}\mathbf{t}\mathbf{i}$	3 15 6	Lex	3 39 1
Exg	3 02 2	Lex Nat	4 20 3
Ext	3 26 2	Lib	4 27 5
		Lic	3 32 2
\mathbf{Fal}	1 36 0	Loc	4 01 1
Fem	4 18 8	Lon	4 07 2
Fer	4 11 3	Lon Lqd	4 26 3
\mathbf{Fin}	2 01 6		
\mathbf{F} it	2 03 1	Lud	3 40 1
${\it f}L$	1 17 1	Mal	3 01 3
Foc	4 23 I	MAN	3 21 1
\mathbf{For}	4 26 1	Mas	4181
\mathbf{Fre}	2 01 4, 4 13 2	Mat	4 18 8
\mathbf{Ftt}	3 18 1	Met	4 27 1
		Mod	3 07 4
Gas	4 26 3	Mod Ben	3 18 1
Gra	4 18 1	Mod Mal	3 18 1
Gru	1 34 3	Mom	2 02 4
Gru Abe	1 34 3	Mov	4 08 3
	0.01 -	MPN	3 21 1
Ha, Hb, Hc,		Msc	4 18 8
Hom	3 26 5	Mul	3191
\mathbf{Hor}	4 27 1	Mut	3 11 1
т	4.04.0	Mut	3111
I	4 04 2	Ncs	3 12 2
Iac	4 11 4	Nim	3 24 4, 3 24 5
Ict	3 40 1	Nne	3 20 1
Ifa	4 27 3	Nos	2 01 3
Ilo	3111	Num	1 17 1
Ini	2016		
inf -	1 33 1	\mathbf{Obv}	4 10 4
Inq	3 00 5	Osc	4 13 1
Inq!	3 17 0	T) 4 3 T	2 21 -
Int	1 17 1	PAN	3 21 1
10	3 08 1, 3 08 2, 3 08 3	Par	3 07 1
Ipu -	4 20 4	Pat	4188
Ise	3 28 2	Pau	3 19 1
Itg	3 18 1	\mathbf{Pec}	3 38 2

Dan	4 97 1	Dan	3 10 1
Per	4 27 1 3 06 1	$egin{array}{c} \mathbf{Rsp} \\ \mathbf{Rtt} \end{array}$	4 23 6
Pet		Nut	4 20 0
Pha Dla	4 13 2	Sat	3 30 2
Pla	4016	Sca	4 27 1
Pla _*	4 04 3 3 34 1	Sci	3 13 0
Plt	3 30 1	Sec	2 01 1
Pol Pon	3 05 3	Sed	3 01 3, 3 04 6
Pos	1 23 2	Sfi	4 22 2, 4 24 2
Pot	3 22 1	Sin	3112
PPN	3 21 1	Soc	3 40 9
Prb	3404	Spa	4 01 1
Prd	3 40 1	Sph	4 02 1
Pre	2 02 3	Ssc	4 23 1
Pri	1 20 2	Ste	4 24 1
	1360	Sua	4 27 3
Prp Pst	2022, 2038, 2039,	sup	1 33 1
ISU	3 31 1		
Put	3 05 2	Tan	3 01 3
Fut	3 03 2	Tdo	4 27 2
Qia	3 12 0	${f Tem}$	2 02 5
\mathbf{Qqm}	3 01 3		
\mathbf{QUA}	3 21 3	UAN	3 21 2
\mathbf{QUP}	3 21 4	$\mathbf{Ult^{um}}$	3 03 4
\mathbf{Qus}	1 36 0	um	3 03 4
Da Dh Da	. 4 11 1, 4 11 2	Und	4 14 3, 4 18 2
Rad	4 22 1, 4 24 2	Uni 	3 07 1
rad	4 24 8	Unq	3 21 2
Ras	4247	UPN	3 21 2
ras	4 24 8	Usd	4 08 4
Rat	1 22 3	Usq	2 03 1
Ret	4016	Uti	3 40 5
Ret*	4 04 3	Utr	3 02 1
Rea	1 23 1	Uul	4 14 3
Rea ³	4041	Veo	4 03 2
Rep	3 06 1	Ver	1360
Res	4111	Via	4 21 1
Rev	4 23 2	Vib	4134
Rfc	4 02 4	Vin	3 40 1
Rgu	4 27 2	VIII	3 37 3
Rig	4 26 3	Vid Vol	4 06 2, 4 06 4
Rot	4 23 3	Vol	3 29 3
		4 UI	J 20 0









